# Classical Encryption Techniques



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Audio/Video recordings of this lecture are available at:

http://www.cse.wustl.edu/~jain/cse571-11/

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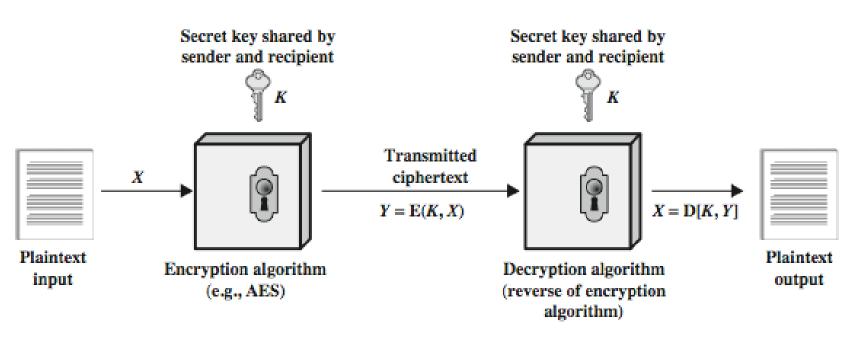


- 1. Symmetric Cipher Model
- 2. Substitution Techniques
- 3. Transposition Techniques
- 4. Product Ciphers
- 5. Steganography

These slides are based on Lawrie Brown's slides supplied with William Stalling's book "Cryptography and Network Security: Principles and Practice," 5<sup>th</sup> Ed, 2011.

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# **Symmetric Cipher Model**



$Y = \mathrm{E}(\mathrm{K}, X)$	)
X = D(K,	Y)

K=Secret Key Same key is used for encryption and decryption.  $\Rightarrow$  Single-key or private key encryption.

# **Some Basic Terminology**

- **Plaintext** original message
- **Ciphertext** coded message
- **Cipher** algorithm for transforming plaintext to ciphertext
- **Key** info used in cipher known only to sender/receiver
- **Encipher (encrypt)** converting plaintext to ciphertext
- **Decipher (decrypt)** recovering ciphertext from plaintext
- **Cryptography** study of encryption principles/methods
- Cryptanalysis (code breaking) study of principles/ methods of deciphering ciphertext without knowing key
- **Cryptology** field of both cryptography and cryptanalysis

# **Cryptography Classification**

### □ By type of encryption operations used

- Substitution
- > Transposition
- > Product
- By number of keys used
  - > Single-key or private
  - > Two-key or public
- □ By the way in which plaintext is processed
  - > Block

### > Stream

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# Cryptanalysis

Objective: To recover key not just message

### □ Approaches:

- Cryptanalytic attack
- > Brute-force attack
- □ If either succeed all key use is compromised

### □ Brute-force attack:

Key Size (bits)	Number of Alternative Keys	Time requi	red at 1 decryption/µs	Time required at 10 <sup>6</sup> decryptions/μs
32	$2^{32} = 4.3 \times 10^9$	2 <sup>31</sup> μs	= 35.8 minutes	2.15 milliseconds
56	$2^{56} = 7.2 \times 10^{16}$	2 <sup>55</sup> μs	= 1142 years	10.01 hours
128	$2^{128} = 3.4 \times 10^{38}$	2 <sup>127</sup> μs	$= 5.4 \times 10^{24}$ years	$5.4 \times 10^{18}$ years
168	$2^{168} = 3.7 \times 10^{50}$	2 <sup>167</sup> µs	$= 5.9 \times 10^{36}$ years	$5.9 \times 10^{30}$ years
26 characters (permutation)	$26! = 4 \times 10^{26}$	$2 \times 10^{26} \ \mu s$	$= 6.4 \times 10^{12}$ years	$6.4 \times 10^6$ years
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### **Substitution**

□ Caesar Cipher: Replaces each letter by 3rd letter on

### □ Example:

meet me after the toga party PHHW PH DIWHU WKH WRJD SDUWB

### □ Can define transformation as:

a b c d e f g h i j k l m n o p q r s t u v w x y z D E F G H I J K L M N O P Q R S T U V W X Y Z A B C

#### ■ Mathematically give each letter a number abcdefghijklmnopqrstuvwxyz 012345678910111213141516171819202122232425

□ Then have Caesar cipher as:

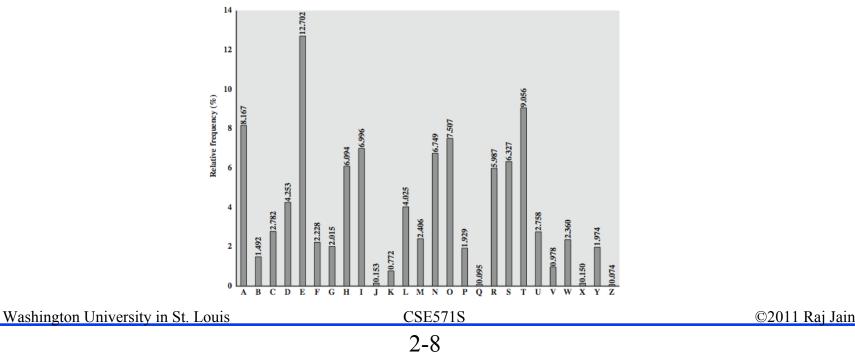
$$c = E(k, p) = (p + k) \mod (26)$$
  
 $p = D(k, c) = (c - k) \mod (26)$ 

□ Weakness: Total 26 keys

# **Substitution: Other forms**

Random substitution:
 Plain: abcdefghijklmnopqrstuvwxyz
 Cipher: DKVQFIBJWPESCXHTMYAUOLRGZN
 The key is 26 character long
 => 26! (= 4 x 10<sup>26</sup>) Keys in place of 26 keys

Letter frequencies to find common letters: E,T,R,N,I,O,A,S



# **Substitution: Other forms (Cont)**

Use two-letter combinations: Playfair Cipher
 Use multiple letter combinations: Hill Cipher

# **Poly-alphabetic Substitution Ciphers**

- Use multiple ciphers. Use a key to select which alphabet (code) is used for each letter of the message
- Vigenère Cipher: Example using keyword deceptive key: deceptivedeceptivedeceptive plaintext: wearediscoveredsaveyourself ciphertext:ZICVTWQNGRZGVTWAVZHCQYGLMGJ

### **One-Time Pad**

- If a truly random key as long as the message is used, the cipher will be secure
- □ Called a One-Time pad
- Is unbreakable since ciphertext bears no statistical relationship to the plaintext
- Since for any plaintext & any ciphertext there exists a key mapping one to other
- □ Can only use the key **once** though
- Problems in generation & safe distribution of key

### **Transposition (Permutation) Ciphers**

Rearrange the letter order without altering the actual letters
 Rail Fence Cipher: Write message out diagonally as:

```
mematrhtgpry
```

etefeteoaat

- Giving ciphertext: MEMATRHTGPRYETEFETEOAAT
- Row Transposition Ciphers: Write letters in rows, reorder the columns according to the key before reading off.
  Key: 4312567
  Column Out 4 3 1 2 5 6 7
  Plaintext: a t t a c k p
  o s t p o n e
  d u n t i l t
  w o a m x y z
  Ciphertext: TTNAAPTMTSUOAODWCOIXKNLYPETZ
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# **Product Ciphers**

Use several ciphers in succession to make harder, but:

- > Two substitutions make a more complex substitution
- > Two transpositions make more complex transposition
- > But a substitution followed by a transposition makes a new much harder cipher
- □ This is a bridge from classical to modern ciphers

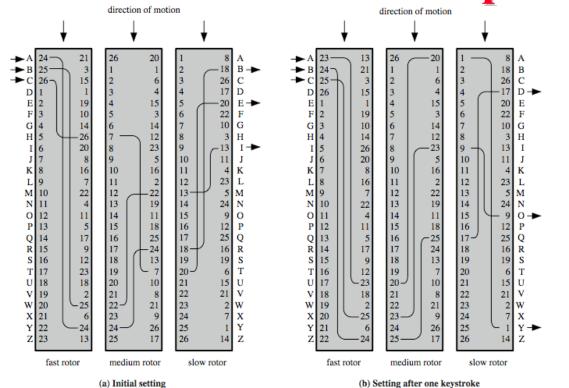
### **Rotor Machines**

- Before modern ciphers, rotor machines were most common complex ciphers in use
- □ Widely used in WW2
  - German Enigma, Allied Hagelin, Japanese Purple
- Implemented a very complex, varying substitution cipher
- Used a series of cylinders, each giving one substitution, which rotated and changed after each letter was encrypted



Hagelin Rotor Machine

### **Rotor Machine Principle**



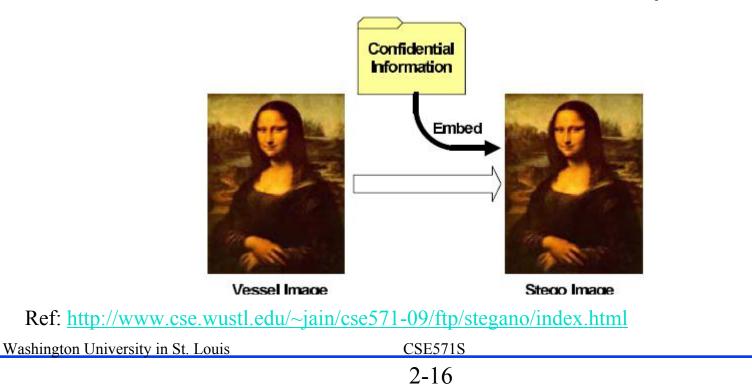
- A becomes Y (First rotor). Y becomes R (2<sup>nd</sup> rotor). R becomes B (3<sup>rd</sup> rotor).
- □ After each letter, first rotor moves 1 position. After each full rotation of 1<sup>st</sup> rotor, 2<sup>nd</sup> rotor moves by 1 position.

• Cycle length = 
$$26^3$$

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## Steganography

- □ Hide characters in a text, hide bits in a photograph
- Least significant bit (lsb) of a digital photograph may be a message.
- Drawback: high overhead to hide relatively few info bits
- Advantage: Can obscure encryption use



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- 1. The key methods for cryptography are: Substitution and transposition
- 2. Letter frequency can be used to break substitution
- 3. Substitution can be extended to multiple letters and multiple ciphers. Mono-alphabetic=1 cipher, Poly-alphabetic=multiple ciphers
- 4. Examples: Caesar cipher (1 letter substitution), Playfair (2letter), Hill (multiple letters), Vigenere (poly-alphabetic).
- 5. Multiple stages of substitution and transposition can be used to form strong ciphers.

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### **Homework 2**

- □ Submit solution to problem 2.18
- 2.18 This problem explores the use of a one-time pad version of the Vigenere cipher. In this scheme, the key is a stream of random numbers between 0 and 26. For example, if the key is 3 19 5..., then the first letter of the plaintext is encrypted with a shift of 3 letters, the second with a shift of 19 letters, the third with a shift of 5 letters, and so on.
- A. Encrypt the plain text sendmoremoney with the key stream 9 0 1 7 23 15 21 14 11 11 2 8 9
- B. Using the ciphertext produced in part (a), find a key so that the cipher text decrypts to the plain text cashnotneeded.