

Networking Layer Protocols for Internet of Things: 6LoWPAN and RPL



Raj Jain
Washington University in Saint Louis
Saint Louis, MO 63130
Jain@cse.wustl.edu

These slides and audio/video recordings of this class lecture are at:

<http://www.cse.wustl.edu/~jain/cse570-23/>

Student Questions



- ❑ 6LoWPAN
 - Adaptation Layer
 - Address Formation
 - Compression
- ❑ RPL
 - RPL Concepts
 - RPL Control Messages
 - RPL Data Forwarding

Note: This is part 3 of a series of class lectures on IoT.

Student Questions

Recent Protocols for IoT

| | | | |
|----------|---|---|---|
| Session | MQTT, SMQTT, CoRE, DDS, AMQP, XMPP, CoAP, IEC, IEEE 1888, ... | Security | Management |
| Network | Encapsulation 6LoWPAN, 6TiSCH, 6Lo, Thread... Routing RPL, CORPL, CARP | IEEE 1888.3, TCG, Oath 2.0, SMACK, SASL, EDSA, ace, DTLS, Dice, ... | IEEE 1905, IEEE 1451, IEEE 1377, IEEE P1828, IEEE P1856 |
| Datalink | Wi-Fi, Bluetooth Low Energy, Z-Wave, ZigBee Smart, DECT/ULE, 3G/LTE, NFC, Weightless, HomePlug GP, 802.11ah, 802.15.4e, G.9959, WirelessHART, DASH7, ANT+, LTE-A, LoRaWAN, ISA100.11a, DigiMesh, WiMAX, ... | | |

Student Questions

Ref: Tara Salman, Raj Jain, "A Survey of Protocols and Standards for Internet of Things," Advanced Computing and Communications, Vol. 1, No. 1, March 2017, http://www.cse.wustl.edu/~jain/papers/iot_accs.htm
 Washington University in St. Louis <http://www.cse.wustl.edu/~jain/cse570-23/> ©2023 Raj Jain

IEEE 802.15.4

- ❑ Wireless Personal Area Network (WPAN)
- ❑ Allows mesh networking.
Full-function nodes can forward packets to other nodes.
- ❑ A PAN coordinator (like Wi-Fi Access Point) allows nodes to join the network.
- ❑ Nodes have 64-bit addresses
- ❑ The coordinator assigns a 16-bit short address for use during the association
- ❑ Maximum frame size is 127 bytes
- ❑ More details in CSE 574 wireless networking course
http://www.cse.wustl.edu/~jain/cse574-20/j_12wpan.htm

Student Questions

- ❑ IEEE 802.15.4 is about personal area networks. Isn't a 64-bit address too large? Is it possible to have such many nodes in a personal area network?

The number of unique PAN objects in the world is significant. So, we need a sizeable global address. For any one network, the number is small, so we use a 16-bit short address.

EUI64 Addresses

- ❑ **Ethernet addresses:** 48-bit MAC

| | | | |
|----------------------|--------------------|-------------------------------------|--------------------------|
| Unicast Multicast | Universal Local | Organizationally Unique ID (OUI) | Manufacturer Assigned |
|----------------------|--------------------|-------------------------------------|--------------------------|

1b 1b 22b 24b

- ❑ **IEEE 802.15.4 Addresses:** 64-bit Extended Unique Id (EUI)

| | | | |
|----------------------|--------------------|-------------------------------------|--------------------------|
| Unicast Multicast | Universal Local | Organizationally Unique ID (OUI) | Manufacturer Assigned |
|----------------------|--------------------|-------------------------------------|--------------------------|

1b 1b 22b 40b

- ❑ **Local bit** was incorrectly assigned. $L=1 \Rightarrow$ Local, but all-broadcast address = all 1's is not local
IETF RFC4291 changed the meaning so that $L=0 \Rightarrow$ Local
The 2nd bit is now called the Universal bit (U-bit)
 \Rightarrow U-bit formatted EUI64 addresses.

| | | | |
|----------------------|--------------------|-------------------------------------|--------------------------|
| Unicast Multicast | Local Universal | Organizationally Unique ID (OUI) | Manufacturer Assigned |
|----------------------|--------------------|-------------------------------------|--------------------------|

1b 1b 22b 40b

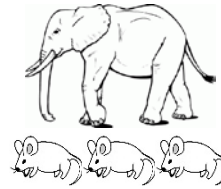
Student Questions

- ❑ So, to sum up, in EUI 64, local = 0; in Ethernet 48-bit MAC, local = 1, except for the all-broadcast address. Is it correct?

Yes.

6LowPAN

- ❑ IPv6 over Low Power Wireless Personal Area Networks
- ❑ How to transmit IPv6 datagrams (elephants) over low-power IoT devices (mice)?



❑ Issues:

1. **IPv6 address formation:** 128-bit IPv6 from 64-bit EUI64
2. **Maximum Transmission Unit (MTU):** IPv6 at least 1280 bytes vs. IEEE 802.15.4 standard packet size is 127 bytes

| | | |
|-----------------|-----------------|---------|
| 802.15.4 Header | Security Option | Payload |
| 25B | 21B | 81B |

3. **Address Resolution:** 128b or 16B IPv6 addresses. 802.15.4 devices use 64-bit (no network prefix) or 16-bit addresses
4. **Optional mesh routing in the datalink layer**
⇒ Need destination and intermediate addresses.

Student Questions

Ref: G. Montenegro, et al., "Transmission of IPv6 Packets over IEEE 802.15.4 Networks," RFC 4944, Sep 2007, <http://tools.ietf.org/pdf/rfc4944>

Washington University in St. Louis

<http://www.cse.wustl.edu/~jain/cse570-23/>

©2023 Raj Jain

6LowPAN Adaptation Layer

5. MAC-level retransmissions versus end-to-end:

- Optional hop-by-hop ack feature of 802.15.4 is used, but the max number of retransmissions is kept low (to avoid overlapping L2 and L4 retransmissions)

6. Extension Headers: 8b *or less* Shannon-coded dispatch

⇒ header type

- 10_2 : Mesh addressing header (2-bit dispatch)
- $11x00_2$: Destination Processing Fragment header (5-bit)
- 01010000_2 : Hop-by-hop LowPAN Broadcast header (8-bit)

7. IPv6 and UDP header compression

| | | | | | | | | | | |
|---------------|--------|-------|------------|------------------|---------|-----------|---------|-----------|---------|--------------|
| Frame Control | Seq. # | Adrs | [Security] | Disp bits | Ext Hdr | Disp bits | Ext Hdr | Disp bits | Ext Hdr | IPv6 Payload |
| 2B | 1B | 0-20B | 0-21B | | | | | | | |

Student Questions

- ❑ Which extension headers are always present in the 6LowPan packet, and which are optional?

All of the extension headers in IPv6 are optional. However, some are necessary for certain situations. 6LowPAN compresses whatever is present in the IPv6 header.

- ❑ How does the security part work? Does it contain something like a private key?

Any secret or private information cannot be included in the headers that are not encrypted. It is open information that the receiver can use but has no meaning for others without complete context, e.g., an initialization vector (IV).

IPv6 Address Formation

- ❑ **Link-Local IPv6 address** = FE80::U-bit formatted EUI64
- ❑ **Example:**
 - EUI64 Local Address = 40::1 = 0100 0000::0000 0001
 - U-bit formatted EUI64 = 0::1
 - IPv6 Link-local address = FE80::1 = 1111 1110 1000 0000::1
- ❑ IEEE 802.15.4 allows nodes to have 16-bit **short addresses**, and each PAN has a 16-bit **PAN ID**.
1st bit of Short address and PAN ID is Unicast/Multicast
The 2nd bit of Short Address and PAN ID is Local/Universal.
You can broadcast to all members of a PAN or all PANs.
- ❑ IPv6 Link Local Address = FE80 :: PAN ID : Short Address
Use 0 if the PAN ID is unknown.
2nd bit of PAN ID should always be zero since it is always local (assigned by the local admin).

Student Questions

- ❑ Based on RFC 4291, the EUI64 is generated by adding two octets (FF: FE) in the middle of a MAC address, and the seventh bit in EUI64 is the bit that indicates if it's local or universal. Is it the case in this example? If this is the case, should I assume that when solving HW 12A?
RFC4291 explains how to get an EUI address from a 48-bit MAC address. In HW 12A, EUI is given, and you do not need to construct it. Use it exactly as given. IEEE introduced EUI by making it a more significant bit similar to MAC addresses. IETF uses U-Bit formatted EUI. IEEE writes bits as seen on the wire. IETF writes bits as stored inside "their" computer. We use IEEE bit order, but when talking about IETF protocols, we use U-bit encoded EUI.
- ❑ Could you explain the example one more time? *Sure.*
- ❑ Is it possible to abbreviate '0::1' as '::1'? *Yes.*
- ❑ How many bits are there in EUI64 Address, U-bit formatted EUI64, IPv6 Link-local address? Is that 64, 64, 128? *Yes.*

IPv6 Address Formation

- ❑ **Link-Local IPv6 address** = FE80::U-bit formatted EUI64
- ❑ **Example:**
 - EUI64 Local Address = 40::1 = 0100 0000::0000 0001
 - U-bit formatted EUI64 = 0::1
 - IPv6 Link-local address = FE80::1 = 1111 1110 1000 0000::1
- ❑ IEEE 802.15.4 allows nodes to have 16-bit **short addresses**, and each PAN has a 16-bit **PAN ID**.
1st bit of Short address and PAN ID is Unicast/Multicast
The 2nd bit of Short Address and PAN ID is Local/Universal.
You can broadcast to all members of a PAN or all PANs.
- ❑ IPv6 Link Local Address = FE80 :: PAN ID : Short Address
Use 0 if the PAN ID is unknown.
2nd bit of PAN ID should always be zero since it is always local (assigned by the local admin).

Student Questions

- ❑ How to change EUI64 Address into U-bit formatted EUI64? Does that set the second bit to 0, or are there other rules?
Complement the 2nd bit.

Homework 12A

- ❑ [8 points] What is the IPv6 Link-Local address for a IEEE 802.15.4 node whose EUI64 address in hex is 0000::0002. Indicate your final answer in hex without using ::
- ❑ EUI64 in Binary =
- ❑ U-bit EUI64 Binary =
- ❑ U-bit EUI64 Hex =
- ❑ IPv6 Link Local Address =

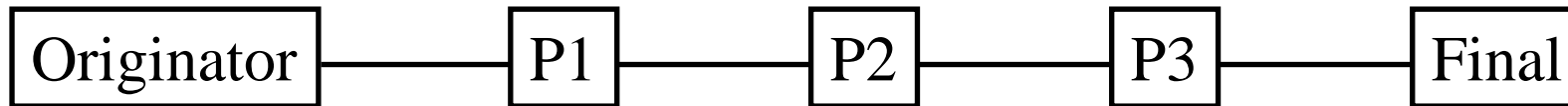
Student Questions

- ❑ The answer for IPV6 Local address should be in Hex, correct?

Yes.

Mesh Addressing Header

- ❑ Dispatch = 10_2 (2 bits) \Rightarrow Mesh Addressing Header
- ❑ MAC header contains per-hop source and destination
- ❑ Original source and destination addresses are saved in the Mesh addressing header
- ❑ A 4-bit hops-left field is decremented at each hop



| | | | | | |
|----------|----|----|-----------|--------------------|---------------|
| Dispatch | V | F | Hops Left | Originator Address | Final Address |
| 10 | | | | | |
| 2b | 1b | 1b | 4b | 16b/64b | 16b/64bit |

V=0 \Rightarrow Originator address is EUI64, V=1 \Rightarrow 16bit

F=0 \Rightarrow Final address is EUI64, F=1 \Rightarrow 16-bit

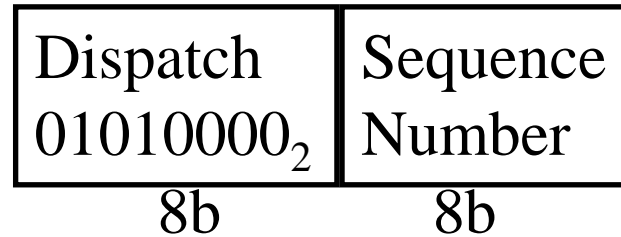
Student Questions

- ❑ Why 10 indicates the end of the mesh header? What other cases might we have?
Shannon coding: Most frequent items should be coded with the least bits. Other dispatch items will be 11... or 01...
- ❑ Does the “number of hops” limit the IOT device’s reach?

The hop limit comes from the IP header.

6LowPAN Broadcast Header

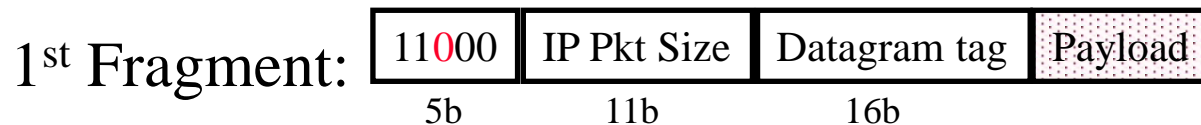
- ❑ For Mesh broadcast/multicast
- ❑ A new sequence number is put in every broadcast message by the originator



Student Questions

6LowPAN Fragment Header

- ❑ Dispatch = 11x00 (5 bits) \Rightarrow Fragment Header
- ❑ Full packet size in the first fragment's fragment header
- ❑ Datagram tag = sequence number
 \Rightarrow Fragments of the same packet
- ❑ Fragment Offset in multiples of 8 bytes

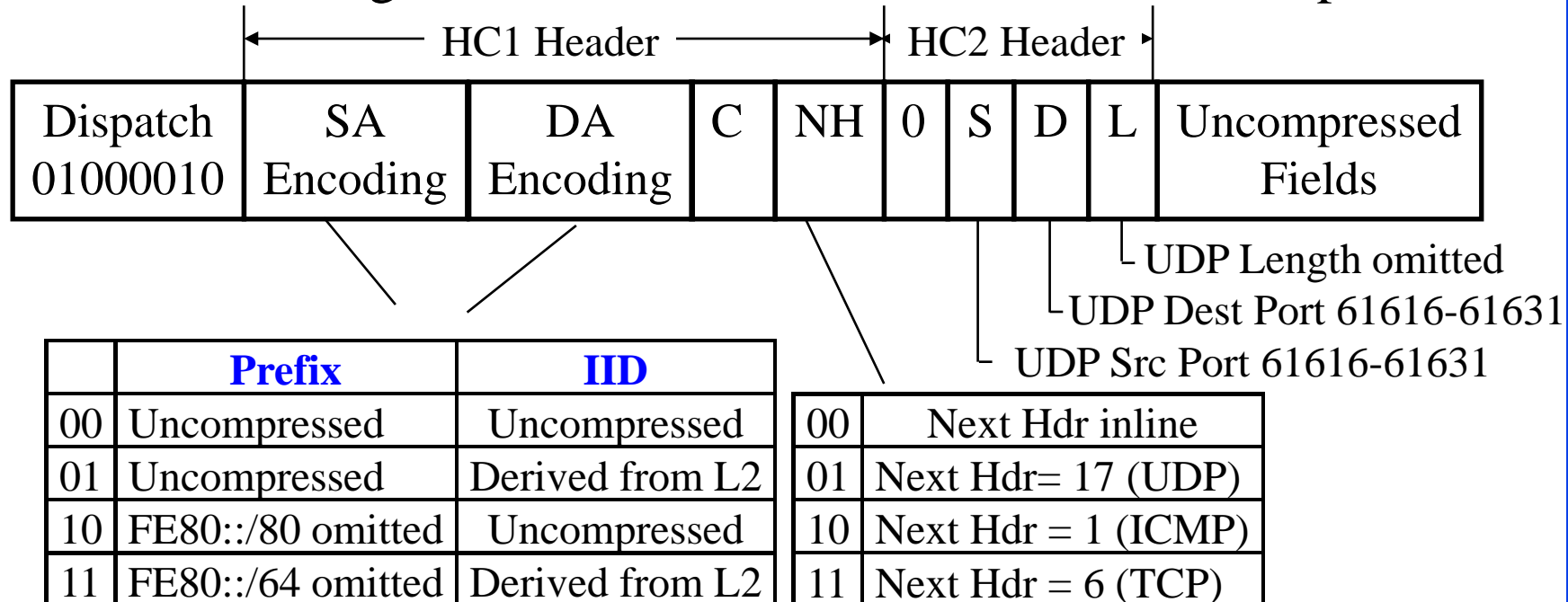


Student Questions

- ❑ A fragment offset of 8 bytes indicates that no fragment size can be larger than 8 bytes.
*No. Offset = Distance from the beginning of the datagram. $X*8$ is the maximum offset. X is encoded in 8 bits, so the maximum offset is $(2^8-1)*8$ bytes.*
- ❑ A Datagram offset of 8 bits means we can't have more than 2^8 fragments per packet.
See above.
- ❑ Does fragment offset correspond to datagram offset in the image? *Yes.*

IP+UDP Header Compression: Stateless

- ❑ Called **HC1-HC2 compression** (not recommended)
- ❑ The IP version field is omitted
- ❑ Flow label field if zero is omitted and C=1
- ❑ Only 4b UDP ports are sent if between 61616-61631 (F0Bx)
- ❑ The UDP length field is omitted. IP addresses are compressed.



Student Questions

- ❑ If the flow label is not omitted, would it be within the HC1 or HC2 header?

It is omitted if it is zero. Otherwise, it will be in "Uncompressed fields."

- ❑ The x in "F0Bx" is 4 bits? If so, how can we represent the 61631-61616= 16 ports with 4 bits?

61631-61616= 15₁₀=1111₂

- ❑ Does the dispatch number 0100 0010 mean anything here?

This indicates the type of 6LowPAN header.

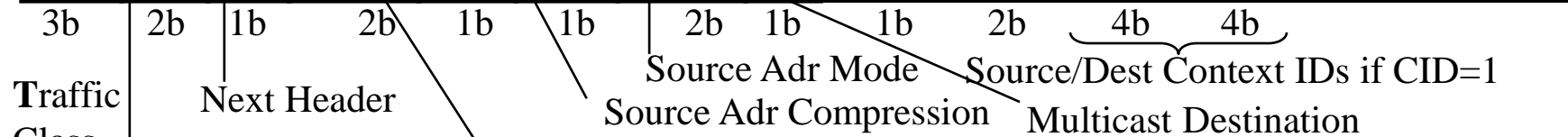
- ❑ The 80 and 64 bits are omitted from the LSB, correct?
- ❑ In the second table, if the next header=NH=17, then what is 01?

17, 1, and 6 are IP payload types. NH=01 indicates the payload is UDP. When uncompressed, it will be encoded as 17.

Context Based Compression

- ❑ HC1 works only with **link-local** addresses
- ❑ Need globally routable IPv6 addresses for outside nodes
- ❑ IPHC uses a 3b dispatch code and a 13-bit base header

| | | | | | | | | | | | | |
|-------------|----|----|--------------|-----|-----|-----|---|-----|-----|-----|-----|-----------------------------|
| Disp 011 | TF | NH | Hop Limit | CID | SAC | SAM | M | DAC | DAM | SCI | DCI | Uncompressed IPv6 fields |
|-------------|----|----|--------------|-----|-----|-----|---|-----|-----|-----|-----|-----------------------------|



| | |
|----|---|
| 00 | ECN+DSCP+4b pad+ 20b Flow label (4 Bytes) |
| 01 | ECN +2b pad + 12b Flow label (2 Bytes), DSCP omitted |
| 10 | ECN+DSCP (1B), Flow label omitted |
| 11 | ECN+DSCP+Flow label omitted |

| SAC DAC | SAM DAM | Address |
|------------|------------|-------------------------------|
| 0 | 00 | No compression |
| 0 | 01 | First 64-bits omitted |
| 0 | 10 | First 112 bits omitted |
| 0 | 11 | 128 bits omitted. Get from L2 |
| 1 | 00 | Unspecified Address :: |
| 1 | 01 | First 64 bits from context |
| 1 | 10 | First 112 bits from context |
| 1 | 11 | 128 bits from context and L2 |

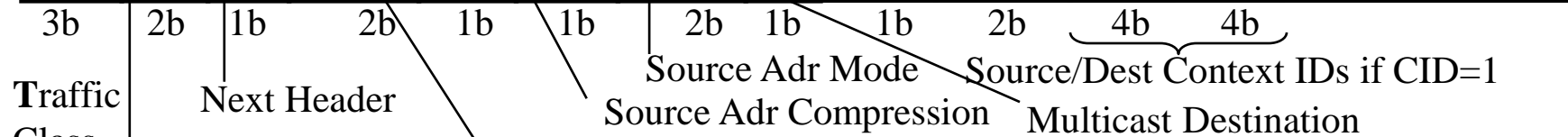
Student Questions

- ❑ The hop limit of 64 is represented as 10? is this correct? What does CID represent?
CID=0 ⇒ SCI/DCI absent
- ❑ If I want to broadcast, then what field should I change?
M=1, DAC=0, DAM=11
or M=1, DAC=1, DAM=11
- ❑ I am confused as to why it is compressed. Could you show an example that the uncompressed bits take more space to represent the same information?
Addresses are only a few bits long. IPv4 addresses are 128-bit long.
- ❑ SAC and DAC are 1 bit, so 2 bits in total. Why does the table use only 1 bit to represent DAC and SAC?
Their interpretation is the identical/similar.
- ❑ Same question with SAM and DAM.
Same as above.
- ❑ Why is it safe to omit some bits?
Those bits can be reconstructed.

Context Based Compression

- ❑ HC1 works only with **link-local** addresses
- ❑ Need globally routable IPv6 addresses for outside nodes
- ❑ IPHC uses a 3b dispatch code and a 13-bit base header

| | | | | | | | | | | | | |
|-------------|----|----|--------------|-----|-----|-----|---|-----|-----|-----|-----|-----------------------------|
| Disp 011 | TF | NH | Hop Limit | CID | SAC | SAM | M | DAC | DAM | SCI | DCI | Uncompressed IPv6 fields |
|-------------|----|----|--------------|-----|-----|-----|---|-----|-----|-----|-----|-----------------------------|



Traffic Class, Flow Label

Next Header uses LowPAN_NHC

Predefined hop limit = uncompressed (00), 1, 64, 255

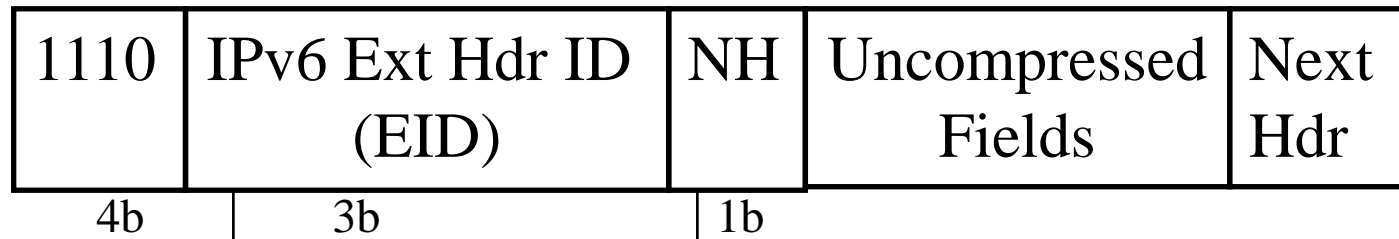
| | |
|----|---|
| 00 | ECN+DSCP+4b pad+ 20b Flow label (4 Bytes) |
| 01 | ECN +2b pad + 12b Flow label (2 Bytes), DSCP omitted |
| 10 | ECN+DSCP (1B), Flow label omitted |
| 11 | ECN+DSCP+Flow label omitted |

| SAC DAC | SAM DAM | Address |
|------------|------------|-------------------------------|
| 0 | 00 | No compression |
| 0 | 01 | First 64-bits omitted |
| 0 | 10 | First 112 bits omitted |
| 0 | 11 | 128 bits omitted. Get from L2 |
| 1 | 00 | Unspecified Address :: |
| 1 | 01 | First 64 bits from context |
| 1 | 10 | First 112 bits from context |
| 1 | 11 | 128 bits from context and L2 |

Student Questions

Context Based Compression (Cont)

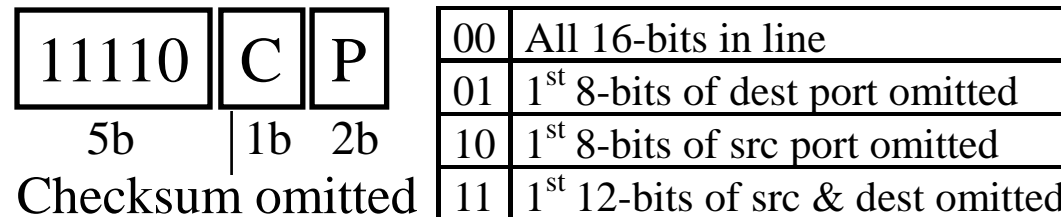
- ❑ If the next header uses LowPAN_NHC
 - For IPv6 base extension headers:



0 = Uncompressed
1 = LowPAN_NHC encoded

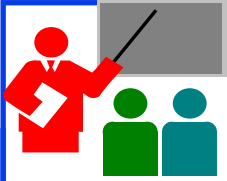
| EID | Header |
|-----|--------------------------|
| 0 | IPv6 Hop-by-Hop Options |
| 1 | IPv6 Routing |
| 2 | IPv6 Fragment |
| 3 | IPv6 Destination Options |
| 4 | IPv6 Mobility Header |
| 5 | Reserved |
| 6 | Reserved |
| 7 | IPv6 Header |

LowPAN_NHC UDP Header:



Student Questions

Ref: J. Hui and P. Thubert, "Compression Format for IPv6 Datagrams over IEEE 802.15.4-Based Networks," IETF RFC 6282, Sep 2011, <http://tools.ietf.org/pdf/rfc6282>



6LowPAN: Summary

❑ 3 New Headers:

- Mesh addressing: Intermediate addresses
- Hop-by-Hop: Mesh broadcasts
- Destination processing: Fragmentation

❑ Address Formation: 128-bit addresses by prefixing FE80::

❑ Header compression:

- HC1+HC2 header for link-local IPv6 addresses
- IPHC compression for all IPv6 addresses

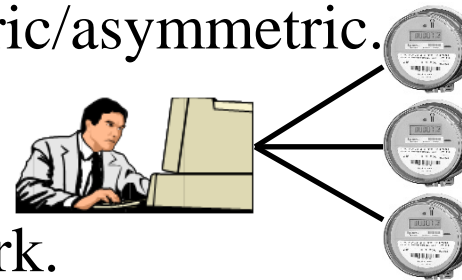
Student Questions

- ❑ If the MTU gets bigger and bigger in the future, does that mean 6LowPAN will become useless?

6LowPAN compresses the IPv6 header, which will remain large even if the MTU is larger.

Routing Protocol for Low-Power and Lossy Networks (RPL)

- ❑ Developed by IETF Routing over Low-Power and Lossy Networks (ROLL) working group
- ❑ Low-power and Lossy Networks (LLN) Routers have constraints on processing, memory, and energy.
⇒ Can't use OSPF, OLSR, RIP, ~~AODV, DSR~~, etc
- ❑ LLN links have high loss rates, low data rates, and instability
⇒ expensive bits, dynamically formed topology
- ❑ Covers both wireless and wired networks
Requires **bidirectional** links. Maybe symmetric/asymmetric.
- ❑ Ideal for n-to-1 (**data sink**) communications, e.g., meter reading
1-to-n and 1-to-1 possible with some extra work.
- ❑ Multiple LLN instances on the same physical networks



Ref: T. Winder, Ed., et al., "RPL: IPv6 Routing Protocol for Low-Power and Lossy Networks," IETF RFC 6550, Mar 2012,
<https://ietf.org/doc/rfc6550/>

Student Questions

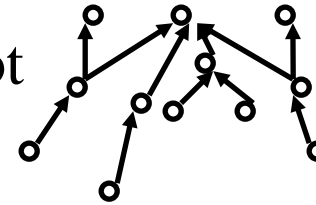
- ❑ You stated, "Can't use OSPF, OLSR, RIP, AODV, DSR." However, ZigBee, which is used in IoT, uses AODV. You also mentioned that in your wireless course (Slide 13-19).

You are right. This slide has been corrected.

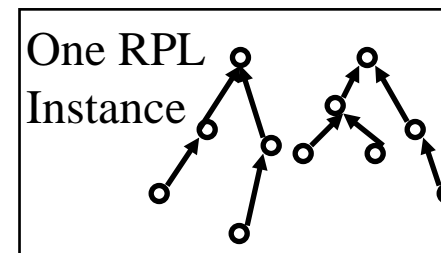
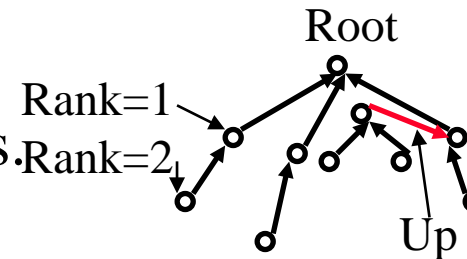
RPL Concepts

- ❑ **Directed Acyclic Graph (DAG):** No cycles
- ❑ **Root:** No outgoing edge
- ❑ **Destination-Oriented DAG (DODAG):** Single root
- ❑ **Up:** Towards the root
- ❑ **Down:** Away from root
- ❑ **Objective Function:** Minimize energy, latency, ...
- ❑ **Rank:** Distance from the root using specified objective
- ❑ **RPL Instance:** One or more DODAGs.
A node may belong to multiple RPL instances.
- ❑ **DODAG ID:** IPv6 Adr of the root
- ❑ **DODAG Version:** Current version of the DODAG. Every time a new DODAG is computed with the same root, its version incremented.

DAG



DODAG



Student Questions

- ❑ I'm not sure what RPL Instance is.
RPL Instance=DODAG
- ❑ Does 'A node may belong to multiple RPL instances' mean 'A node can be in more than one DODAG'?

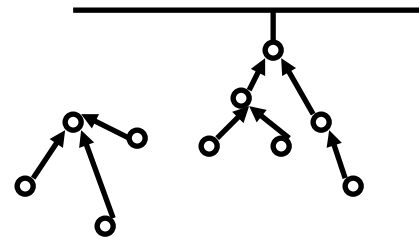
Yes.

- ❑ If it does, is there one more root?

Yes, each DODAG will have its root.

RPL Concepts (Cont)

- ❑ **Goal:** Reachability goal, e.g., connected to a database
- ❑ **Grounded:** Root can satisfy the goal
- ❑ **Floating:** Not grounded. Only in-DODAG communication.
- ❑ **Parent:** Immediate successor towards the root
- ❑ **Sub-DODAG:** Sub tree rooted at this node
- ❑ **Storing:** Nodes keep routing tables for sub-DODAG
- ❑ **Non-Storing:** Nodes know only their parent. Do not keep a routing table.



Student Questions

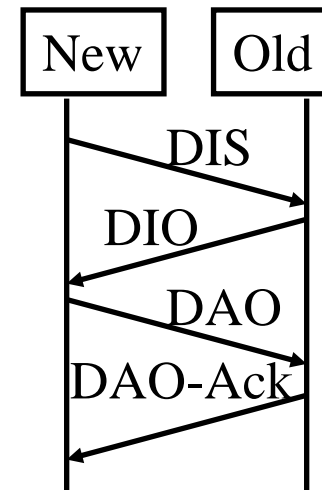
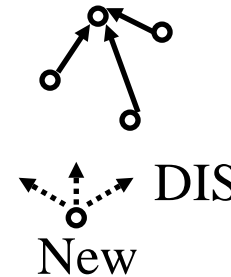
- ❑ Could you please explain the concept "parent" and "sub-DODAG" again?

Sub-DODAG = Subtree

Parent = Next node towards the root

RPL Control Messages

- DODAG Information Object (DIO):**
 - Downward RPL instance multicasts
 - Allows other nodes to discover an RPL instance and join it
- DODAG Information Solicitation (DIS):**
 - Link-local **multicast** request for DIO (neighbor discovery).
Do you know of any DODAGs?
- Destination Advertisement Object (DAO):**
 - From child to parent or root.
 - Can I join you as a child on DODAG #x?
- DAO Ack:** Yes, you can! Or Sorry, you can't!
- Consistency Check:** Challenge/response messages for security



Student Questions

- DIO: an announcement from above; does this mean a parent acknowledges a child? **DIO=Parent is announcing/ broadcasting. Not acking.**
- DIS: request from the bottom. Does this mean a child is looking for a parent? **Yes.**
- The image on this slide is confusing! It looks like a 4-way handshake!
It is a 4-way handshake.
- Is RPL a stateful protocol?
No. Stateful = one packet determines the fate of the next.
- Can you explain slide 20 again? **Sure.**

Ref: S. Kuryla, "RPL:IPv6 Routing Protocol for Low Power and Lossy Networks,"

<http://cnds.eecs.jacobs-university.de/courses/nds-2010/kuryla-rpl.pdf>

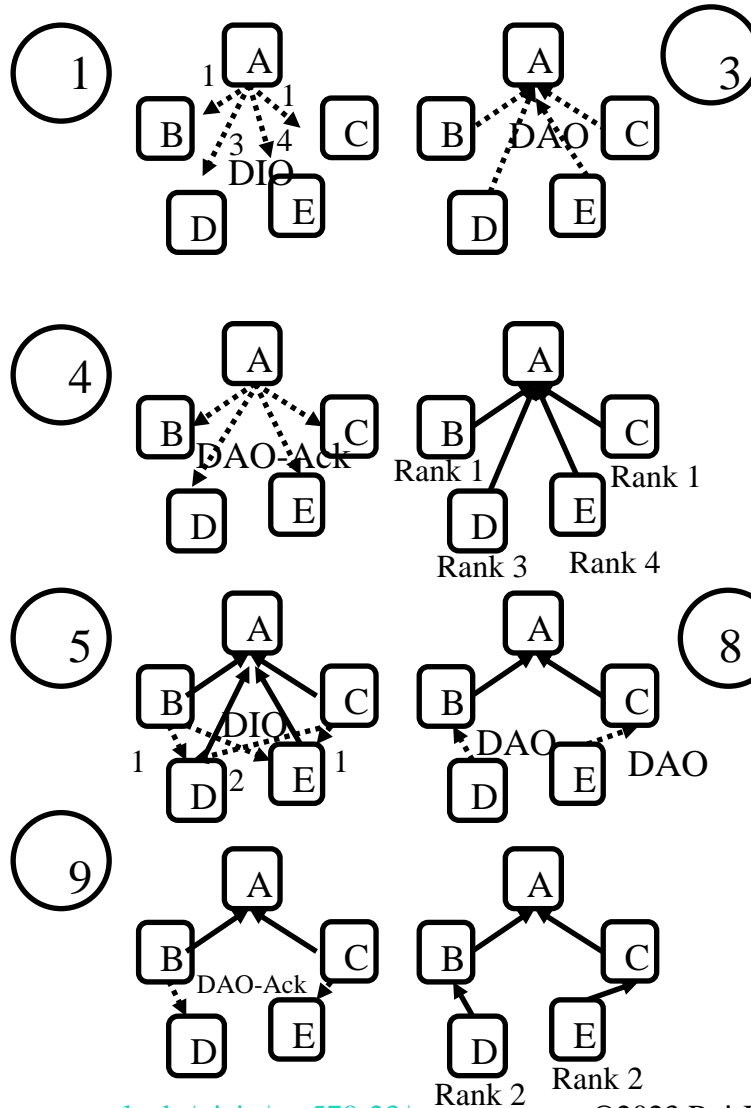
Washington University in St. Louis

<http://www.cse.wustl.edu/~jain/cse570-23/>

©2023 Raj Jain

DODAG Formation Example

1. A multicasts DIOs that it's a member of DODAG ID itself with Rank 0.
2. B, C, D, and E hear and determine that their rank (distance) is 1, 1, 3, 4, respectively, from A
3. B, C, D, E send DAOs to A.
4. A accepts all
5. B and C multicast DIOs
6. D hears those and determines that its distance from B and C is 1, 2
7. E hears both B, and C and determines that its distance from B and C is 2, 1
8. D sends a DAO to B
E sends a DAO to C
9. B sends a DAO-Ack to D
C sends a DAO-Ack to E

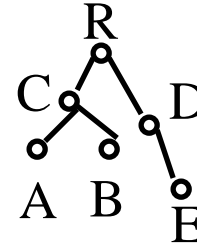


Student Questions

- Does it matter how we break the tie here?
Yes. The final answers will be different. But one answer is equal to the other. For consistency, choose the smaller number or alphabet.
- Could you explain this example in more detail, please? *Sure.*
- In number #6, why are the distances from D to B, C 1,2 each? Is it an assumption?
No. Distance is guessed from the signal strength.

RPL Data Forwarding

- ❑ Case 1: To the root (n-to-1)
 - Address to root and give to parent
- ❑ Case 2: A to B
 - 2A: Storing (Everyone keeps a routing table)
 - ❑ Forward up from A to the common parent
 - ❑ Forward down from common parent to B
 - 2B: Non-storing (No routing tables except at root)
 - ❑ Forward up from A to root
 - ❑ Root puts a source route and forwards down
- ❑ Case 2: Broadcast from the root (1-to-n)
 - 2A: Storing (everyone knows their children)
 - ❑ Broadcast to children
 - 2B: Non-Storing (Know only parents but not children)
 - ❑ Root puts a source route for each leaf and forwards



Student Questions

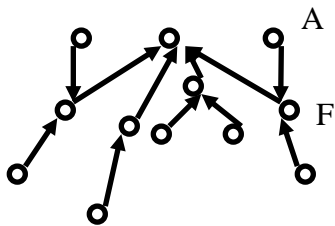
- ❑ Could you explain the difference between Storing and Non-Storing in the Broadcast case?

Non-storing ⇒ Alzheimer. You do not know your children or grandchildren. So, I don't know whether you should forward the packet any further. The root will tell you.

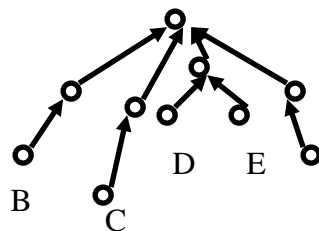
Homework 12B

[10 points]

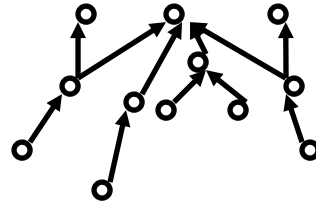
1. Which of the following is not a DODAG, and why?
2. What is the direction of Link A to F? (Up or Down):
3. Assuming each link has a distance of 1, what is the rank of node B?
4. Show the paths from B to C if the DODAG is non-storing.
5. Show the paths from D to E if the DODAG is storing.



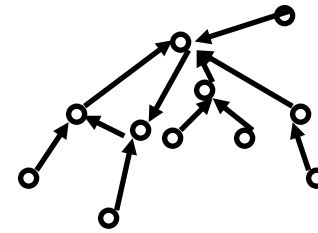
(a)



(b)

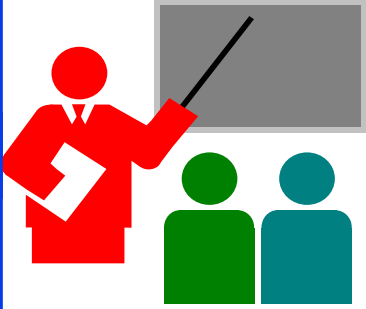


(c)



(d)

Student Questions

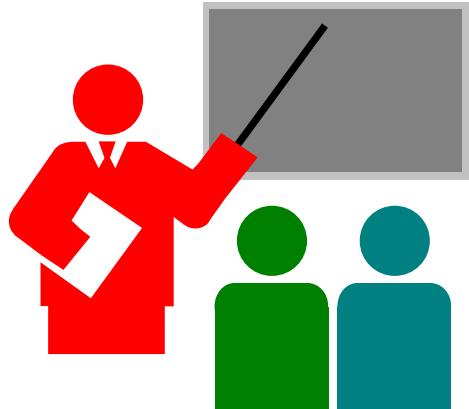


RPL Summary

1. An RPL instance consists of one or more DODAGs
2. DIO are broadcast downward,
DAOs are requests to join upward.
DIS are DIO solicitations
DAO-ack are responses to DAO
3. Non-storing nodes do not keep any routing table and send everything upwards toward the root

Student Questions

Summary



1. 6LowPAN is designed for IPv6 over IEEE 802.15.4
Frame size and address sizes are primary issues.
Header compression is the key mechanism
2. RPL is designed primarily for data collection
There is no assumption about IEEE 802.15.4 or wireless or
frame size
Routing is the primary issue.
Forming a spanning tree-like DODAG is the solution.

Student Questions

Reading List

- ❑ O. Hersent et al., “The Internet of Things: Key Applications and Protocols,” Wiley, 2013, 344 pp., ISBN: 9781119994350 (Safari Book)
- ❑ G. Montenegro, et al., “Transmission of IPv6 Packets over IEEE 802.15.4 Networks,” RFC 4944, Sep 2007, <http://tools.ietf.org/pdf/rfc4944>
- ❑ J. Hui and P. Thubert, “Compression Format for IPv6 Datagrams over IEEE 802.15.4-Based Networks,” IETF RFC 6282, Sep 2011, <http://tools.ietf.org/pdf/rfc6282>
- ❑ T. Winder, Ed., et al., "RPL: IPv6 Routing Protocol for Low-Power and Lossy Networks," IETF RFC 6550, Mar 2012, <https://ietf.org/doc/rfc6550/>
- ❑ S. Kuryla, “RPL: IPv6 Routing Protocol for Low Power and Lossy Networks,” <http://cnds.eecs.jacobs-university.de/courses/nds-2010/kuryla-rpl.pdf>

Student Questions

Wikipedia Links

- ❑ <http://en.wikipedia.org/wiki/6LoWPAN>
- ❑ http://en.wikipedia.org/wiki/IEEE_802.15.4
- ❑ http://en.wikipedia.org/wiki/MAC_address
- ❑ <http://en.wikipedia.org/wiki/IPv6>
- ❑ http://en.wikipedia.org/wiki/IPv6_address
- ❑ http://en.wikipedia.org/wiki/Organizationally_unique_identifier
- ❑ http://en.wikipedia.org/wiki/IPv6_packet
- ❑ http://en.wikipedia.org/wiki/Link-local_address

Student Questions

References

- ❑ N. Kushalnagar et al., "IPv6 over Low-Power Wireless Personal Area Networks (6LoWPANs): Overview, Assumptions, Problem Statement, and Goals", IETF RFC 4919, Aug 2007, <http://www.rfc-editor.org/rfc/pdf/rfc4919.txt.pdf>
- ❑ G. Montenegro, N. Kushalnagar, J. Hui, D. Culler, "Transmission of IPv6 Packets over IEEE 802.15.4 Networks," IETF RFC 4944, <https://tools.ietf.org/pdf/rfc4944>
- ❑ J. Hui, Ed., P. Thubert, "Compression Format for IPv6 Datagrams over IEEE 802.15.4-Based Networks," IETF RFC 6282, Sept 2011, <https://tools.ietf.org/html/rfc6282>
- ❑ E. Kim, et al., "Design and Application Spaces for IPv6 over Low-Power Wireless Personal Area Networks (6LoWPANs)," IETF RFC 6568, Apr 2012, <http://www.rfc-editor.org/rfc/pdf/rfc6568.txt.pdf>
- ❑ E. Kim et al., "Problem Statement and Requirements for IPv6 over Low-Power Wireless Personal Area Network (6LoWPAN) Routing," IETF RFC 6606, May 2012, <http://www.rfc-editor.org/rfc/pdf/rfc6606.txt.pdf>
- ❑ Z. Shelby, et al., "Neighbor Discovery Optimization for IPv6 over Low-Power Wireless Personal Area Networks (6LoWPANs), IETF RFC 6775, Nov. 2012, <http://www.rfc-editor.org/rfc/pdf/rfc6775.txt.pdf>

Student Questions

References (Cont)

- ❑ "Routing Requirements for Urban Low-Power and Lossy Networks," IETF RFC 5548, May 2009, <https://ietf.org/doc/rfc5548/>
- ❑ "Industrial Routing Requirements in Low-Power and Lossy Networks," IETF RFC 5673, Oct 2009, <https://ietf.org/doc/rfc5673/>
- ❑ "Home Automation Routing Requirements in Low-Power and Lossy Networks," IETF RFC 5826, Apr 2010, <https://ietf.org/doc/rfc5826/>
- ❑ "Building Automation Routing Requirements in Low-Power and Lossy Networks," IETF RFC 5867, Jun 2010, <https://ietf.org/doc/rfc5867/>
- ❑ "The Trickle Algorithm," IETF RFC 6206, Mar 2011, <https://ietf.org/doc/rfc6206/>
- ❑ "RPL: IPv6 Routing Protocol for Low-Power and Lossy Networks," IETF RFC 6550, Mar 2012, <https://ietf.org/doc/rfc6550/>
- ❑ "Routing Metrics Used for Path Calculation in Low-Power and Lossy Networks," IETF RFC 6551, Mar 2012, <https://ietf.org/doc/rfc6551/>
- ❑ "Objective Function Zero for the Routing Protocol for Low-Power and Lossy Networks (RPL)," IETF RFC 6552, Mar 2012, <https://ietf.org/doc/rfc6552/>

Student Questions

References (Cont)

- ❑ "The Minimum Rank with Hysteresis Objective Function," IETF RFC 6719, Sep 2012, <https://ietf.org/doc/rfc6719/>
- ❑ "Reactive Discovery of Point-to-Point Routes in Low-Power and Lossy Networks," IETF RFC 6997, Aug 2013, <https://ietf.org/doc/rfc6997/>
- ❑ "A Mechanism to Measure the Routing Metrics along a Point-to-Point Route in a Low-Power and Lossy Network," IETF RFC 6998, Aug 2013, <https://ietf.org/doc/rfc6998/>

Student Questions

Acronyms

- ❑ 6LowPAN IPv6 over Low Power Wireless Personal Area Network
- ❑ AODV Ad-hoc On-demand Distance Vector
- ❑ AQMP Advanced Queueing Message Protocol
- ❑ ARC-EM4 Name of a product
- ❑ ARM Acorn RISC Machine
- ❑ CC Consistency Check
- ❑ CID Context ID
- ❑ CoAP Constrained Application Protocol
- ❑ CoRE Constrained Restful Environment
- ❑ DA Destination Address
- ❑ DAC Destination Address Compression
- ❑ DAG Directed Acyclic Graph
- ❑ DAM Destination Address Mode
- ❑ DAO DODAG Advertisement Object
- ❑ DCI Destination Context ID
- ❑ DDS Data Distribution Service

Student Questions

Acronyms (Cont)

- ❑ DECT Digital Enhanced Cordless Telecommunication
- ❑ DIO DODAG Information Object
- ❑ DIS DODAG Information Solicitation
- ❑ DODAG Destination Oriented Directed Acyclic Graph
- ❑ DSCP Differentiated Services Control Point
- ❑ DSR Dynamic Source Routing
- ❑ DTLS Datagram Transport Level Security
- ❑ ECN Explicit Congestion Notification
- ❑ EID IPv6 Extension Header ID
- ❑ EUI Extended Unique ID
- ❑ GP GreenPHY
- ❑ HC Header Compression
- ❑ HC1-HC2 Header Compression 1 and Header Compression 2
- ❑ ICMP IP Control Message Protocol
- ❑ ID Identifier
- ❑ IEEE Institution of Electrical and Electronic Engineers

Student Questions

Acronyms (Cont)

- ❑ IETF Internet Engineering Task Force
- ❑ IID Interface Identifier
- ❑ IoT Internet of Things
- ❑ IP Internet Protocol
- ❑ IPHC IP Header Compression
- ❑ IPv6 Internet Protocol Version 6
- ❑ ISASecure Security certification by
- ❑ LLN Low-Power and Lossy Networks
- ❑ LoRaWAN Long Range Wide Area Network
- ❑ LTE Long-Term Evolution
- ❑ MAC Media Access Control
- ❑ MTU Maximum Transmission Unit
- ❑ NFC Near Field Communication
- ❑ NH Next Header
- ❑ NHC Next Header Compression
- ❑ OLSR On-Demand Link State Routing

Student Questions

Acronyms (Cont)

- ❑ OSPF Open Shortest Path Forwarding
- ❑ PAN Personal Area Network
- ❑ RFC Request for Comments
- ❑ RIP Routing Information Protocol
- ❑ ROLL Routing over Low-Power and Lossy Networks
- ❑ RPL Routing Protocol for Low-Power and Lossy Networks
- ❑ SA Source Address
- ❑ SAC Source Address Compression
- ❑ SAM Source Address Mode
- ❑ SASL Simple Authentication and Security Layer
- ❑ SCI Source Context ID
- ❑ SMACK Simplified Mandatory Access Control Kernel
- ❑ TCG Trusted Computing Group
- ❑ TCP Transmission Control Protocol
- ❑ TF Traffic Class, Flow Label
- ❑ TinyOS Tiny Operating System

Student Questions

Acronyms (Cont)

- ❑ UDP User Datagram Protocol
- ❑ ULE Ultra Low Energy
- ❑ Wi-Fi Wireless Fidelity
- ❑ WirelessHART Wireless Highway Addressable Remote Transducer Protocol
- ❑ WPAN Wireless Personal Area Network

Student Questions

Scan This to Download These Slides



Raj Jain

<http://rajjain.com>

http://www.cse.wustl.edu/~jain/cse570-23/m_12lpn.htm

Student Questions

Related Modules



CSE567M: Computer Systems Analysis (Spring 2013),

https://www.youtube.com/playlist?list=PLjGG94etKypJEKjNAa1n_1X0bWWNyZcof

CSE473S: Introduction to Computer Networks (Fall 2011),

https://www.youtube.com/playlist?list=PLjGG94etKypJWOSPMh8Azcg5e_10TiDw



Wireless and Mobile Networking (Spring 2016),

https://www.youtube.com/playlist?list=PLjGG94etKypKeb0nzyN9tSs_HCd5c4wXF

CSE571S: Network Security (Fall 2011),

<https://www.youtube.com/playlist?list=PLjGG94etKypKvzfVtutHcPFJXumyyg93u>



Video Podcasts of Prof. Raj Jain's Lectures,

<https://www.youtube.com/channel/UCN4-5wzNP9-ruOzQMs-8NUw>

Student Questions