High-Speed Networking: Trends and Issues



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- □ Industry trends
- □ High-speed network design
- □ A Simple rule of thumb
- **Trends in traffic**
- □ Trends in network topology

Trend: Telecommunication and Networking

□ From computerization of telephone traffic switching to telephonization of computer traffic switching.



Trend: Networking is Critical

- □ Communication more critical than computing
 ⇒ Bus performance vs ALU speed
 ⇒ I/O performance vs SPECMarks
- □ User Location:
 - 1960: Computer room
 - 1980: Desktop
- □ System Extent:
 - 1980: 1 Node within 10 m
 - 1990: 100 nodes within 10 km

1970: Terminal room 1990: Mobile

- □ Last 10 years: Individual computing Next 10 years: Cooperative computing
- Past: Corporate networks
 Future:
 - Intercorporate networks
 - National Info Infrastructures
 - International Info Infrastructures







Trend: Standardization

- □ Religion must be forgotten
 ⇒ Improve on someone else's ideas as naturally as yours
- □ Can't succeed alone
 - \Rightarrow Innovation + Technology partnerships
- To impact: Participate in standardization
 Publication is too late and insufficient
- □ Vertical vs horizontal specialization \Rightarrow Switch, router, host, applications



Challenge: Economy of Scale

- Technology is far ahead of the applications.
 Invention is becoming the mother of necessity.
 We have high speed fibers, but not enough video traffic.
- □ Low-cost is the primary motivator. Not necessity.
 ⇒ Buyer's market (Like \$99 airline tickets to Bahamas.) Why? vs Why not?
- □ Ten 100-MIPS computer are cheaper than one 1000-MIPS computer \Rightarrow Parallel computing, not supercomputing
- □ Ethernet was and still is cheaper than 10 one-Mbps links.
- No FDDI if it is 10 times as expensive as Ethernet.
 10/100 Ethernet adapters = \$50 over 10 Mbps
- □ Q: Given ATM or 100 Mbps Ethernet at the same cost, which network will you buy?

A: Ethernet. Proven Technology.









- □ Efficiency = Maximum throughput/Media bandwidth
- \square Efficiency is a decreasing function of α
 - = Propagation delay /Transmission time
 - = (Distance/Speed of light)/(Transmission size/Bits/sec)
 - = Distance*Bits/sec/(Speed of light)(Transmission size)
- □ Bit rate-distance-transmission size tradeoff.
- Most people cannot visualize bit rate but can see distance easily.

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Lessons

- For any given access method: the throughput (or efficiency) goes down as either the bit rate is increased, distance is increased, or frame size is decreased.
- □ If you scale the bit rate and packet size by the same factor, all tilizations, delays remain same.
- □ If you increase the bit rate by a factor of 10 but decrease the distance by a factor of 10, ff remains same.
- □ If you increase the bit rate by a factor of 10 but increase the frame size by a factor of 10, ff remains same.
- Designing a high-speed network is somewhat similar to designing a l-speed long-distance network.

Networking to Mars

- □ Distance*speed = constant
- 1 Gb/s between Boston and San Francisco is similar to 56 kb/s to Mars
- □ Earth-Mars Distance/Boston-SF Distance = 49×10^{6} Miles/3128 Miles = 1 Gb/s/56 kb/s
- □ Rule of Thumb: Don't do on a high-speed network, what you wouldn't do on a network to Mars.



What You Wouldn't Do on A Network to Mars?

Media Access:

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- □ Transmit and wait to hear others (e.g., Ethernet)
- Hold token while your frame goes around the ring (e.g., IEEE 802.5)
- □ Hold entire path while using only a part of it (e.g., FDDI) \Rightarrow Spatial reuse



What You Wouldn't Do on A Network to Mars?

Transport or logical-link layer:

- □ Drop all packets if one is lost \Rightarrow (out-of-order caching)
- □ Retransmit all packets when just one is lost \Rightarrow (Selective retransmission)
- □ Wait for a packet to be resent to you if it is lost \Rightarrow (Forward Error Correction)
- □ Wait until last minute to order \Rightarrow (Anticipation, prefetching)
- □ Wait for a three-way (or two-way) handshake before sending first byte \Rightarrow (Implicit handshake)
- □ Summary: Minimize delay vs maximize throughput
 ⇒ Generation gap

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Trends in Applications

- □ Little Voice
- □ AT&T: 125 to 130 M calls/day @ 5 min/call 64 kbps = 28.8 Gbps = 1/1000 of one fiber
- \Box 200 Million X 24 hr/day X 64 kbps = 12.8 Tbps
- □ Survey of 1750 businesses:



Video Characteristics

- \Box Size: 1 Hr uncompressed HDTV= 540 GB = \$150/sec
- \Box 1 Hr compressed HDTV= 9 GB = \$2.5/sec
 - \Rightarrow Needs to be compressed for storage
 - \Rightarrow Variable bit rate
- □ Holding time: At 1 Gbps:
 - -10 Mb image =10 ms
 - 1 hour compressed VHS movie =10 secs or less
 - \Rightarrow Bursty short-lived traffic

Electro-optic Bottleneck

- \Box Bandwidth of fiber = 25 THz/window
- \Box Bandwidth of electronics = 1-10 Gbps
- □ Switching bottleneck ⇒ Optical switching ⇒ All-optical networks
- Switches more expensive than media: Less switches and more links
- □ Higher connectivity, less hops
- Distributed media shared switching (like WANs) and not distributed
- □ switching shared media (like LANs)





- □ Variable bandwidth/station
- \Box Cost \propto bandwidth
- □ Incremental upgradability
- □ Natural spatial reuse

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- □ High-speed links iff economy of scale.
- □ Bursty, short holding time traffic.
- □ Shared-switch distributed-media. No shared-media access.
- □ Speed-distance-transmission size tradeoff \Rightarrow Don't do on a high-speed network what you wouldn't do on a network to Mars.

Further Reading

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