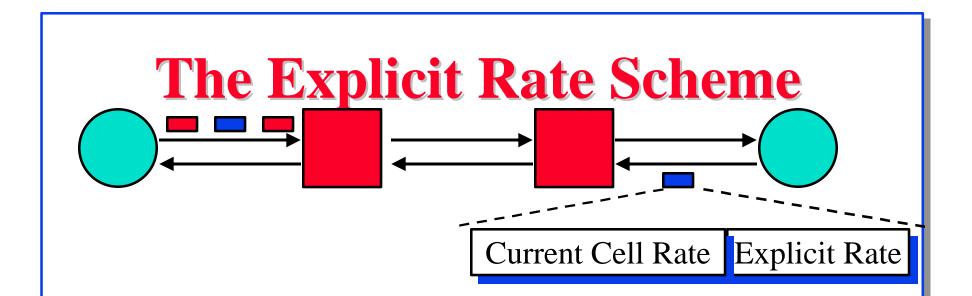




- q Presentation at ATM Forum
- q Modeling MPEG2 Transport Streams over VBR background
- q Virtual Source/Virtual Destination Design Analysis



- q Sources send one RM cell every n cells
- q The RM cells contain "Explicit rate"
- q Destination returns the RM cell to the source
- q The switches adjust the rate down
- q Source adjusts to the specified rate

1. ATM Forum Presentation

- q "Real-Time ABR: Proposal for a New Work Item," ATM Forum Contribution 96-1760, December 1996, <u>ftp://netlab.ohio-state.edu/pub/jain/atm96-1760.txt</u>
- q Contribution co-sponsored by Samsung and Lucent
 Technologies
- q AT&T seems to be working on it also
- q Accepted as a work item for Traffic Management
 V5.0

Video over ABR: How?

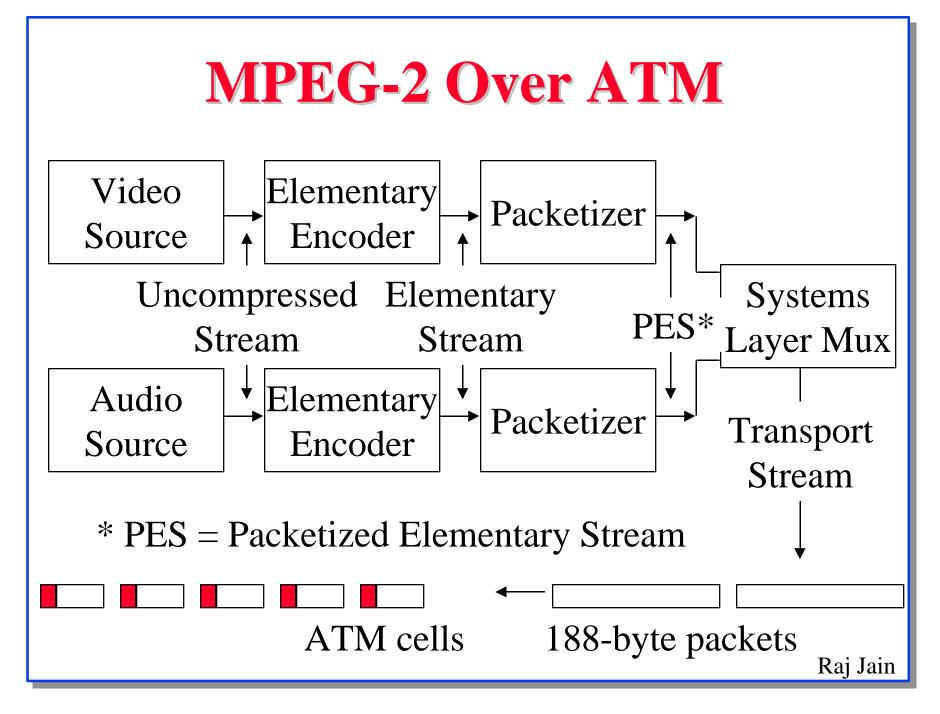
- q Compression parameters can be dynamically adjusted to match the available bandwidth
 ⇒ real-time ABR or rt-ABR
- q With proper switch algorithm,
 ABR queues in the switches are very small
 ⇒ Negligible delay in the network
- Any switch algorithm with fast transient response and queue control can loosely guarantee low delay through the switch

Scheduling and Buffering Issues

- q Weighted max-min fairness: Allocate rates to flows in proportion to their weights
 - \Rightarrow Higher rate sources are treated preferentially
- q Buffering at the sources and acceptable loss \rightarrow Equivalent bandwidth
 - \Rightarrow Equivalent bandwidth
 - \Rightarrow MCR
 - \Rightarrow Minimum acceptable quality is guaranteed
- q Internet does not provide MCR. ABR does.rt-ABR video will be much better

2. MPEG2 Streams over VBR

- q MPEG2 over ATM Overview
- q Modeling MPEG2 Transport Streams over VBR
- q Simulation Results for Terrestrial Networks
- q Simulation Results for Satellite Networks
- q Ref: "Performance of TCP over ABR with Long-Range Dependent VBR Background Traffic Over Terrestrial and Satellite ATM Networks," ATM Forum Contribution, 97-0177, February 1997, <u>ftp://netlab.ohio-state.edu/pub/jain/atm97-0177.txt</u>

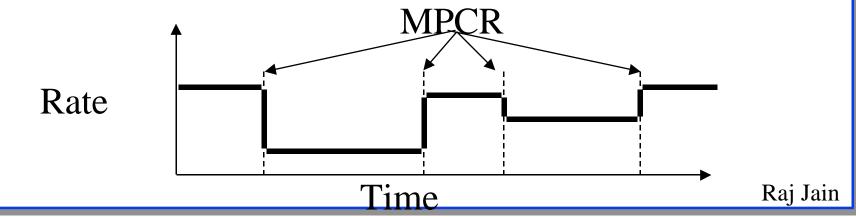


Elementary Stream Ρ В В В В В q Elementary stream: Sequence of I, P, B frames q Individually coded I frames - Large Transmission time = 4 to 5 frame display time q Predictively coded **P** frames - Medium Transmission time = 0.5-1 frame display time q Bidirectionally coded **B** frames - Small

Transmission time = 0.2 frame display time

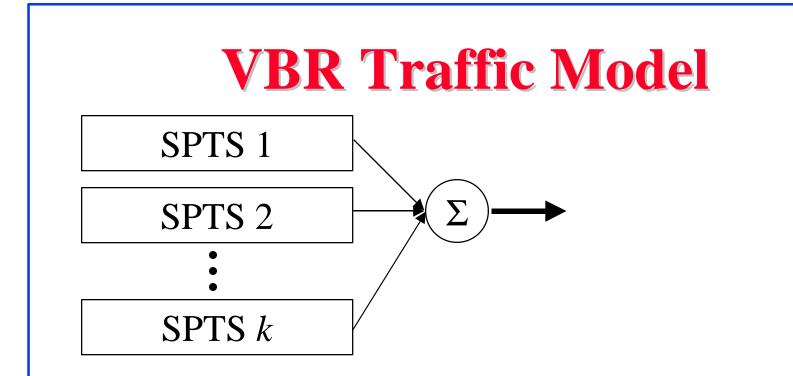
Timestamps in MPEG2

- q Frames may contain a presentation timestamp.
- q To synchronize the clocks, a sample of system clock is sent every 80µs to 100 ms
 MPEG2 Program Clock Reference (MPCR)
 We use MPCR instead of PCR (Peak Cell Rate)
- q MPCRs are used by a phase lock loop
 - \Rightarrow Rate between MPCRs must be constant



MPEG2 Traffic Characteristics

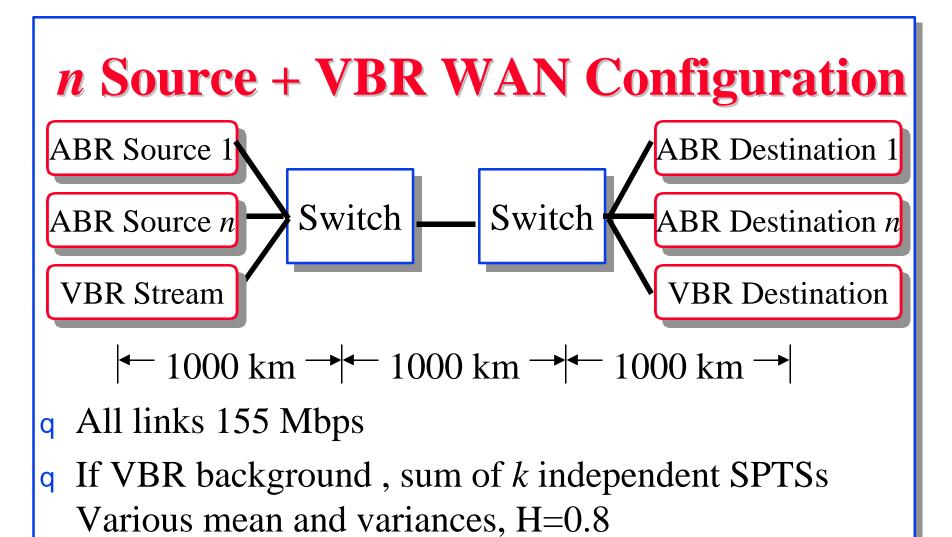
- q Single Program Transport Stream
- q Piecewise CBR
- q Rate changes only at MPCRs
- q Inter-MPCR interval is random
 Standard allows 80µs to 100 ms interval
 Most implementations change only 20 to 100 ms
- q Rate values have a long-range dependence



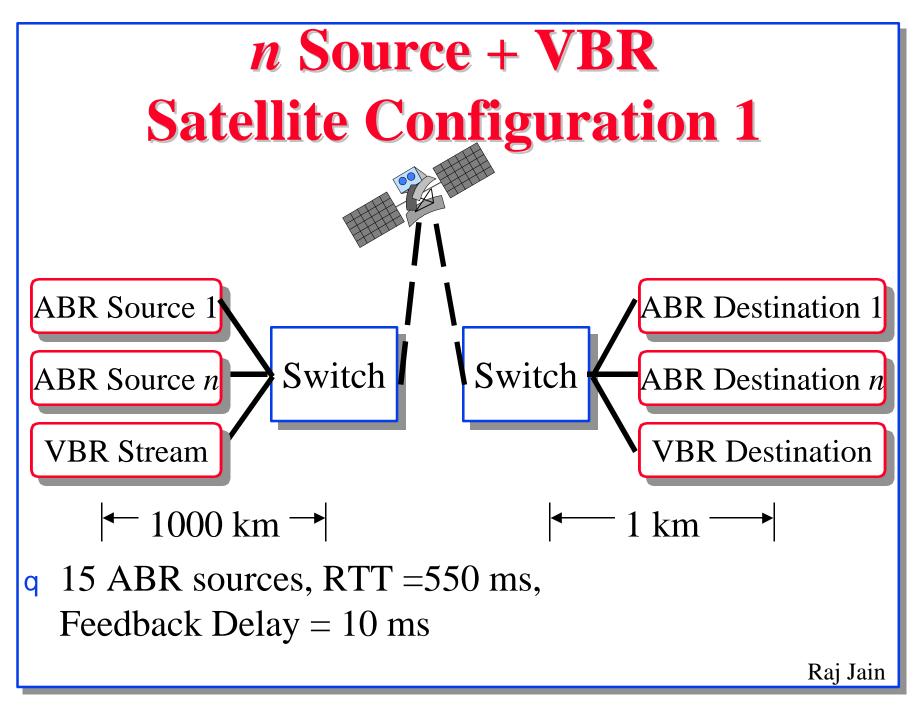
- q VBR background = Sum of k transport streams
- q Each transport stream has
 - q a random inter-MPCR interval = Uniform(20,100)
 - q a random long-range dependent rates (Fractional Gaussian Noise)

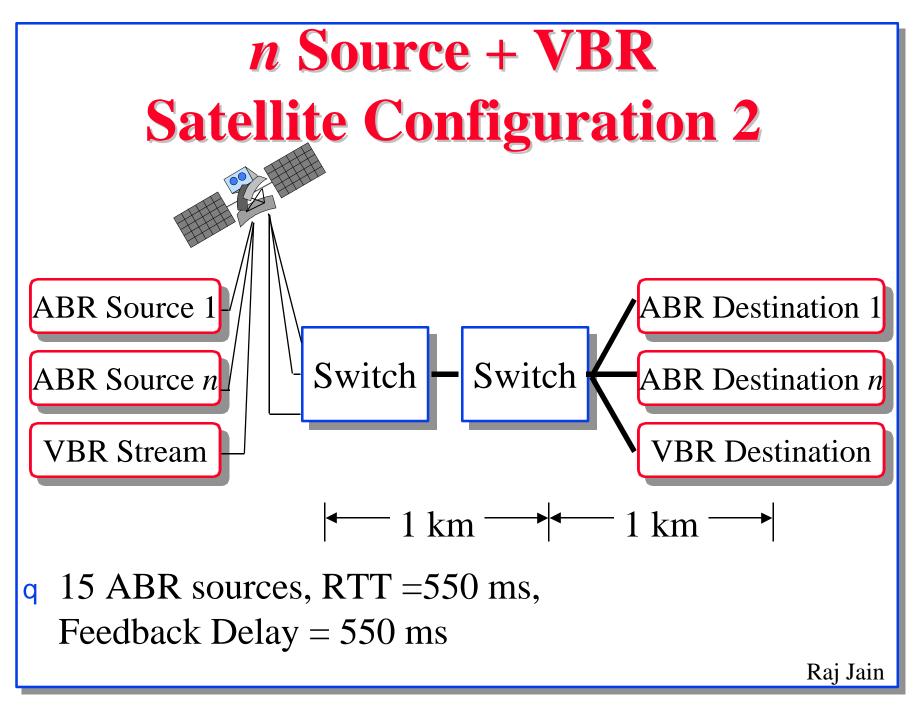
VBR Model (Cont)

q Maximum bandwidth demand = 15 Mbps Minimum bandwidth demand = 0 Mbps \Rightarrow Random numbers below 0 or above 15 are ignored (Pruning) (Alternative choices: clipping or exponentiation were rejected). (a) Gaussian (c) Exponentiation (d) Pruning (b)Clipping $Min{Max{0, x}, 15} Min{2^x, 15}$ Raj Jain



- q All traffic unidirectional; Large file transfer application
- q 15 ABR sources, RTT =30 ms, Feedback Delay = 10 ms Raj Jain



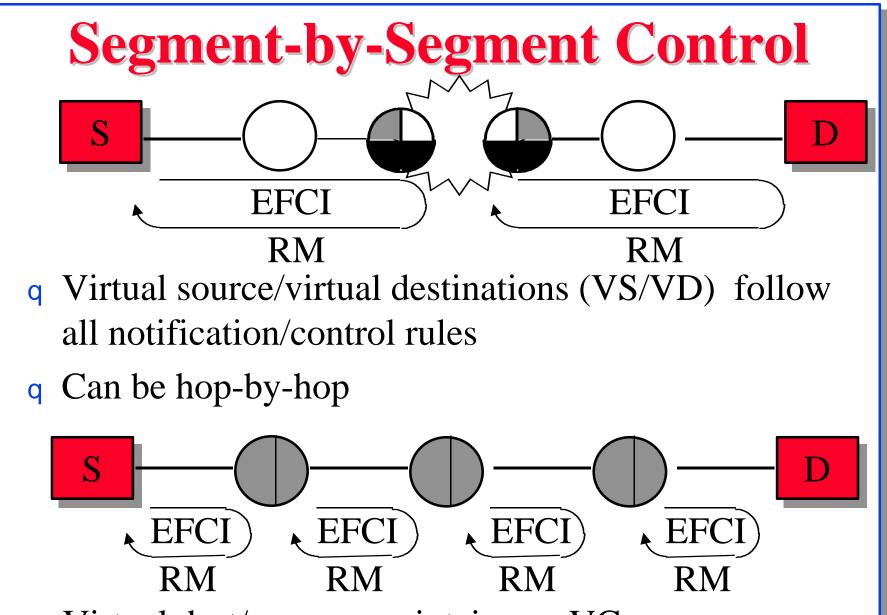


Summary of Results

- q MPEG2 compressed video = piecewise CBR, long-range dependent rate, random inter-MPCR intervals
- ABR with appropriate switch algorithm can handle the randomness in ABR capacity
- q With ERICA+ and Infinite TCP Traffic:
 - q Queue lengths $< 3 \times$ Feedback delay
 - q Efficiency close to the maximum possible.
 - Queues are similar to those with deterministic
 VBR

3. Virtual Source/Virtual Destination

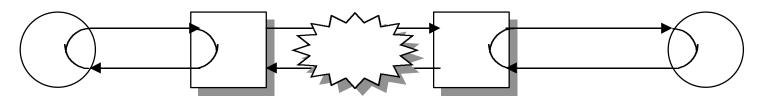
- q Overview of VS/VD
- q Implementation Guidelines
- **q** Simulation results
- q Ref: "Virtual Source/Virtual Destination: Design Considerations," ATM Forum Contribution, 96-1759, December 1996, <u>ftp://netlab.ohio-</u> <u>state.edu/pub/jain/atmf/atm96-1759.ps</u>



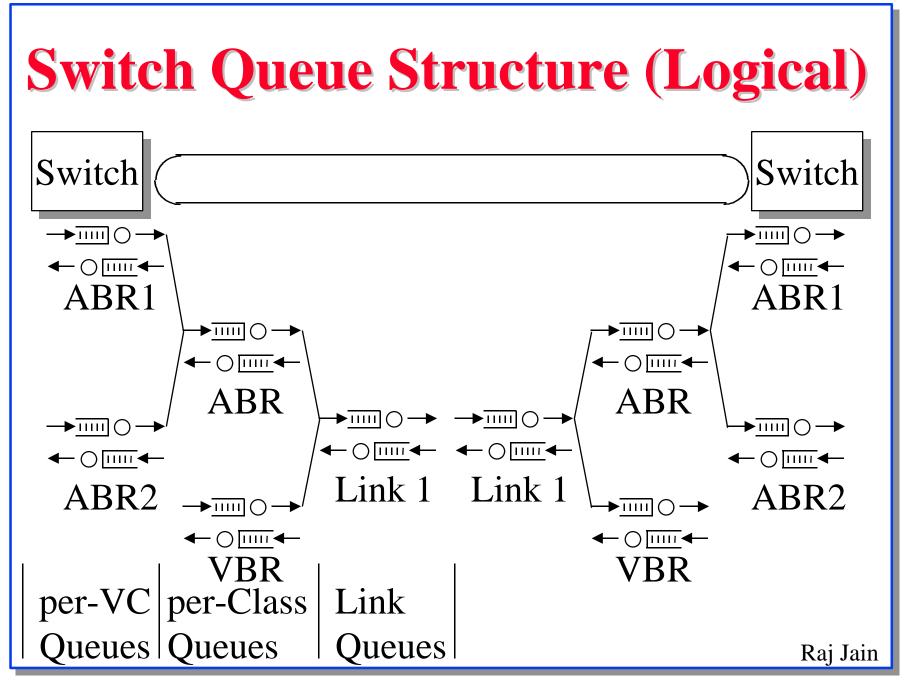
q Virtual dest/sources maintain per-VC queues.

Why Implement VS/VD?

- q Isolates users from the networkOr, isolates different networks
- q Allows proprietary protocol in the intermediate cloud
- q Shorter control loops improve performance



q Little cost to implement VS/VD if per-VC queueing and scheduling is already in the switch. (Queues shared by multiple VCs aren't sufficient.)

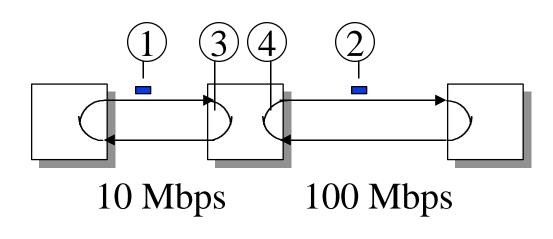


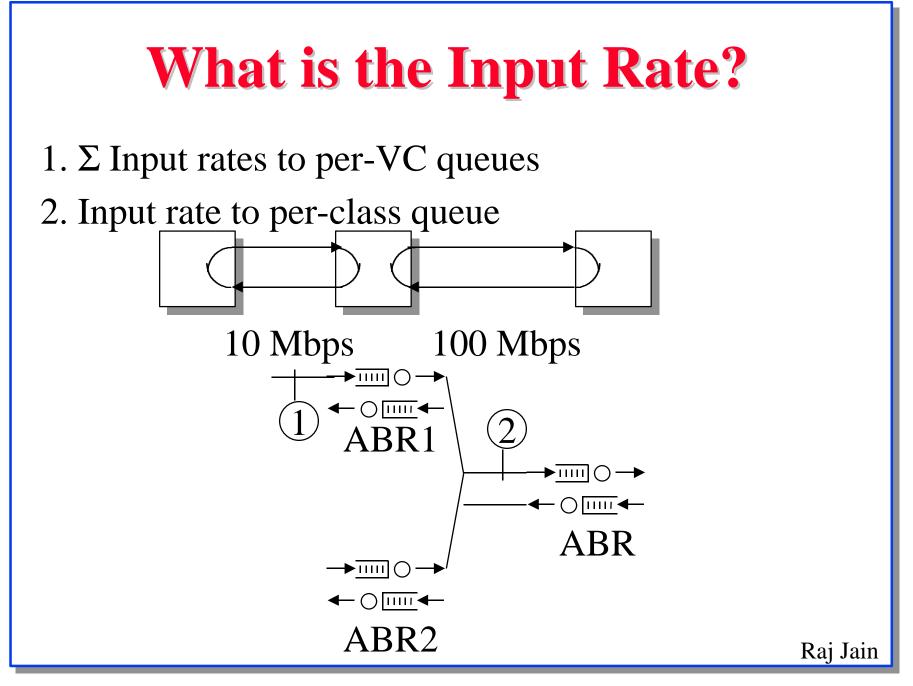
Design Decisions

- **q** What is the VC's rate?
- **q** What is the input rate?
- q Does a link affect current loop or previous loop?
- q When to calculate the VAL?

What is the VC's Rate?

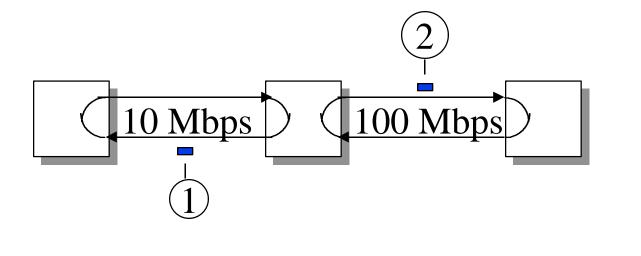
- 1. CCR in FRM1
- 2. CCR in FRM2 = ACR_2
- 3. Measured source rate in the previous loop =VC's input rate to per-VC queue (Not yet analyzed)
- 4. Measured source rate in the next loop =VC's input rate to per-class queue





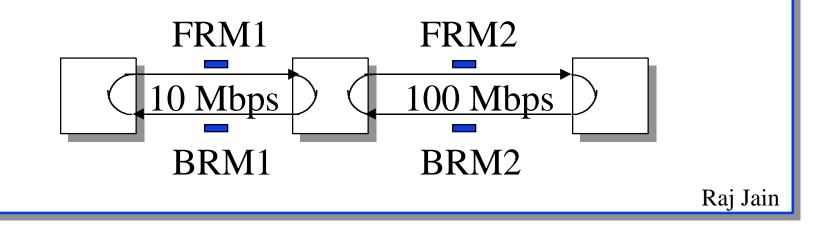
Effect of link congestion

- q Which link affects which loop?E.g., Effect of Link 2 congestion:
 - 1. Change $\text{ER}_1 \Rightarrow$ Previous loop only
 - 2. Change $ACR_2 \Rightarrow$ Next loop only
 - 3. Change ER_1 and $\text{ACR}_2 \Rightarrow \text{Both loops}$



Allocated Rate Update Frequency

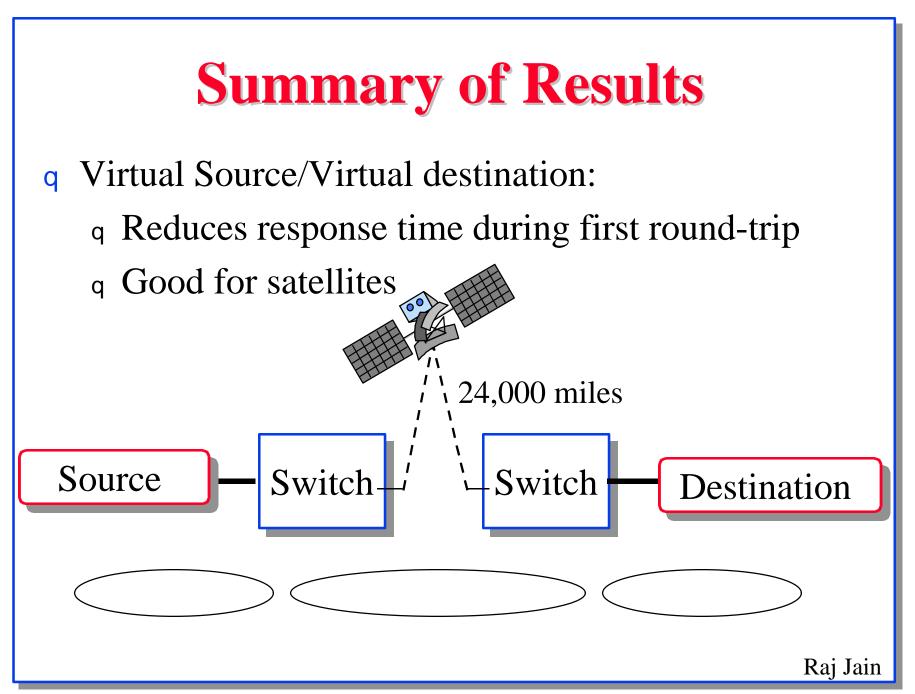
- q When should the rate allocated to a VC be calculated? (Applies only to the previous loop)This is normally done on receiving a BRM in a switch or on turning around an FRM in a destination
 - 1. On receiving BRM2
 - 2. On turning around FRM1
 - 3. Both



Design Decisions: Summary

q Four Decisions:

- 1. What is the VC's rate: 4 alternatives
- 2. What is the input rate: 2 alternatives
- 3. Effect of link congestion: 3 alternatives
- 4. Allocated rate update frequency: 3 alternatives
- q Total $4 \times 2 \times 3 \times 3 = 72$ combinations
- q Some of these combinations do not work
- q Recommendation: Measured VC rate from per-Class Queue, per-class input rate, Control both loops, VC's allocation updated at FRM₁ and at BRM₂



- vS/VD does improve the stability of the network.
 Some cases that diverged with basic ERICA converge with VS/VD.
- vS/VD increases throughput slightly due to reduced response time and reduced convergence time.
- q The effect of VS/VD depends upon the switch algorithm.
- q In VS/VD situations, ACR and actual rates are very different. Cannot rely on CCR field.
 Must measure VC's rate.

Summary

- q Real-time ABR accepted by the industry as a work-item for the next version of ATM Forum Traffic Management
- q MPEG2 Video is piece-wise CBR
- q Developed VS/VD implementation guidelines
 - VS/VD may help in satellite paths.
- q Results are quickly being communicated to industry.

q