

Part II

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TACACS

- Terminal Access Controller Access-Control System
- □ Routing nodes in ARPAnet were called IMPS.
- □ IMPs with dial up access were called TIPs.
- □ BBN developed TACACS for ARPANET
- AAA server is a process in a UNIX server called TACACS daemon.
- □ Uses UDP port 49
- ❑ Username and passwords were sent in clear for authentication
 ⇒ No longer used
- Cisco adopted TACACS for terminal servers extended TACACS or XTACACS

TACACS+

- Terminal Access Controller Access-Control System Plus
- Cisco's further improved version of TACACS and XTACACS
- □ Not compatible with TACACS
- □ Payload is encrypted
- Described in draft-grant-tacacs-02.txt, Jan 1997.
- Uses TCP port 49

RADIUS

- **RFC 2138, June 2000**
- **UDP** port 1812
- □ Why UDP?
 - ➤ In case of server failure, the request must be re-sent to backup ⇒ Application level retransmission required
 - > TCP takes to long to indicate failure
 - > Stateless protocol

RADIUS Packet Format

Code	Identifier	Length	Authenticator	Attributes
1 B	1 B	2B	16B	

Codes:

- 1 = Access Request
- 2 = Access Accept
- 3 = Access Reject
- 4 = Accounting request
- 5 = Accounting Response
- 11 = Access Challenge
- 12 = Server Status (experimental)
- 13 = Client Status (Experimental)

255 = Reserved

RADIUS Packet Format (Cont)

- I6B Authenticator is used to authenticate the reply from the RADIUS server
- In Access-request packets 16B random number is send as authenticator
- □ Password in packet
 - = MD5(Shared secret | authenticator) \oplus password
- **Response** Authenticator

= MD5(Code|ID|Length|Request Auth|Attributes| Shared secret)

□ All attributes are TLV encoded.

RADIUS Accounting

- **RFC 2866, June 2000**
- □ Client sends to the server:
 - > Accounting Start Packet at service beginning
 - > Accounting Stop Packet at end
- □ All packets are acked by the server
- Packet format same as in authentication

RADIUS Server Implementations

- **Public domain software implementations:**
- □ FreeRADIUS
- GNU RADIUS
- JRadius
- OpenRADIUS
- Cistron RADIUS
- **BSDRadius**
- TekRADIUS

Problems with RADIUS

- ❑ Does not define standard failover mechanism ⇒ varying implementations
- Original RADIUS defines integrity only for response packets
- RADIUS extensions define integrity for EAP sessions
- Does not support per-packet confidentiality
- Billing replay protection is assumed in server. Not provided by protocol.
- □ IPsec is optional
- □ Runs on UDP ⇒ Reliability varies between implementation. Billing packet loss may result in revenue loss.
- RADIUS does not define expected behavior for proxies, redirects, and relays ⇒ No standard for proxy chaining

Problems with RADIUS (Cont)

- Does not allow server initiated messages
 No On-demand authentication and unsolicited disconnect
- Does not define data object security mechanism
 ⇒ Untrusted proxies can modify attributes
- Does not support error messages
- Does not support capability negotiation
- □ No mandatory/non-mandatory flag for attributes
- Servers name/address should be manually configured in clients ⇒ Administrative burden ⇒ Temptation to reuse shared secrets

Diameter Base Protocol

- **RFC 3588, Sep 2003**
- Defines standard failover algorithm
- Runs over TCP and Stream Control Transmission Protocol (SCTP)
- **D** PDU format incompatible with RADIUS
- □ Can co-exist with RADIUS in the same network
- **Supports:**
 - > Delivery of attribute-value pairs (AVPs)
 - Capability negotiation
 - Error notification
 - > Ability to add new commands and AVPs
 - > Discovery of servers via DNS
- Dynamic session key derivation via TLS Washington University in St. Louis

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Diameter Base Protocol (Cont)

- □ All data is delivered in the form of AVPs
- □ AVPs have mandatory/non-mandatory bit
- \Box Peer-to-peer protocol \Rightarrow any node can initiate request.
- Documents: Base, transport profile, applications
- Applications: NAS, Mobile IP, Credit control (prepaid, post-paid, credit-debit), 3G, EAP, SIP

AAA Transport Profile

- **RFC 3539, June 2003**
- □ Network Access Identifier (NAI) = User ID
- ❑ Application driven vs. network driven: Network is not the bottleneck for AAA messages ⇒ Application driven. No congestion issues.
- □ Slow Failover: TCP time outs \Rightarrow slow
- Use of Nagle Algorithm: Many AAA messages are combined in one TCP message
- Multiple Connections: Max 256 requests in progress between a client and a server
- Duplicate Detection: Servers and clients recognize duplicate request or responses and discard them.

A single request when duplicated can result in success and failure responses.
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AAA Transport Profile (Cont)

- Invalidation of Transport Parameter Estimates: Timeouts should account for network congestion
- Inability to use fast re-transmit: most AAA protocols are always close to initial window set to 1 or 2
- **Congestion Avoidance:**
- \Box Delayed Acks: application driven \Rightarrow explicit acks
- Premature failover: some implementation switch to backup server prematurely
- ❑ Head of line blocking: TCP queue may build up after a packet loss ⇒ hold up other AAA requests on the same connection
- □ Connection load balancing:

AAA Key Management Principles

- □ RFC 4962, July 2007 (Housley Criteria)
- ❑ Ability to negotiate crypto algorithms
 ⇒ Support multiple algorithm
- □ Ability to negotiate key derivation function is not required
- □ At least one suite of mandatory algorithms must be selected
- □ Use strong fresh session keys.
- Session keys must not be dependent on one another
 - \Rightarrow Knowing a session key, Can't find another session key
 - \Rightarrow Use nonce to ensure each session key is fresh.
- □ Include replay detection mechanism
- □ Authenticate all parties
- Lower layer identifiers used for authorization should be authenticated

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AAA Key Management Principles (Cont)

- Both peer and authenticator must be authorized
 Detect unauthorized authenticator
- Peer, Authenticator, Authentication server should have a common view of authorizations
- □ Cipher suite selection should be securely confirmed ⇒ detect roll-back attacks
- All keys should be uniquely named and key name should disclose key value
- □ Prevent domino effect ⇒ Compromise of a single entity must not compromise key material at other entities in other branches (may compromise children entities)
- Bind key to its context: use, who has access, life time. All entities with access to keying material should have the same context.

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- □ TACACS and TACACS+ are legacy AAA protocols
- RADIUS provides good security but lacks sophisticated mechanisms required for failover
- Diameter is a replacement for RADIUS. Fixes most known shortcomings of RADIUS.