Multimedia Networking

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Audio/Video recordings of this lecture are available on-line at:

http://www.cse.wustl.edu/~jain/cse473-20/

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http://www.cse.wustl.edu/~jain/cse473-20/i_9mmn.htm

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- Multimedia Networking Applications
- Skype
- Real-Time Transport Protocol (RTP)
- Session Initiation Protocol (SIP)

Note: This class lecture is based on Chapter 9 of the textbook (Kurose and Ross) and the figures provided by the authors.

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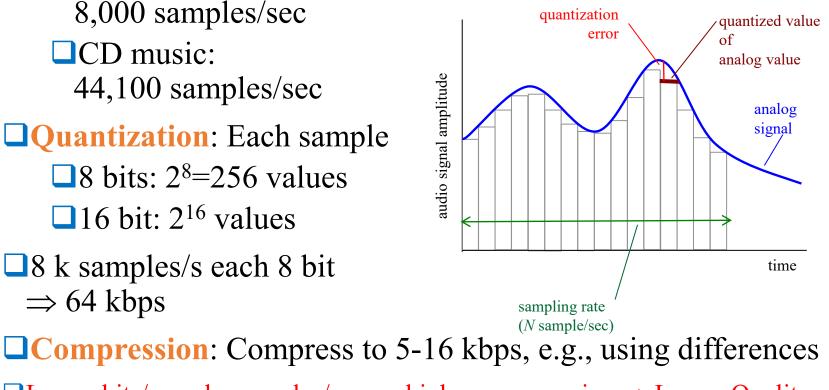


- 1. Audio Digitization
- 2. Playout Buffers
- 3. Streaming Using UDP
- 4. Streaming Using HTTP

Audio Digitization

Sampling: Analog audio signal sampled at constant rate

- **Telephone**: 8,000 samples/sec **CD** music: 44,100 samples/sec **Quantization**: Each sample $\square 8$ bits: $2^8 = 256$ values
 - \square 16 bit: 2¹⁶ values
- ■8 k samples/s each 8 bit \Rightarrow 64 kbps



 \Box Lower bits/sample, samples/sec, or higher compression \Rightarrow Lower Quality

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Multimedia Networking Applications

Streaming Stored Multimedia

Stored Media: Fast rewind, pause, fast forward
 Streaming: simultaneous play out and download

Continuous play out: Delay jitter smoothed by playout buffer

Streaming Live Multimedia: IPTV and Internet Radio

□No fast-forward

- High data rate to large number of users
 - \Rightarrow multicast or P2P,

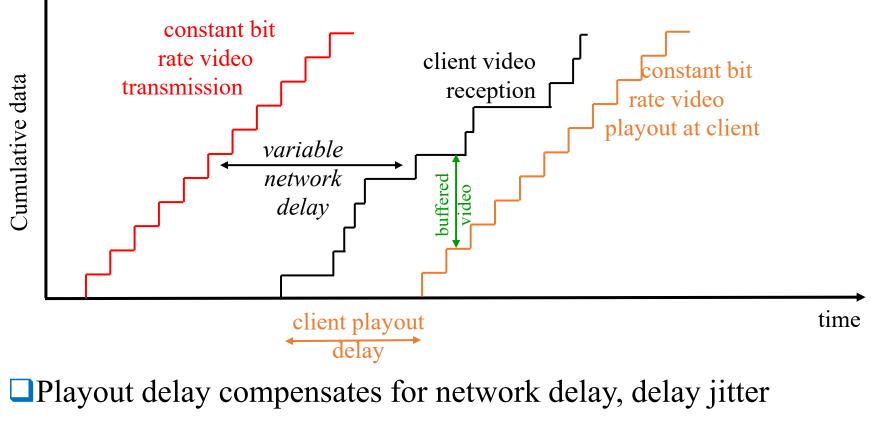
Delay jitter controlled by caching,

Real-Time Interactive Multimedia: Internet Telephone, Video Conferencing

Delay<400 ms.

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Playout Buffers

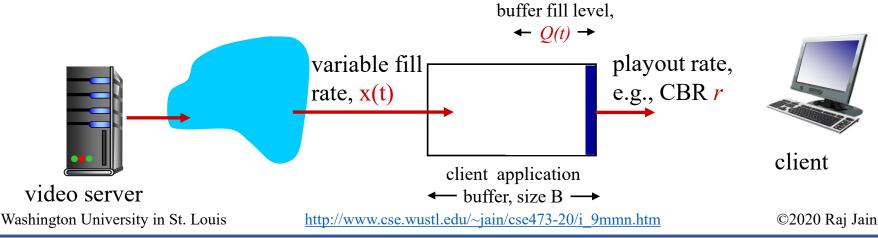


 $\Box Delay > Playout Delay \Rightarrow Packet late \Rightarrow Same as a lost packet$

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Client-side Buffering

- 1. Initial fill of buffer until playout begins at t_p
- 2. Fill rate x(t) varies and playout rate r is constant
- 3. x < r: Buffer eventually empties causing freezing of video
- 4. x > r: buffer will not empty, Flow control to avoid overflow
- 5. *Tradeoff: Large initial playout delay* \Rightarrow *B*uffer starvation less likely but larger delay until user begins watching

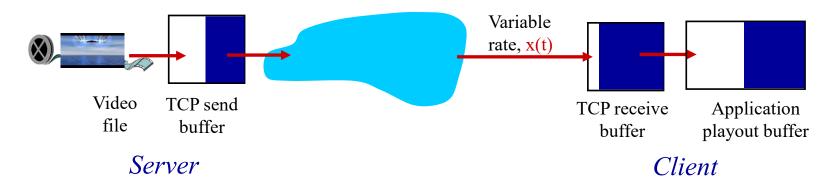


Streaming Using UDP

Server sends at rate appropriate for client
 Often: Send rate = Encoding rate = Constant
 Transmission rate can be oblivious to congestion levels
 Short playout delay (2-5 seconds) to remove network jitter
 Application level error recovery
 UDP may *not* go through firewalls

Streaming Using HTTP

Multimedia file retrieved via HTTP GETSend at maximum possible rate under TCP



□Fill rate fluctuates due to TCP congestion control, retransmissions (in-order delivery)

Larger playout delay to smooth TCP delivery rateHTTP/TCP passes more easily through firewalls

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Review Multimedia Applications

- 1. Audio is sampled, digitized, and compressed
- 2. Initial playout delay helps overcome the jitter in delay
- 3. UDP results in lower jitter but may not go through firewall
- 4. HTTP uses TCP and so the delay variation can be large

Ref: Section 9.1 and 9.2, Review Questions R1-R8, Problems P1-P5Washington University in St. Louishttp://www.cse.wustl.edu/~jain/cse473-20/i_9mmn.htm



- 1. VoIP Packet Losses
- 2. VoIP with Fixed Playout Delay
- 3. VoIP with Adaptive Playout Delay
- 4. Recovering From Packet Loss
- 5. Skype

Voice-over-IP (VoIP)

End-end-delay Requirement: needed to maintain "conversational" aspect

Higher delays noticeable, impair interactivity

- \Box < 150 ms: good
- □> 400 ms: bad
- Includes application-level (packetization, playout), network delays

Alternating talk spurts, silent periods.

□64 kbps during talk spurt

□Packets generated only during talk spurts

□20 ms chunks at 8 kbytes/sec: 160 bytes of data

Application sends a segment every 20 ms during talk spurt

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VoIP Packet Losses

Network Loss: IP datagram lost due to network congestion (router buffer overflow)

Delay Loss: IP datagram arrives too late for playout

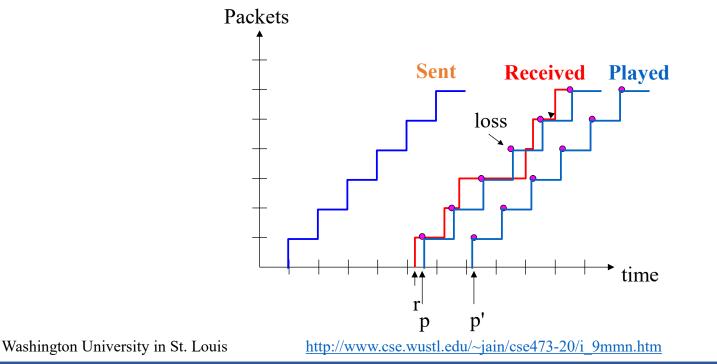
□typical maximum tolerable delay: 400 ms

Loss Tolerance: *P*acket loss rates up to 10% can be concealed

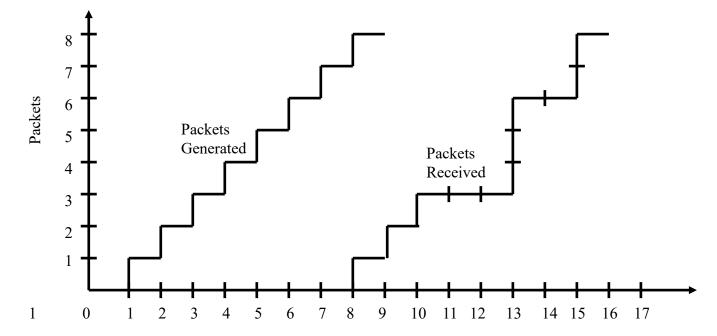
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VoIP with Fixed Playout Delay

- Example: Packets sent every 20 ms during talk spurt.
- First packet received at time r
- □ If playout begins at p, 4th packet will arrive too late
- □ If playout begins at p', all packets can be played on time



Homework 9: Playout Delay



- □ [4 points] Consider the packet generation and reception sequence shown above. The first packet is generated at t=1 and is received at t=8.
- A. If Playout delay is zero and playout begins at t=8, which of the packets will not arrive in time? Show in a table.
- □ B. What is the minimum playout delay at the receiver that result in all of the first eight packets arriving in time for their playout?

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Adaptive Playout Delay

Estimate network delay, adjust playout delay at beginning of each talk spurt

Silent periods compressed and elongated

Chunks still played out every 20 ms during talk spurt

Adaptively estimate packet delay: Similar to TCP RTT estimate

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Adaptive Playout Delay

□ t_i = Sending time □ r_i = Receiving time □ Measured delay sample = r_i - t_i □ d_i = Average network delay $d_i = (1-a)d_{i-1} + a(r_i-t_i)$ □ v_i = Variation of the delay $v_i = (1-a)v_{i-1} + a|r_i - t_i - d_i|$ □ p_i = Playout time

$$p_i = t_i + d_i + K v_i$$

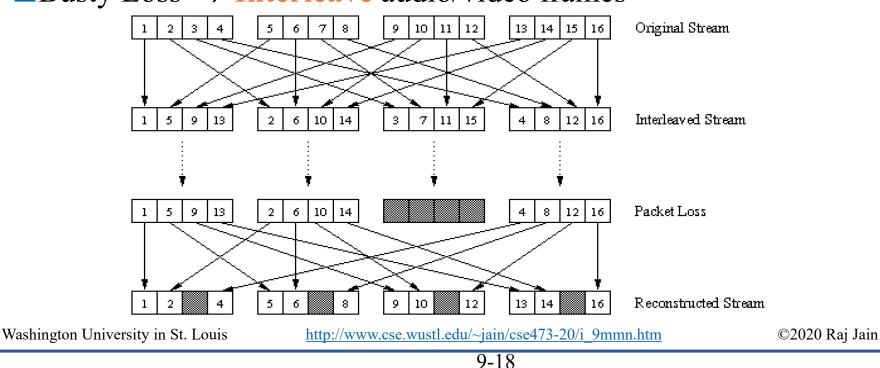
 \Box Here *K* is a constant, say 4.

Sequence numbers and timestamps used to determine talk spurts and silence

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Recovering From Packet LossForward Error Correction

- □Send n+1 packets in place of n packets
- Send a lower-resolution stream in addition
- Play out the old syllable
- $\Box Busty Loss \Longrightarrow Interleave audio/video frames$



Voice-over-IP: Skype

Proprietary

application-layer protocol (inferred via reverse engineering)

- **Encrypted** messages
- **P2P**: Media does not go through a central server
- Clients: Skype peers connect directly to each other for VoIP call
- **Super Nodes** (SN): Skype peers with special functions
- Overlay Network: Among SNs to locate clients

Login server

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Skype clients (SC) Skype login server Super node (SN) Super node overlay network

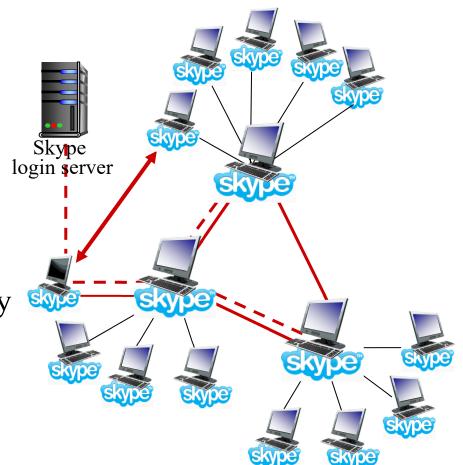
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P2P voice-over-IP: Skype

Skype client operation:

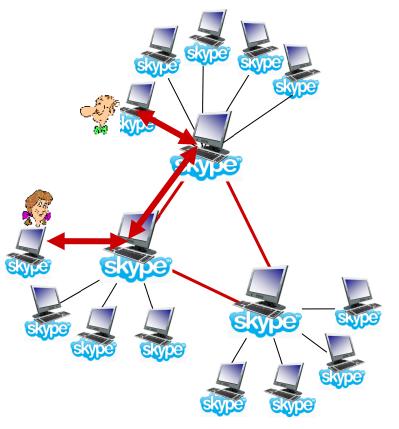
- 1. Joins Skype network by contacting SN (IP address cached) using TCP
- 2. Logs-in: Username, password to centralized Skype login server
- 3. Obtains IP address for callee from SN, SN overlay
- 4. Initiate call directly to callee via SN



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Skype: Super Nodes as Relays

- **Problem**: both Alice, Bob are behind "NATs"
 - NAT prevents outside peer from initiating connection to inside peer
 - □Inside peer *can* initiate connection to outside
- Relay solution: Alice, Bob maintain open connection to their SNs
 - Alice signals her SN to connect to Bob
 - □Alice's SN connects to Bob's SN
 - Bob' s SN connects to Bob over open connection Bob initiated



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- 1. Talk spurts are delayed to overcome jitter
- Playout delay is estimated adaptively using mean and 2. standard deviation
- Forward error correction and interleaving is used to overcome 3 losses and burst errors
- Skype uses super nodes to help connect peers. A login server 4. is used for authentication.
- Skype nodes maintain an outgoing connection with the super 5. nodes. These connections are used for incoming VoIP packets.

Ref: Section 9.3, Review Question R9-R11, Problems P6-P14 Washington University in St. Louis



- 1. Real-Time Transport Protocol (RTP)
- 2. Session Initiation Protocol (SIP)
- 3. H.323 Protocols

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Real-Time Transport Protocol (RTP)

Common sublayer between applications and UDP

Provides sequence numbers, timestamps, and other facilitiesSupports both unicast and multicast

transport layer UDP IP Data Link Physical

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RTP Packet Format

Marker	Payload Type	Sequence Number	Time- stamp	Synchronization Source Identifier	Miscellaneous Fields
1b	7b	16b	32b	32b	

Marker indicates that the packet contains special data required by some applications

SSRC = Synchronization <u>Source</u> = Stream #

Payload	Coding	Rate							
Туре									
0	PCM mu-law	64 kbps							
3	GSM	13 kbps							
7	LPC	2.4 kbps							
26	Motion JPEG								
31	H.261								
33	MPEG2 video								

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Session Initiation Protocol (SIP)

- Application level signaling protocol for voice and video conferencing over Internet
- Allows creating, modifying, terminating sessions with one or more participants
- Carries session descriptions (media types) for user capabilities negotiation
- Supports user location, call setup, call transfers
- Supports mobility by proxying and redirection

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SIP (Cont)

■ SIP Uniform Resource Identifiers (URIs): Similar to email URLs sip:jain@wustl.edu sip:+1-614-292-3989:123@wustl.edu?subject=lecture

□SIP can use UDP or TCP

□SIP messages are sent to SIP servers:

Registrar: Clients register and tell their location to it

Location: Given name, returns possible addresses for a user. Like Directory service or DNS.

Redirect: Returns current address to requesters

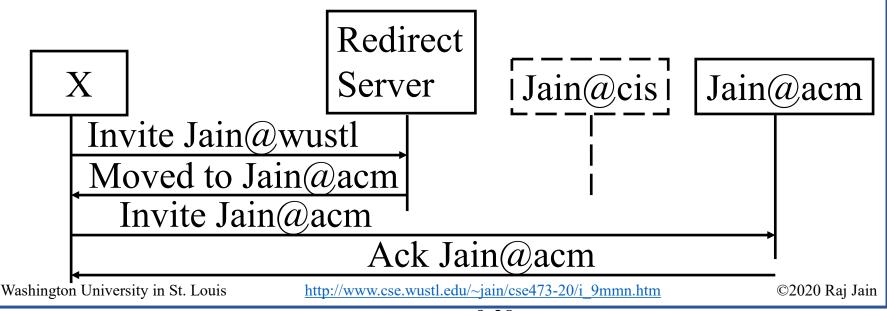
Proxy: Intermediary. Acts like a server to internal client and like a client to external server

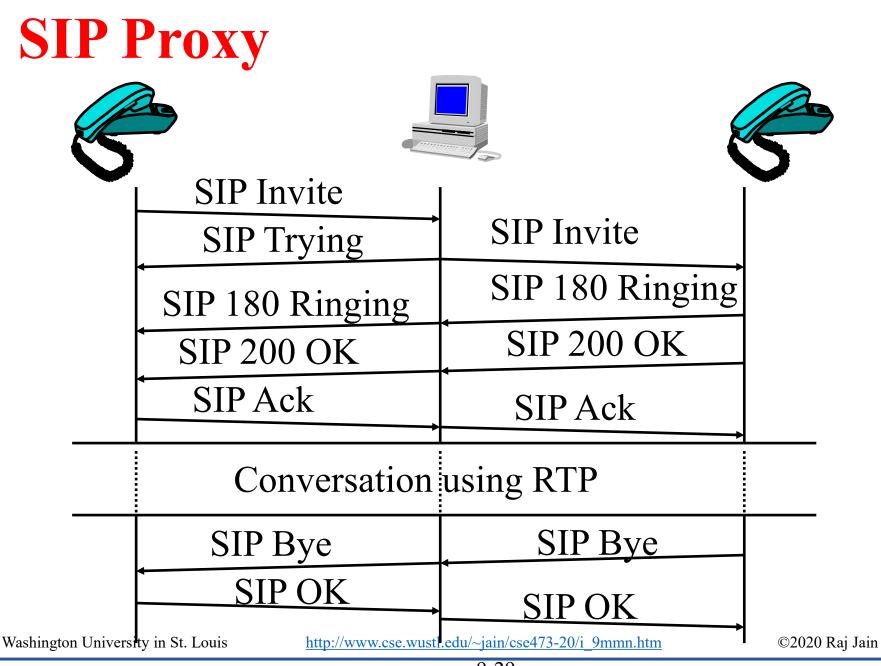
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Locating using SIP

Allows locating a callee at different locations

- Callee registers different locations with Registrar
- SIP Messages: Ack, Bye, Invite, Register, Redirection, ..





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H.323 Protocols

Multimedia over LANs, V1 (June 96), V2(Feb 98)

Provides component descriptions, signaling procedures, call control, system control, audio/video codecs, data protocols

Video Audio	(Control and Management						
H.261 H.263 G.711, G.722 G.723.1, G.72 G.729	2, 28, 22, 28, 20, 20, 20, 20, 20, 20, 20, 20, 20, 20	H.225.0 RAS	H.225.0 Signaling		T.124			
RTP		X.224 Class 0			T.125			
UDP			T.123					
Network (IP)								
Datalink (IEEE 802.3)								

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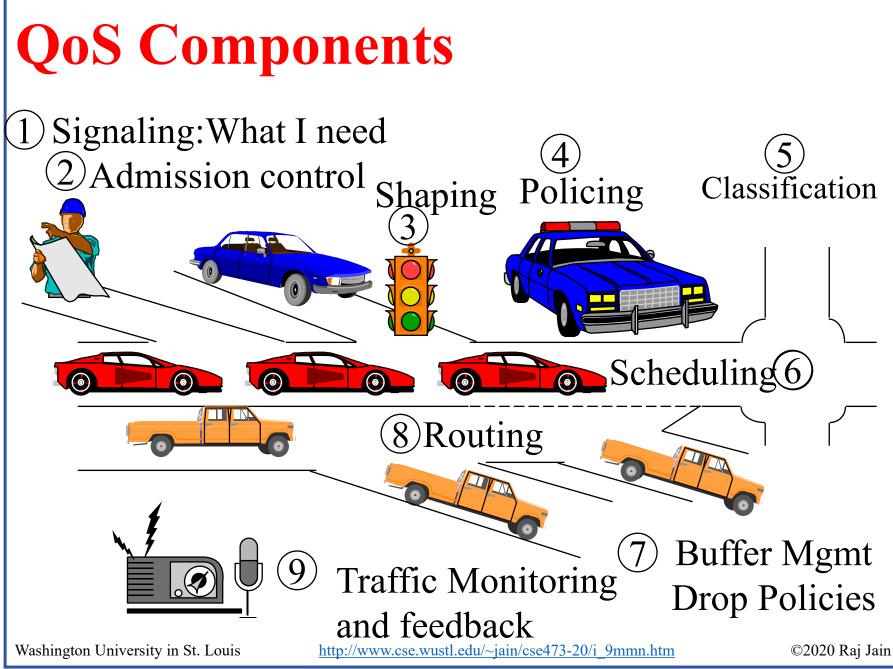
Review Protocols for Real-Time Applications

- 1. RTP is used to transmit multimedia over UDP
- 2. SIP is a signaling (control) protocol to establish multimedia connections
- 3. H.323 is a framework for a group of protocols used for multimedia

Ref: Section 9.4, Review questions R12-R13, Problems P15-P16Washington University in St. Louishttp://www.cse.wustl.edu/~jain/cse473-20/i_9mmn.htm



- 1. QoS Components
- 2. Traffic Shaping
- 3. Token Bucket Shaper
- 4. Traffic Policing
- 5. Differentiated Services



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QoS Components (Cont)

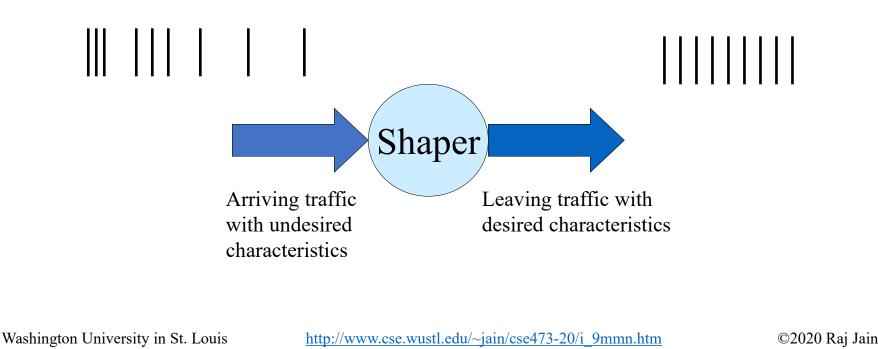
- **1. Signaling**: Users need to tell/negotiate their QoS requirements with the network
- 2. Admission Control: Network can deny requests that it can not meet
- 3. Shaping: Traffic is smoothed out so that it is easier to handle
- **4. Policing**: Ensuring that the users are sending at the rate they agreed to
- 5. Marking/Classification: Packets are classified based on the source, destination, TCP ports (application)
- 6. Scheduling: Different flows get appropriate treatment
- 7. **Drop Policies**: Low priority packets are dropped.
- **8. Routing**: Packets are sent over paths that can meet the QoS
- **9.** Traffic Management: Sources may be asked to reduce their rates to meet the loss rate and delay guarantees

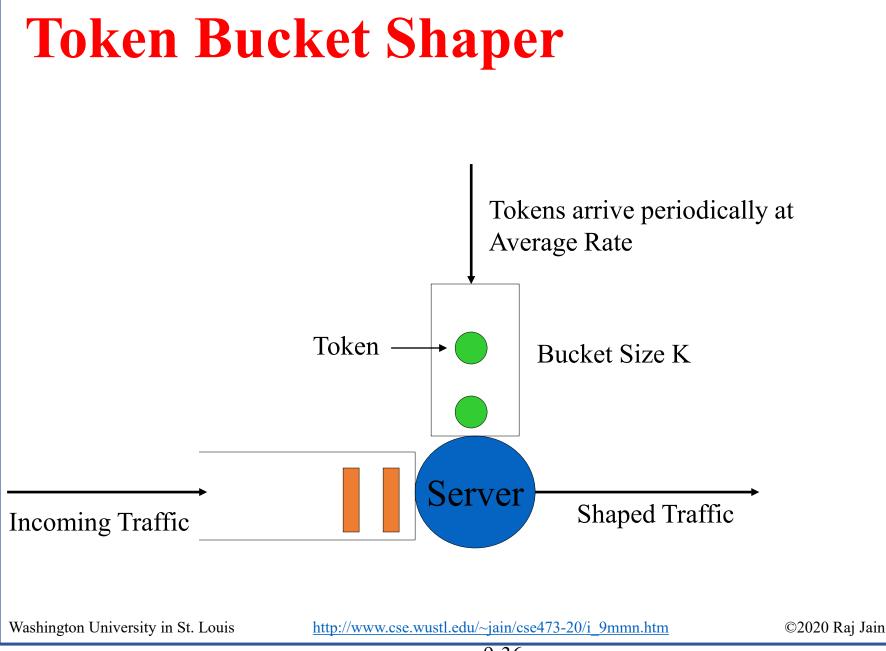
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Traffic Shaping

Altering the traffic characteristics of a given flow is called traffic shaping

The source must shape its traffic prior to sending it to network so it does not violate traffic contract





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Traffic Policing

Users violating the traffic contract can jeopardise the QoS of other connections

- The network must protect well behaving users against such traffic violations
- Policing functions are deployed at the edge (entry) of the network
 Conforming traffic

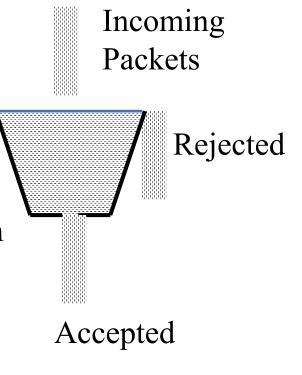
admitted into network

Arriving traffic Non-conforming traffic (dropped)

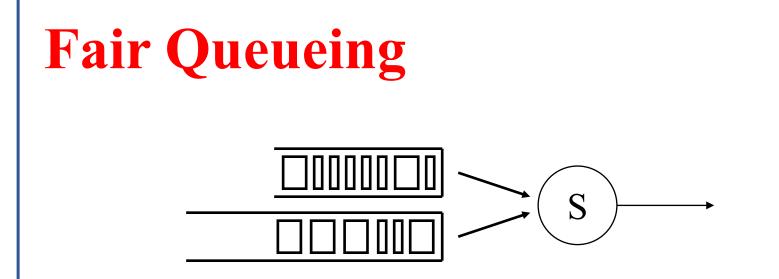
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Peak Rate Policing with Leaky Bucket

- Enforces sustained rate and maximum burst size
- Requires only one counter
 - □ counter is decremented, to a minimum of zero, at the avg rate
 - counter is incremented by one, to a maximum of a limiting value, for each packet arrival
- An arriving packet is non-conforming if counter is at its limit



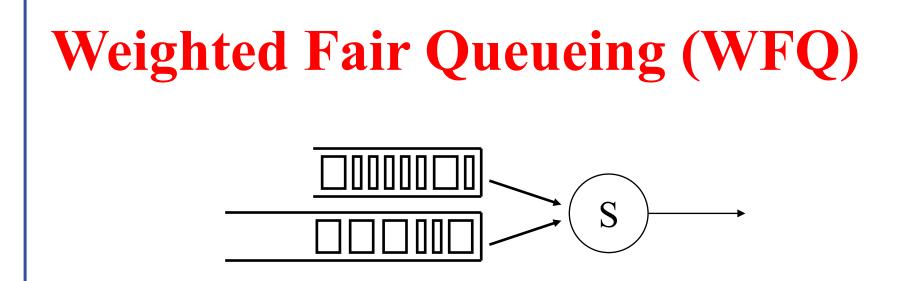
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Bit-level round robin but packet level scheduling

- Count the packet size and determine which packet would finish first. Serve that packet.
- Each flow gets the same number of bits/sec

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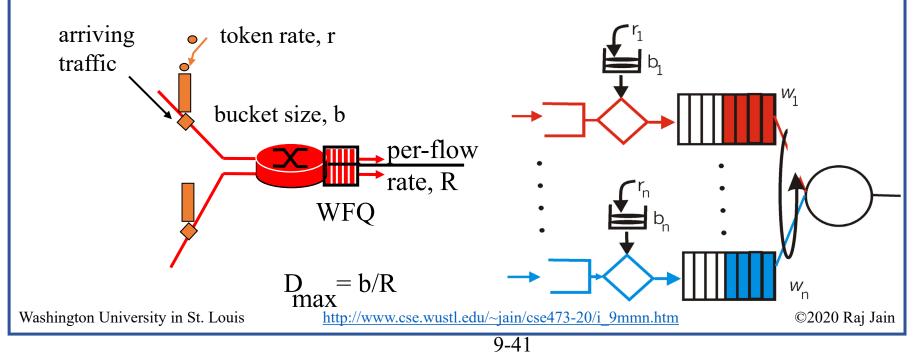


Fair queueing with different weight for each queue
Flow 1 gets x bit/sec
Flow 2 gets y bit/sec
Flow n gets z bit/sec
Here, x, y, z are weights

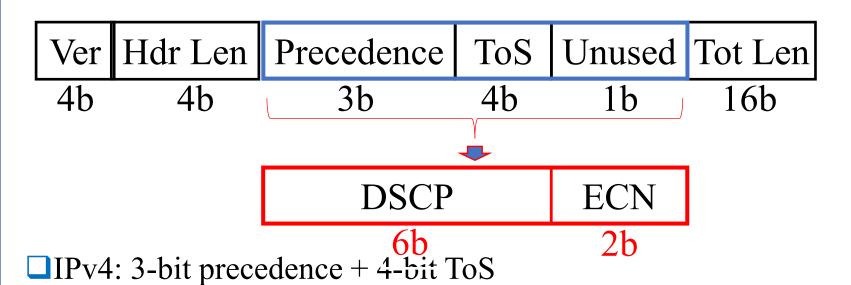
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Maximum Delay with WFQ and Policing

Max Delay d_{max} = b_i/(R w_i/Σ w_j)
 Here,
 b_i=Burst size of ith flow
 R=Service Rate
 W_i=Weight of ith flow



Differentiated Services



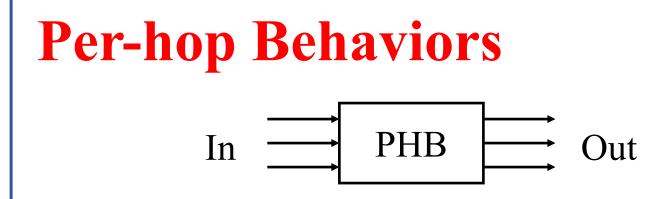
OSPF and integrated IS-IS can compute paths for each ToS

□ Many vendors use IP precedence bits but the service varies \Rightarrow Need a standard \Rightarrow Differentiated Services

 $\Box Edge routers can mark the packets \Rightarrow Set ToS field$

Core routers use ToS field to provide "Per-Hop-Behavior"

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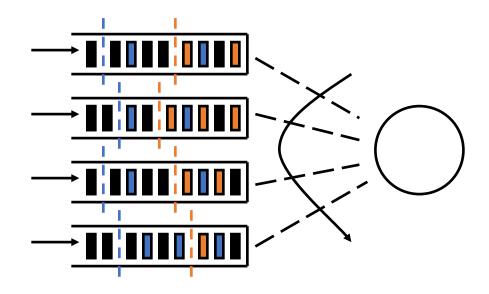
- Externally Observable Forwarding Behavior
- $\Box x\%$ of link bandwidth
- Minimum x% and fair share of excess bandwidth
- Priority relative to other PHBs

Expedited Forwarding

- Also known as "Premium Service"
- □Virtual leased line
- Guaranteed minimum service rate
- Policed: Arrival rate < Minimum Service Rate</p>
- ❑Not affected by other data PHBs
 ⇒ Highest data priority (if priority queueing)
 ❑Code point: 101 110

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Assured Forwarding



PHB <u>Group</u> Four Classes: No particular ordering Three drop preference per class

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Assured Forwarding (Cont)

- DS nodes SHOULD implement all 4 classes and MUST accept all 3 drop preferences. Can implement 2 drop preferences.
- Similar to nrt-VBR/ABR/GFR
- Code Points:

Drop Prec.	Class 1	Class 2	Class 3	Class 4
Low	010 000	011 000	100 000	101 000
Medium	010 010	011 010	100 010	101 010
High	010 100	011 100	100 100	101 100

□ Avoids 11x000 (used for network control)

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Network Support for Multimedia

Approach	Granularity	Guarantee	Mechanisms	Complex	Deployed?
Making best	All traffic	None or	No network	low	everywhere
of best effort	treated	soft	support (all at		
service	equally		application)		
Differentiated	Traffic	None of	Packet market,	med	some
service	"class"	soft	scheduling,		
			policing.		
Per-	Per-	Soft or hard	Packet market,	high	little to
connection	connection	after flow	scheduling,		none
QoS	flow	admitted	policing, call		
			admission		

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Review Network Support for Multimedia

- 1. QoS is obtained using several components including shaping, policing, differentiated services
- 2. Shaping is done by a token bucket
- 3. Policing is done using a leaky bucket
- 4. Differentiated services specifies per-hop behaviors
 - 1. Expedited Forwarding: min service rate
 - 2. Assured Forwarding: 4 classes, 3 drop precedence's

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- 1. Multimedia applications require bounded delay, delay jitter, and minimum throughput
- 2. RTP allows sequencing and time stamping
- 3. SIP allows parameter negotiation and location
- 4. QoS requires shaping, policing, scheduling, etc.
- 5. Diffserv allows different packets to get different service

Ref: Entire Chapter 9, Review Exercises R1-R12, Problems P2-P4,P9, P11, P16, P19Washington University in St. Louishttp://www.cse.wustl.edu/~jain/cse473-20/i_9mmn.htm

Acronyms

ABRAvailable Bit Rate

- CBRConstant Bit Rate
- CD Compact Disk
- DNS Domain Name System
 - DiffServe
- GFR Guaranteed Frame Rate
- HTTP HyperText Transfer Protocol
- **IEEE** Institution of Electrical and Electronics Engineers
 - Internet Protocol
- □IPTV Internet Protocol Television
- IPv4Internet Protocol Version 4
- □IS Integrated Services
- LAN Local Area Network
- ■NAT Network Address Translator
 - Open Shortest Path First
 - Per-Hop Behavior

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DS

DPHB

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Acronyms (Cont)

QoS Quality of Service Registration, Admission, and Status \Box RTCP Real-Time Transport Protocol Control Protocol **Real-Time Transport Protocol Real-Time Streaming Protocol** \Box RTSP **Round Trip Time Skype Clients Session Initiation Protocol** Super Node Synchronization Source **Transmission Control Protocol** Type of Service User Datagram Protocol Uniform Resource Identifiers Uniform Resource Locator $\Box VBR$ Variable Bit Rate

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Acronyms (Cont)

□VoIP Voice over IP

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Related Modules

CSE 567: The Art of Computer Systems Performance Analysis <u>https://www.youtube.com/playlist?list=PLjGG94etKypJEKjNAa1n_1X0bWWNyZcof</u>

CSE473S: Introduction to Computer Networks (Fall 2011), https://www.youtube.com/playlist?list=PLjGG94etKypJWOSPMh8Azcgy5e_10TiDw





CSE 570: Recent Advances in Networking (Spring 2013)

https://www.youtube.com/playlist?list=PLjGG94etKypLHyBN8mOgwJLHD2FFIMGq5

CSE571S: Network Security (Spring 2011),

https://www.youtube.com/playlist?list=PLjGG94etKypKvzfVtutHcPFJXumyyg93u





Video Podcasts of Prof. Raj Jain's Lectures,

https://www.youtube.com/channel/UCN4-5wzNP9-ruOzQMs-8NUw

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