Types of Workloads

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Terminology

Test Workloads for Computer Systems

- > Addition Instruction
- Instruction Mixes
- > Kernels
- Synthetic Programs
- Application Benchmarks: Sieve, Ackermann's Function, Debit-Credit, SPEC

Part II: Measurement Techniques and Tools

Measurements are not to provide numbers but insight - Ingrid Bucher

- 1. What are the different types of workloads?
- 2. Which workloads are commonly used by other analysts?
- 3. How are the appropriate workload types selected?
- 4. How is the measured workload data summarized?
- 5. How is the system performance monitored?
- 6. How can the desired workload be placed on the system in a controlled manner?
- 7. How are the results of the evaluation presented?

Terminology

- Test workload: Any workload used in performance studies. Test workload can be real or synthetic.
- Real workload: Observed on a system being used for normal operations.
- **Synthetic workload**:
 - Similar to real workload
 - > Can be applied repeatedly in a controlled manner
 - No large real-world data files
 - No sensitive data
 - > Easily modified without affecting operation
 - Easily ported to different systems due to its small size
 - > May have built-in measurement capabilities.

Test Workloads for Computer Systems

- 1. Addition Instruction
- 2. Instruction Mixes
- 3. Kernels
- 4. Synthetic Programs
- 5. Application Benchmarks

Addition Instruction

- Processors were the most expensive and most used components of the system
- □ Addition was the most frequent instruction

Instruction Mixes

- □ Instruction mix = instructions + usage frequency
- Gibson mix: Developed by Jack C. Gibson in 1959 for IBM

704 systems.	1.	Load and Store	31.2
	2.	Fixed-Point Add and Subtract	6.1
	3.	Compares	3.8
	4.	Branches	16.6
	5.	Floating Add and Subtract	6.9
	6.	Floating Multiply	3.8
	7.	Floating Divide	1.5
	8.	Fixed-point Multiply	0.6
	9.	Fixed-point Divide	0.2
	10.	Shifting	4.4
	11.	Logical, And, Or, etc.	1.6
	12.	Instructions Not Using Registers	5.3
	13.	Indexing	18.0
		Total	100.0
		Iotai	100.0

Instruction Mixes (Cont)

Disadvantages:

- > Complex classes of instructions not reflected in the mixes.
- > Instruction time varies with:
 - □ Addressing modes
 - □ Cache hit rates
 - Dipeline efficiency
 - Interference from other devices during processormemory access cycles
 - Parameter values
 - □ Frequency of zeros as a parameter
 - □ The distribution of zero digits in a multiplier
 - The average number of positions of preshift in floatingpoint add

Number of times a conditional branch is taken

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Instruction Mixes (Cont)

- □ Performance Metrics:
 - > MIPS = Millions of Instructions Per Second
 - » MFLOPS = Millions of Floating Point Operations Per Second

Kernels

- \Box Kernel = nucleus
- □ Kernel= the most frequent function
- Commonly used kernels: Sieve, Puzzle, Tree Searching, Ackerman's Function, Matrix Inversion, and Sorting.
- Disadvantages: Do not make use of I/O devices

Synthetic Programs

- □ To measure I/O performance lead analysts \Rightarrow Exerciser loops
- The first exerciser loop was by Buchholz (1969) who called it a synthetic program.
- A Sample Exerciser: See program listing Figure 4.1 in the book

Synthetic Programs

- □ Advantage:
 - > Quickly developed and given to different vendors.
 - > No real data files
 - > Easily modified and ported to different systems.
 - > Have built-in measurement capabilities
 - > Measurement process is automated
 - > Repeated easily on successive versions of the operating systems
- Disadvantages:
 - ≻ Too small
 - > Do not make representative memory or disk references
 - > Mechanisms for page faults and disk cache may not be adequately exercised.

 - ≻ CPU-I/O overlap may not be representative.
 ≻ Loops may create synchronizations ⇒ better or worse

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Application Benchmarks

- □ For a particular industry: Debit-Credit for Banks
- Benchmark = workload (Except instruction mixes)
- Some Authors: Benchmark = set of programs taken from real workloads
- Popular Benchmarks

Sieve

- Based on Eratosthenes' sieve algorithm: find all prime numbers below a given number n.
- □ Algorithm:

> Write down all integers from 1 to n

- > Strike out all multiples of k, for k=2, 3, ..., \sqrt{n} .
- **Example:**

> Write down all numbers from 1 to 20. Mark all as prime:

<u>1</u>, <u>2</u>, <u>3</u>, <u>4</u>, <u>5</u>, <u>6</u>, <u>7</u>, <u>8</u>, <u>9</u>, <u>10</u>, <u>11</u>, <u>12</u>, <u>13</u>, <u>14</u>, <u>15</u>, <u>16</u>, <u>17</u>, <u>18</u>, <u>19</u>, <u>20</u>

□ Remove all multiples of 2 from the list of primes:

<u>1</u>, <u>2</u>, <u>3</u>, 4, <u>5</u>, 6, <u>7</u>, 8, <u>9</u>, 10, <u>11</u>, 12, <u>13</u>, 14, <u>15</u>, 16, <u>17</u>, 18, <u>19</u>, 20

Sieve (Cont)

- The next integer in the sequence is 3. Remove all multiples of 3:
- <u>1</u>, <u>2</u>, <u>3</u>, 4, <u>5</u>, 6, <u>7</u>, 8, 9, 10, <u>11</u>, 12, <u>13</u>, 14, 15, 16, <u>17</u>, 18, <u>19</u>, 20
- $\Box 5 \ge \sqrt{20} \Longrightarrow \text{Stop}$
- Pascal Program to Implement the Sieve Kernel: See Program listing Figure 4.2 in the book

Ackermann's Function

- To assess the efficiency of the procedure-calling mechanisms.
 The function has two parameters and is defined recursively.
- □ Ackermann(3, n) evaluated for values of n from one to six.

• Metrics:

- > Average execution time per call
- > Number of instructions executed per call, and
- Stack space per call
- □ Verification: Ackermann(3, n) = 2^{n+3} -3
- □ Number of recursive calls in evaluating Ackermann(3,*n*):

 $(512 \times 4^{n-1} - 15 \times 2^{n+3} + 9n + 37)/3$

Execution time per call.

□ Depth of the procedure calls = $2^{n+3}-4$ ⇒ stack space required doubles when n ← n+1.

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Ackermann Program in Simula

□ See program listing Figure 4.3 in the book

Other Benchmarks

- □ Whetstone
- U.S. Steel
- □ LINPACK
- Dhrystone
- Doduc
- **TOP**
- □ Lawrence Livermore Loops
- Digital Review Labs
- Abingdon Cross Image-Processing Benchmark

Debit-Credit Benchmark

- A de facto standard for transaction processing systems.
- □ First recorded in Anonymous et al (1975).
- In 1973, a retail bank wanted to put its 1000 branches, 10,000 tellers, and 10,000,000 accounts online with a peak load of 100 Transactions Per Second (TPS).
- Each TPS requires 10 branches, 100 tellers, and 100,000 accounts.



Debit-Credit Benchmark (Continued)

- □ Metric: price/performance ratio.
- Performance: Throughput in terms of TPS such that 95% of all transactions provide one second or less response time.
- Response time: Measured as the time interval between the arrival of the last bit from the communications line and the sending of the first bit to the communications line.
- Cost = Total expenses for a five-year period on purchase, installation, and maintenance of the hardware and software in the machine room.
- Cost does not include expenditures for terminals, communications, application development, or operations.

Pseudo-code Definition of Debit-Credit

- □ See Figure 4.5 in the book
- □ Four record types: account, teller, branch, and history.
- Fifteen percent of the transactions require remote access
- Transactions Processing Performance Council (TPC) was formed in August 1988.
- □ TPC BenchmarkTM A is a variant of the debit-credit
- Metrics: TPS such that 90% of all transactions provide two seconds or less response time.

SPEC Benchmark Suite

- Systems Performance Evaluation Cooperative (SPEC): Nonprofit corporation formed by leading computer vendors to develop a standardized set of benchmarks.
- Release 1.0 consists of the following 10 benchmarks: GCC, Espresso, Spice 2g6, Doduc, LI, Eqntott, Matrix300, Fpppp, Tomcatv
- □ Primarily stress the CPU, Floating Point Unit (FPU), and to some extent the memory subsystem ⇒ To compare CPU speeds.
- Benchmarks to compare I/O and other subsystems may be included in future releases.

SPEC (Cont)

- The elapsed time to run two copies of a benchmark on each of the N processors of a system (a total of 2N copies) is measured and compared with the time to run two copies of the benchmark on a reference system (which is VAX-11/780 for Release 1.0).
- □ For each benchmark, the ratio of the time on the system under test and the reference system is reported as SPECthruput using a notation of #CPU@Ratio. For example, a system with three CPUs taking 1/15 times as long as the the reference system on GCC benchmark has a SPECthruput of <u>3@15</u>.
- Measure of the per processor throughput relative to the reference system

SPEC (Cont)

- The aggregate throughput for all processors of a multiprocessor system can be obtained by multiplying the ratio by the number of processors. For example, the aggregate throughput for the above system is 45.
- The geometric mean of the SPECthruputs for the 10 benchmarks is used to indicate the overall performance for the suite and is called SPECmark.



- Synthetic workload are representative, repeatable, and avoid sensitive information
- □ Add instruction most frequent instruction initially
- □ Instruction mixes, Kernels, synthetic programs
- □ Application benchmarks: Sieve, Ackerman, ...
- □ Benchmark standards: Debit-Credit, SPEC

Exercise 4.1

Select an area of computer systems (for example, processor design, networks, operating systems, or databases), review articles on performance evaluation in that area and make a list of benchmarks used in those articles.

Exercise 4.2

Implement the Sieve workload in a language of your choice, run it on systems available to you, and report the results.