LAN Extension and Network Virtualization in Cloud Data Centers



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These slides and audio/video recordings of this class lecture are at: http://www.cse.wustl.edu/~jain/cse570-21/

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- 1. TRILL
- 2. NVGRE
- 3. VXLAN
- 4. NVO3
- 5. Geneve
- 6. EVPN
- 7. GUE

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Geographic Clusters of Data Centers

- Multiple data centers are used to improve availability
- □ Cold-Standby: Data is backed up on tapes and stored off-site.
 In case of disaster, application and data are loaded in standby.
 Manual switchover ⇒ Significant downtime. (1970-1990)
- □ Hot-Standby: Two servers in different geographically close data centers exchange state and data continuously.
 Synchronous or Asynchronous data replication to standby.
 On a failure, the application automatically switches to standby.
 Automatic switchover ⇒ Reduced downtime (1990-2005)
 Only 50% of resources are used under normal operation.
- □ Active-Active: All resources are used. Virtual machines and data can be quickly moved between sites, when needed.

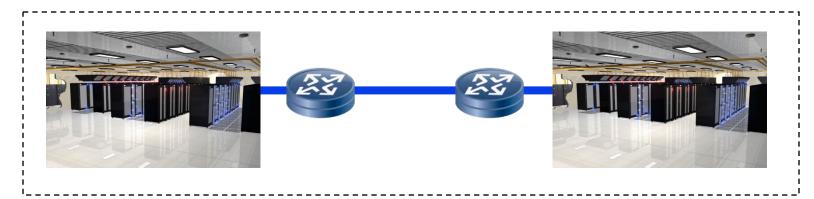
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Ref: G. Santana, "Datacenter Virtualization Fundamentals," Cisco Press, 2014, ISBN: 1587143240
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Data Center Interconnection (DCI)

- Allows distant data centers to be connected in one L2 domain
 - > Distributed applications
 - > Disaster recovery
 - > Maintenance/Migration
 - > High-Availability
 - > Consolidation
- □ Active and standby can share the same virtual IP for switchover.
- Multicast can be used to send state to multiple destinations.



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Challenges of LAN Extension

- **Broadcast storms**: Broadcast, Unknown and Multicast (BUM) may create excessive flood
- Loops: Easy to form loops in a large network.
- □ STP Issues:
 - > High spanning tree diameter (leaf-to-leaf): More than 7.
 - Root can become bottleneck and a single point of failure
 - > Multiple paths remain unused
- **Tromboning**: Dual attached servers and switches generate excessive cross traffic
- Security: Data on LAN extension must be encrypted

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Aggregation

Access

Server

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Ausgregation

Access

Server

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TRILL

- □ Transparent Interconnection of Lots of Links
- Allows a large campus to be a single extended LAN
- LANs allow free mobility inside the LAN but:
 - > Inefficient paths using Spanning tree
 - > Inefficient link utilization since many links are disabled
 - > Inefficient link utilization since multipath is not allowed.
 - ➤ Unstable: small changes in network ⇒ large changes in spanning tree
- □ IP subnets are not good for mobility because IP addresses change as nodes move and break transport connections, but:
 - > IP routing is efficient, optimal, and stable
- □ Solution: Take the best of both worlds
 - ⇒ Use MAC addresses and IP routing

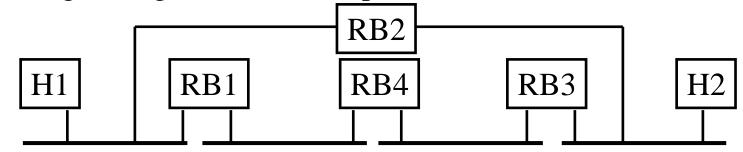
Ref: RFCs 5556, 6325, 6326, 6327, 6361, 6439

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TRILL Architecture

- □ Routing Bridges (RBridges) encapsulate L2 frames and route them to destination RBridges which decapsulate and forward
- Header contains a hop-limit to avoid looping
- RBridges run IS-IS to compute pair-wise optimal paths for unicast and distribution trees for multicast
- RBridge learn MAC addresses by source learning and by exchanging their MAC tables with other RBridges
- Each VLAN on the link has one (and only one) designated RBridge using IS-IS election protocol



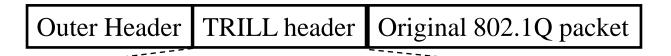
Ref: R. Perlman, "RBridges: Transparent Routing," Infocom 2004

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TRILL Encapsulation Format



Version		Multi- Destination	_	_	_	_	_
2b	2b	1b	5b	6b	16b	16b	

- For outer headers both PPP and Ethernet headers are allowed. PPP for long haul.
- Outer Ethernet header can have a VLAN ID corresponding to the VLAN used for TRILL.
- □ Priority bits in outer headers are copied from inner VLAN

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TRILL Features

- Transparent: No change to capabilities.
 Broadcast, Unknown, Multicast (BUM) support. Auto-learning.
- Zero Configuration: RBridges discover their connectivity and learn MAC addresses automatically
- Hosts can be multi-homed
- VLANs are supported
- Optimized route
- No loops
- Legacy bridges with spanning tree in the same extended LAN

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TRILL: Summary

- □ TRILL allows a large campus to be a single Extended LAN
- □ Packets are encapsulated and routed using IS-IS routing

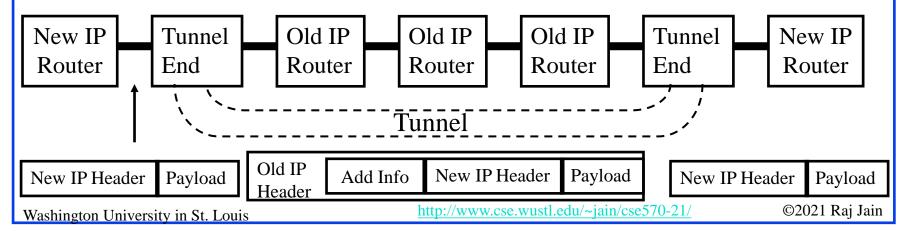
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GRE

- Any new feature in IP requires *encapsulation*, a.k.a. *tunneling*
- □ Generic Routing Encaptulation (RFC 1701/1702)
- \square Generic \Rightarrow X over Y for any X or Y protocols
- □ Given n protocols, we need $O(n^2)$ encapsulation formats, GRE converts this to O(1) format.
- Encapsulations may require the following services:
 - > Stream multiplexing: Which recipient at the other end?
 - > Source Routing: what path to take?
 - > Packet Sequencing



GRE (Cont)

- □ GRE provides all of the above encapsulation services
- Over IPv4, GRE packets use a protocol type of 47
- Optional: Checksum, Loose/strict Source Routing, Key
- Key is used either to authenticate the source or to distinguish different substreams
- Recursion Control: # of additional encapsulations allowed. $0 \Rightarrow$ Restricted to a single provider network \Rightarrow end-to-end
- Offset: Points to the next source route field to be used
- □ IP or IPSec are commonly used as delivery headers

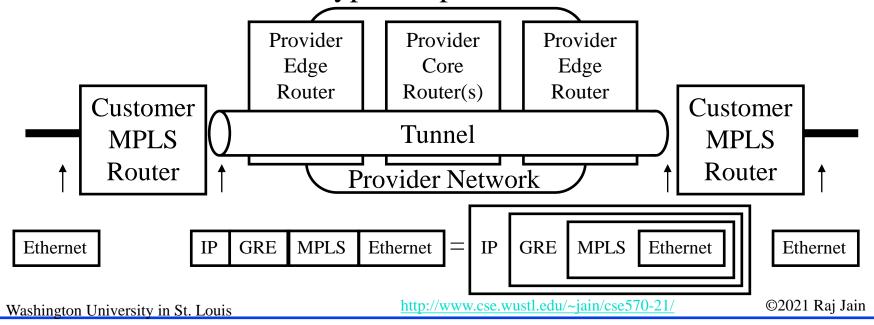
Delivery Header	GRE Header	Payload

Check- sum Present	Routing Present	Present	Seq. # Present	Source	Recursion Control	Flags	Ver. #	Prot. Type			(Opt)	#	Source Routing List (Opt)
1b Washingto	1b 1b 1b 1b 1 University in St. Louis		3b	5b ://www.	3b 16b 16b 16b ww.cse.wustl.edu/~jain/cse570-21/				32b	32b ©20	Variable 21 Raj Jain		

EoMPLSoGRE

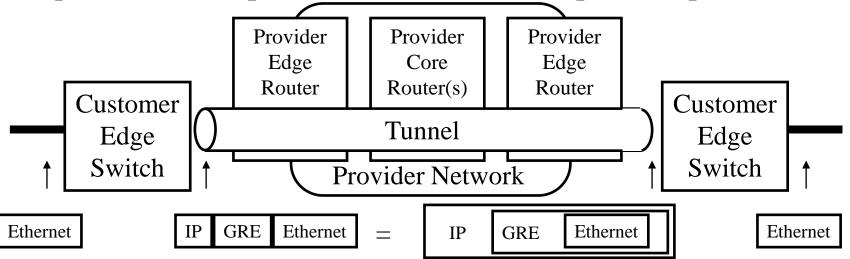
- Ethernet over MPLS over GRE (point-to-point)

 VPLS over MPLS over GRE (Multipoint-to-multipoint)
- Used when provider offers only L3 connectivity Subscribers use their own MPLS over GRE tunnels
- VPLSoGRE or Advanced-VPLSoGRE can also be used
- □ GRE offers IPSec encryption option



NVGRE

- Ethernet over GRE over IP (point-to-point)
- □ A unique 24-bit Virtual Subnet Identifier (VSID) is used as the lower 24-bits of GRE key field \Rightarrow 2²⁴ tenants can share
- □ Unique IP multicast address is used for BUM (Broadcast, Unknown, Multicast) traffic on each VSID
- Equal Cost Multipath (ECMP) allowed on point-to-point tunnels



Ref: P. Garg, Y. Wang, et al. "NVGRE: Network Virtualization Using Generic Routing Encapsulation Encapsulation", RFC 7637, IETF, September 2015.

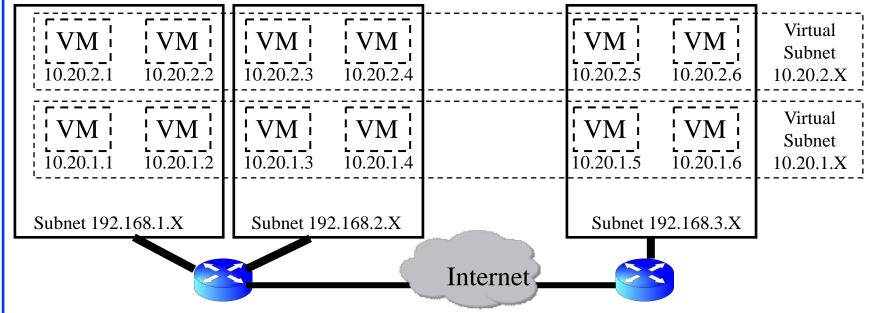
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NVGRE (Cont)

- ☐ In a cloud, a pSwitch or a vSwitch can serve as tunnel endpoint
- □ VMs need to be in the same VSID to communicate
- VMs in different VSIDs can have the same MAC address
- □ Inner IEEE 802.1Q tag, if present, is removed.



Ref: Emulex, "NVGRE Overlay Networks: Enabling Network Scalability," Aug 2012, 11pp.,

 $\underline{\text{http://www.emulex.com/artifacts/074d492d-9dfa-42bd-9583-69ca9e264bd3/elx_wp_all_nvgre.pdf}$

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NVO3

- Network Virtualization Overlays using L3 techniques
- □ Problem: Data Center Virtual Private Network (DCVPN) in a multi-tenant datacenter
- ☐ Issues:
 - > Scale in Number of Networks: Hundreds of thousands of DCVPNs in a single administrative domain
 - > Scale in Number of Nodes: Millions of VMs
 - > VM (or pM) Migration
 - > Support both L2 and L3 VPNs
 - > Dynamic provisioning
 - > Addressing independence: Each tenant should be able select its address space
 - ➤ Virtual Private ⇒ Other tenants do not see your frames
 - > Optimal Forwarding: VMs should not be tied to a single designated router that may be far away.

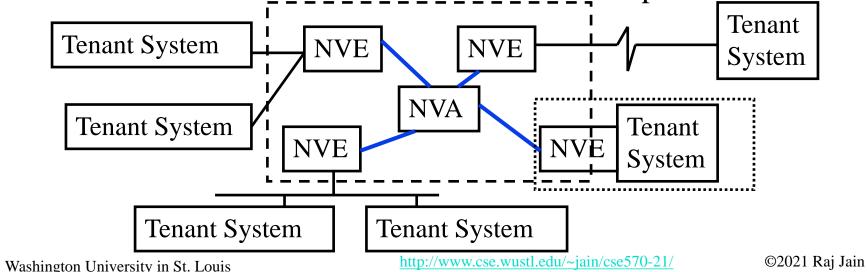
Ref: T. Narten, Ed., "Problem Statement: Overlays for Network Virtualization," IETF RFC 7364, Oct 14, 23 pp.

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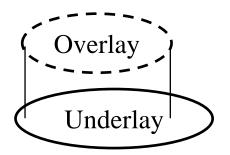
NVO3 Terminology

- ☐ Tenant System (TS): VM or pM
- □ Virtual Network (VN): L2 or L3 Tenant networks
- Network Virtualization Edges (NVEs): Entities connecting TSs (virtual/physical switches/routers)
- □ NVEs could be in vSwitches, external pSwitches or span both.
- Network Virtualization Authority (NVA): Manages forwarding info for a set of NVEs
- NVA could be distributed or centralized and replicated.



NVO3 Terminology (Cont)

- □ Virtual Network (VN): Provides L2/L3 services to a set of tenants
- VN Context ID: A field in the header that identifies a VN instance (VNI).
- □ Overlay header = inner header = Virtual Network Header
- □ Underlay header = outer header = Physical Network Header
- □ **Tenant Separation**: A tenant's traffic cannot be seen by another tenant



Student Questions

Current NVO Technologies

- BGP/MPLS IP VPNs: Widely deployed in enterprise networks. Difficult in data centers because hosts/hypervisors do not implement BGP.
- BGP/MPLS Ethernet VPNs: Deployed in carrier networks. Difficult in data centers.
- **802.1Q**, PB, PBB VLANs
- □ **Shortest Path Bridging**: IEEE 802.1aq
- □ Virtual Station Interface (VSI) Discovery and Configuration Protocol (VDP): IEEE 802.1Qbg
- Address Resolution for Massive numbers of hosts in the Data Center (ARMD): RFC6820
- **TRILL**
- L2VPN: Provider provisioned L2 VPN
- □ Proxy Mobile IP: Does not support multi-tenancy
- □ LISP: RFC 6830

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VXLAN

- □ Virtual eXtensible Local Area Networks (VXLAN)
- L3 solution to isolate multiple tenants in a data center (L2 solution is Q-in-Q and MAC-in-MAC)
- Developed by VMware. Supported by many companies in IETF NVO3 working group
- □ Problem:
 - > 4096 VLANs are not sufficient in a multi-tenant data center
 - ➤ Tenants need to control their MAC, VLAN, and IP address assignments ⇒ Overlapping MAC, VLAN, and IP addresses
 - ➤ Spanning tree is inefficient with large number of switches
 ⇒ Too many links are disabled
 - > Better throughput with IP equal cost multipath (ECMP)

Ref: M. Mahalingam, D. G. Dutt, et al. "VXLAN: A Framework for Overlaying Virtualized Layer 2 Networks over Layer 3 Networks," IETF RFC 7348, August 2014.

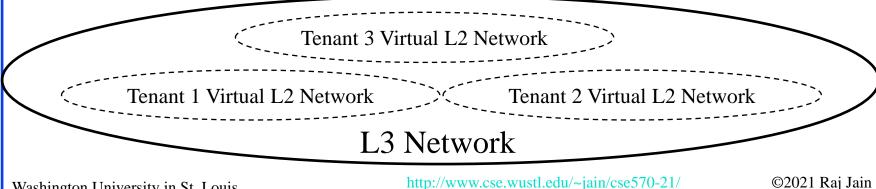
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VXLAN Architecture

- □ Create a virtual L2 overlay (called VXLAN) over L3 networks
- 2²⁴ Virtual Network Instances (VNIs)
- Only VMs in the same VXLAN can communicate
- vSwitches serve as VTEP (VXLAN Tunnel End Point).
 - ⇒ Encapsulate L2 frames in UDP over IP and send to the destination VTEP(s).
- Segments may have overlapping MAC addresses and VLANs but L2 traffic never crosses a VNI

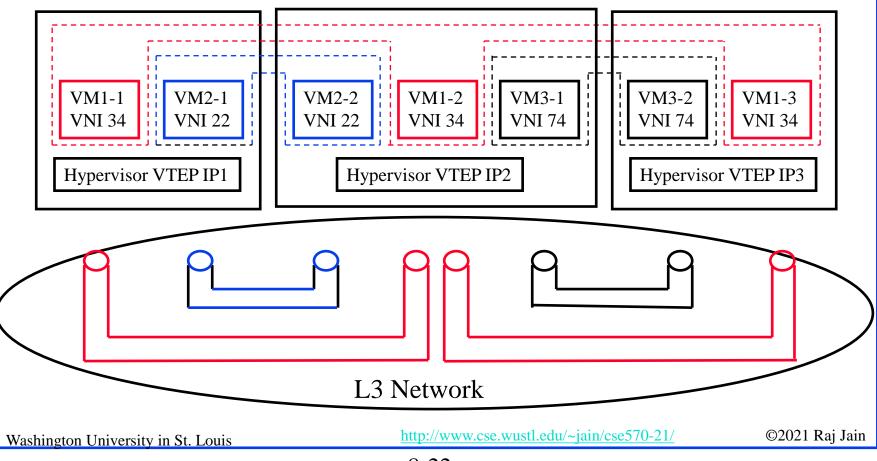


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VXLAN Deployment Example

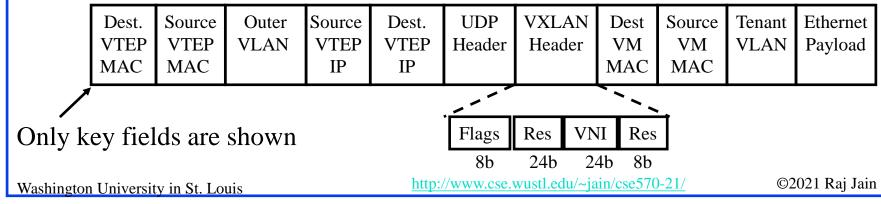
Example: Three tenants. 3 VNIs. 4 Tunnels for unicast.

+ 3 tunnels for multicast (not shown)



VXLAN Encapsulation Format

- Outer VLAN tag is optional.
 Used to isolate VXLAN traffic on the LAN
- □ Source VM ARPs to find Destination VM's MAC address. All L2 multicasts/unknown are sent via IP multicast. Destination VM sends a standard IP unicast ARP response.
- Destination VTEP learns Inner-Src-MAC-to-Outer-Src-IP mapping ⇒ Avoids unknown destination flooding for returning responses



VXLAN Encapsulation Format (Cont)

- ☐ Internet Group Multicast Protocol (IGMP) is used to prune multicast trees
- □ 7 of 8 bits in the flag field are reserved. One flag bit is set if VNI field is valid
- □ UDP source port is a hash of the inner MAC header
 ⇒ Allows load balancing using Equal Cost Multi Path using L3-L4 header hashing
- VMs are unaware that they are operating on VLAN or VXLAN
- VTEPs need to learn MAC address of other VTEPs and of client VMs of VNIs they are handling.
- A VXLAN gateway switch can forward traffic to/from non-VXLAN networks. Encapsulates or decapsulates the packets.

Student Questions



VXLAN: Summary

- VXLAN solves the problem of multiple tenants with overlapping MAC addresses, VLANs, and IP addresses in a cloud environment.
- A server may have VMs belonging to different tenants
- □ No changes to VMs. Hypervisors responsible for all details.
- Uses UDP over IP encapsulation to isolate tenants

Student Questions

Stateless Transport Tunneling Protocol (STT)

- Ethernet over TCP-Like over IP tunnels. GRE, IPSec tunnels can also be used if required.
- Designed for large storage blocks **64kB**. Fragmentation allowed.
- Most other overlay protocols use UDP and disallow fragmentation ⇒ Maximum Transmission Unit (MTU) issues.
- □ TCP-Like: Stateless TCP ⇒ Header identical to TCP (same protocol number 6) but no 3-way handshake, no connections, no windows, no retransmissions, no congestion state
 ⇒ Stateless Transport (recognized by standard port number).
- □ Internet draft expired ⇒ Of historical interest only.
 New work on Geneve.

Ref: B. Davie and J. Gross, "A Stateless Transport Tunneling Protocol for Network Virtualization (STT)," Sep 2013, http://tools.ietf.org/html/draft-davie-stt-04

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Geneve

- Generic Network Virtualization Encapsulation
- Best of NVGRE, VXLAN, and STT
- □ Generic \Rightarrow Can virtualize any (L2/L3/...) protocol over IP
- Tunnel Endpoints: Process Geneve headers and control packets
- □ Transit Device: do not need to process Geneve headers or control packets

Student Questions

Ref: J. Gross, et al, "Geneve: Generic Network Virtualization Encapsulation" IETF Internet Draft, draft-ietf-nvo3-geneve-14, Sep. 12, 2019, https://tools.ietf.org/pdf/draft-ietf-nvo3-geneve-14.txt

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Geneve Frame Format

- □ **Highly Extensible**: Variable number of variable size options
- Any vendor can extend it in its own way by getting an "Option Class" from IANA (Internet Assigned Number Authority)
- □ Options are encoded in a TLV (Type-Length-Value) format

L2IPUDPGeneveGeneveHeaderHeaderHeaderHeaderPayload

		_										
Version	Option Length	OAM Frame	Critical Option Present	Reserved	Payload Protocol Type	Net	rtual twork ID	Reserved	Options	•••	Options	S
2b	6b	1b	1b	бb	16b	2	24 <u>b</u>	8b		```		
						ption Class	Optio Type		ed Value Length		Option Value	
						l 6b	8b	3b	5b			
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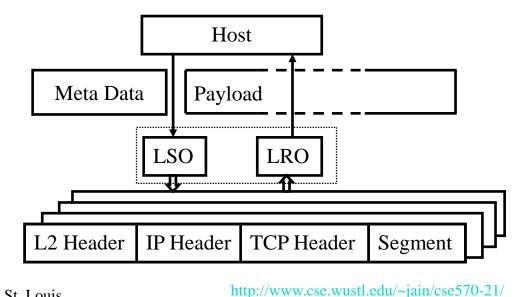
Geneve Frame Format (Cont)

- □ Option Length (6 bits): Length of options field in 4B (does not include the rest of the Geneve header)
- □ OAM Frame (1 bit): Control packet. Does not contain user data. Must be passed on to the control plane CPU
- □ Critical Options Present (1 bits): One or more options are critical.
 - Drop the packet if you don't understand a critical option
- □ Payload Protocol Type (16 bits): 0x6558 for Ethernet
- □ Virtual Network ID (24 bits): Tenant ID
- Option Class (16 bits): Who designed this option. Vendor, technologies, organizations, ...
- \square Option Type (8 bits): msb (most significant bit) =1 => Critical
- □ Option Value Length (5 bits): in units of 4-bytes

Student Questions

LSO and **LRO**

- Large Send Offload (LSO): Host hands a large chunk of data to NIC and meta data. NIC makes MSS size segments, adds checksum, TCP, IP, and MAC headers to each segment.
- □ Large Receive Offload (LRO): NICs attempt to reassemble multiple TCP segments and pass larger chunks to the host. Host does the final reassembly with fewer per packet operations.



Student Questions

Geneve Implementation Issues

- **Fragmentation**: Use Path MTU (Maximum Transmission Unit) discovery to avoid fragmentation on the path
- □ DSCP (Differentiated Services Control Point): DSCP bits in the outer header may or may not be the same as in the inner header. Decided by the policy of the network service provider
- ECN (Explicit Congestion Notification): ECN bits should be copied from inner header on entry to the tunnel and copied back to the inner header on exit from the tunnel
- **Broadcast and Multicast**: Use underlying networks multicast capabilities if available. Use multiple point to point tunnels if multicast is not available.

Student Questions

Geneve Implementation Issues (Cont)

- LSO: Replicate all Geneve headers and options on all outgoing packets.
- □ LRO: Merge all packets with the identical Geneve headers
- □ Option Order: Not significant. Options can be in any order.
- Inner VLAN: Tunnel endpoints decide whether to differentiate packets with different inner VLAN values.

Student Questions

Geneve Summary

- 1. UDP over IP encapsulation
- 2. Geneve header is extensible by vendors
- 3. Generally variable length headers are considered hard for hardware implementation
- 4. Vendor extensibility requires a system to register options and may result in interoperability issues
- 5. All of this is subject to change since it is in the draft stage.

Student Questions

Generic UDP Encapsulation (GUE)

- Using UDP to encapsulate IP protocols
- Allows using efficient hardware implementations of UDP
- □ Generic ⇒ Any IP payload
- □ Optional data in header: VNI, Authentication, Security, congestion control, vendor extensions, etc.
- Allows carrying Non-TCP/Non-UDP IP payloads over networks that filter all non-TCP/non-UDP packets.



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Ref: T. Herbert, L. Yong, and O. Zia, "Generic UDP Encapsulation," IETF Draft, March 7, 2019, draft-ietf-intarea-gue-07
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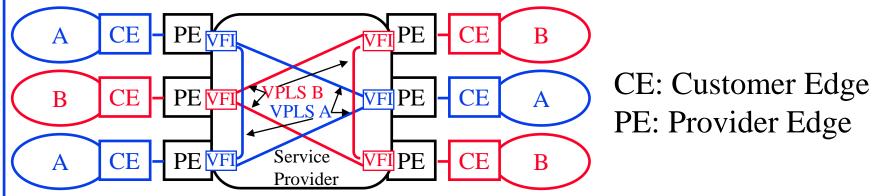
GUE Packet Format

- Source Port: 6080 or any other port
- □ Dest Port: 6080 or any other port
- □ Len: Length of UDP header and payload in 4B unit
- Checksum: Standard UDP checksum
- □ Version: 0 and 1 are standardized. Version 0 shown below.
- □ Hlen: Extenstion Field + Private Data length in 32-bit words Total GUE header length = HLEN*4+4 bytes
- □ Protocol/Control Type: $C=1 \Rightarrow$ Control, $C=0 \Rightarrow$ Protocol type of the payload
- □ Flags indicate the presence of various extension fields and private data
- □ Version 1: No Extension or private data. First 4 bits of IP header are version 0100 or 0110. First 2 bits indicate GUE version.

Source Port	Dest Port	Len	Check sum	Ver	С	Hlen	Proto/ Ctype	Flags	Extension Fields (Opt)	Private Data (Opt)
16b	16b	16b	16b	2b	1b	5b	8b	16b		
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Virtual Private LAN Service (VPLS)

- □ Allows *multi-point* Ethernet services over MPLS networks using a *full-mesh point-to-point pseudo-wires*
- □ Virtual Forwarding Instance (VFI): A virtual LER instance in the provider edge router specific to each customer LAN
- □ VPLS Instance: Set of VFIs and PWs connecting them. Creates a single "VLAN" broadcast domain connecting VFIs.
- Widely deployed but does not meet data center requirements
 ⇒ BGP MPLS-Based Ethernet VPN (EVPN)



Ref: G. Santana, "Datacenter Virtualization Fundamentals," Cisco Press, 2014, ISBN: 1587143240

K. Kompella and Y. Rekhter, "Virtual Private LAN Service (VPLS) Using BGP for Auto-Discovery and Signaling," IETF RFC 4761,

Jan 2007, https://tools.ietf.org/pdf/rfc4761
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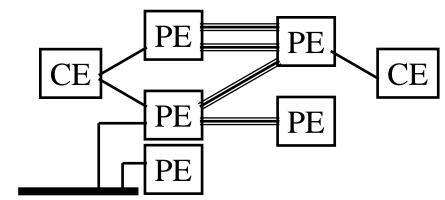
EVPN Features

- Multicast Optimization: MP2MP LSPs
- Ease of Provisioning: Auto-discovery of PEs, Customer Site ID, Automated designated Forwarder election among PEs, MPLS parameters based on VLAN parameters
- New Service Interface: Port = VLAN, Multiple VLANs per port, VLAN bundles are treated as one VLAN
- □ Fast Convergence
- □ Flood Suppression
- □ Flexible VPN topologies and Policies

Student Questions

EVPN Redundancy Features

- □ Flow-based load balancing: CE connected to multiple PEs. Select path by hashing 7-tuple (L2, L3, L4 addresses, VLAN)
- □ Flow-based multi-pathing: Multiple LSPs between PEs
- □ Geo-Redundant PE Nodes: CE connected to multiple PEs in different PoPs
- Optimal Traffic Forwarding: Single/multi-homed CE to Single/multi-homed
 CE. Packet not forwarded to PEs not connected to the destination CE
- Flexible Redundancy Grouping: PEs are grouped for redundancy.
- Multi-homed Network: Entire Ethernet is connected to multiple PEs.



Ref: J. Rabadan, et al., "Applicability of EVPN to NVO3 Networks," IETF Draft, July 8, 2019, https://tools.ietf.org/html/draft-ietf-nvo3-evpn-applicability

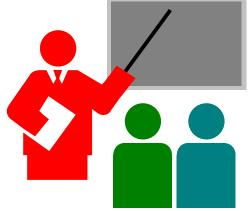
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A. Sajassi, et al., "Requirements for Ethernet VPN (EVPN)," May 2014, https://tools.ietf.org/html/rfc7209
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Student Questions

Summary



- 1. TRILL uses "Routing Bridges" to transport Ethernet packets on a campus network. RBs use IS-IS to find the shortest path.
- 2. NVO3 is a generalized framework for network virtualization and partitioning for multiple tenants over L3. It covers both L2 and L3 connectivity.
- 3. NVGRE uses Ethernet over GRE for L2 connectivity.
- 4. VXLAN uses Ethernet over UDP over IP
- 5. Geneve uses Any protocol over UDP over IP encapsulaton.

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References

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Wikipedia Links

- □ http://en.wikipedia.org/wiki/Generic_Routing_Encapsulation
- http://en.wikipedia.org/wiki/Locator/Identifier_Separation_Protocol
- □ http://en.wikipedia.org/wiki/Large_segment_offload
- □ http://en.wikipedia.org/wiki/Large_receive_offload
- □ https://en.wikipedia.org/wiki/Virtual_Extensible_LAN

Student Questions

Acronyms

□ ARMD Address Resolution for Massive numbers of

hosts in the Data center

ARP Address Resolution Protocol

BGP Border Gateway Protocol

■ BUM Broadcast, Unknown, Multicast

CPU Central Processing Unit

DC Data Center

DCI Data Center Interconnection

DCN Data Center Networks

DCVPN Data Center Virtual Private Network

DSCP Differentiated Services Control Point

□ ECMP Equal Cost Multi Path

■ EoMPLSoGRE Ethernet over MPLS over GRE

■ ECN Explicit Congestion Notification

EVPN Ethernet Virtual Private Network

□ GRE Generic Routing Encapsulation

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■ IANA Internet Address and Naming Authority

□ ID Identifier

□ IEEE Institution of Electrical and Electronic Engineers

□ IETF Internet Engineering Task Force

☐ IGMP Internet Group Multicast Protocol

□ IP Internet Protocol

□ IPSec IP Security

□ IPv4 Internet Protocol V4

□ IS-IS Intermediate System to Intermediate System

■ LAN Local Area Network

LISP Locator ID Separation Protocol

■ LRO Large Receive Offload

□ LSO Large Send Offload

■ MAC Media Access Control

MPLS Multi Protocol Label Switching

MSS Maximum Segment Size

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■ MTU Maximum Transmission Unit

NIC
Network Interface Card

NV Network Virtualization

NVA Network Virtualization Authority

NVEs
Network Virtualization Edge

NVGRE Network Virtualization Using GRE

■ NVO3 Network Virtualization over L3

OAM Operation, Administration and Management

OTV Overlay Transport Virtualization

□ PB Provider Bridges

PBB Provider Backbone Bridge

□ pM Physical Machine

pSwitch Physical Switch

QoS Quality of Service

□ RB Routing Bridge

□ RFC Request for Comment

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□ RS Routing System

□ STT Stateless Transport Tunneling Protocol

TCP Transmission Control Protocol

□ TLV Type-Length-Value

□ TRILL Transparent Routing over Lots of Links

□ TS Tenant System

UDP User Datagram Protocol

□ VDP VSI Discovery and Configuration Protocol

VLAN Virtual Local Area Network

VM Virtual Machine

■ VN Virtual Network

□ VNI Virtual Network Instance/Virtual Network Context ID

VPLS Virtual Private LAN Service

VPLSoGRE Virtual Private LAN Service over GRE

VPN Virtual Private Network

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VRRP Virtual Router Redundancy Protocol

□ VSI Virtual Station Interface

■ VSID Virtual Subnet Identifier

□ vSwitch Virtual Switch

□ VTEP VXLAN Tunnel End Point

□ VXLAN Virtual Extensible Local Area Network

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https://www.youtube.com/playlist?list=PLjGG94etKypJEKjNAa1n_1X0bWWNyZcof

CSE473S: Introduction to Computer Networks (Fall 2011),

https://www.youtube.com/playlist?list=PLjGG94etKypJWOSPMh8Azcgy5e_10TiDw





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Video Podcasts of Prof. Raj Jain's Lectures,

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