# Multimedia Networking

#### Raj Jain

Washington University in Saint Louis Saint Louis, MO 63130 Jain@wustl.edu

Audio/Video recordings of this lecture are available on-line at:

http://www.cse.wustl.edu/~jain/cse473-19/

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/



- Multimedia Networking Applications
- Skype
- Real-Time Transport Protocol (RTP)
- Session Initiation Protocol (SIP)

**Note**: This class lecture is based on Chapter 9 of the textbook (Kurose and Ross) and the figures provided by the authors.

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

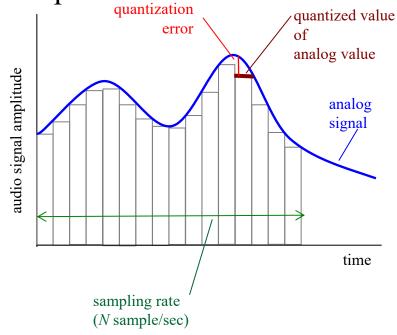


# Overview Multimedia Applications

- 1. Audio Digitization
- 2. Playout Buffers
- 3. Streaming Using UDP
- 4. Streaming Using HTTP

# **Audio Digitization**

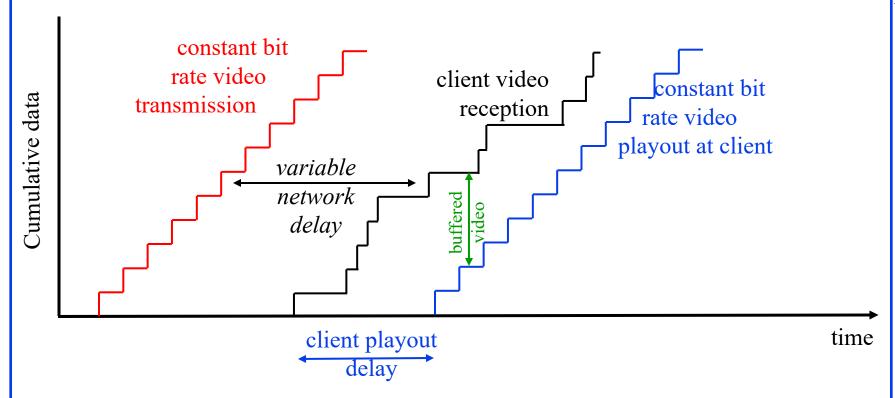
- □ **Sampling**: Analog audio signal sampled at constant rate
  - Telephone:8,000 samples/sec
  - CD music:44,100 samples/sec
- Quantization: Each sample
  - > 8 bits:  $2^8$ =256 values
  - > 16 bit: 2<sup>16</sup> values
- 8 k samples/s each 8 bit
  - $\Rightarrow$  64 kbps
- □ Compression: Compress to 5-16 kbps, e.g., using differences



#### **Multimedia Networking Applications**

- **□** Streaming Stored Multimedia
  - > Stored Media: Fast rewind, pause, fast forward
  - > Streaming: simultaneous play out and download
  - Continuous play out: Delay jitter smoothed by playout buffer
- □ Streaming Live Multimedia: IPTV and Internet Radio
  - > No fast-forward
  - > High data rate to large number of users
    - $\Rightarrow$  multicast or P2P,
  - > Delay jitter controlled by caching,
- □ Real-Time Interactive Multimedia: Internet Telephone, Video Conferencing
  - > Delay<400 ms.

#### **Playout Buffers**



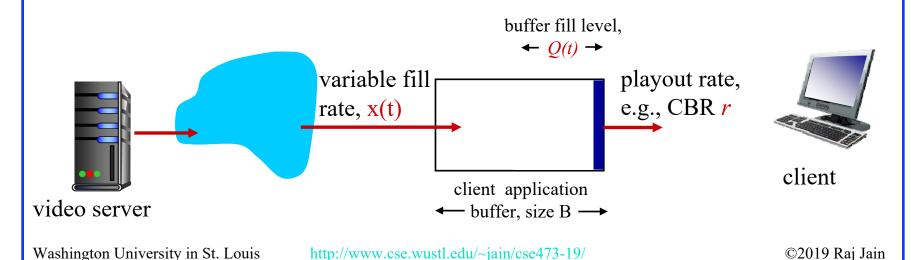
- □ Playout delay compensates for network delay, delay jitter
- □ Delay > Playout Delay  $\Rightarrow$  Packet late  $\Rightarrow$  Same as a lost packet

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

## **Client-side Buffering**

- 1. Initial fill of buffer until playout begins at  $t_p$
- 2. Fill rate x(t) varies and playout rate r is constant
- 3. x < r: Buffer eventually empties causing freezing of video
- 4. x > r: buffer will not empty, Flow control to avoid overflow
- 5. Tradeoff: Large initial playout delay  $\Rightarrow$  Buffer starvation less likely but Larger delay until user begins watching

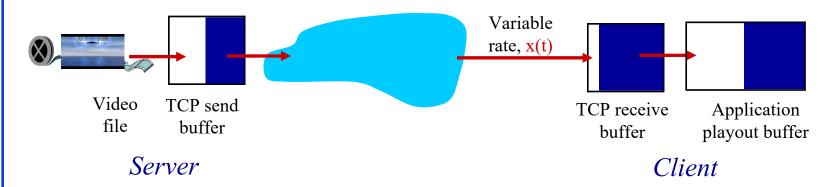


# **Streaming Using UDP**

- Server sends at rate appropriate for client
  - > Often: Send rate = Encoding rate = Constant
  - > Transmission rate can be oblivious to congestion levels
- □ Short playout delay (2-5 seconds) to remove network jitter
- Application level error recovery
- □ UDP may *not* go through firewalls

# **Streaming Using HTTP**

- Multimedia file retrieved via HTTP GET
- Send at maximum possible rate under TCP



- □ Fill rate fluctuates due to TCP congestion control, retransmissions (in-order delivery)
- □ Larger playout delay to smooth TCP delivery rate
- HTTP/TCP passes more easily through firewalls

#### Review



#### Multimedia Applications

- 1. Audio is sampled, digitized, and compressed
- 2. Initial playout delay helps overcome the jitter in delay
- 3. UDP results in lower jitter but may not go through firewall
- 4. HTTP uses TCP and so the delay variation can be large



- VoIP Packet Losses
- 2. VoIP with Fixed Playout Delay
- 3. VoIP with Adaptive Playout Delay
- 4. Recovering From Packet Loss
- 5. Skype

#### Voice-over-IP (VoIP)

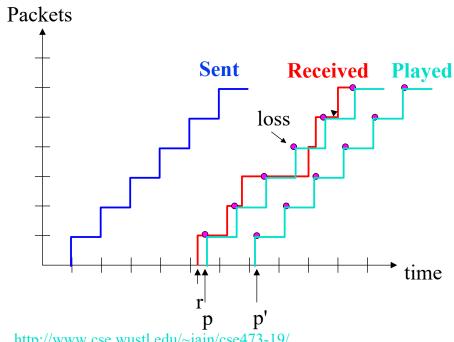
- End-end-delay Requirement: needed to maintain "conversational" aspect
  - > Higher delays noticeable, impair interactivity
  - > < 150 ms: good
  - > 400 ms: bad
  - Includes application-level (packetization, playout), network delays
- □ Alternating talk spurts, silent periods.
  - > 64 kbps during talk spurt
  - > Packets generated only during talk spurts
  - > 20 ms chunks at 8 Kbytes/sec: 160 bytes of data
- Application sends a segment every 20 ms during talk spurt

#### **VoIP Packet Losses**

- Network Loss: IP datagram lost due to network congestion (router buffer overflow)
- □ Delay Loss: IP datagram arrives too late for playout
  - > typical maximum tolerable delay: 400 ms
- Loss Tolerance: Packet loss rates between 1% and 10% can be concealed

#### **VoIP** with Fixed Playout Delay

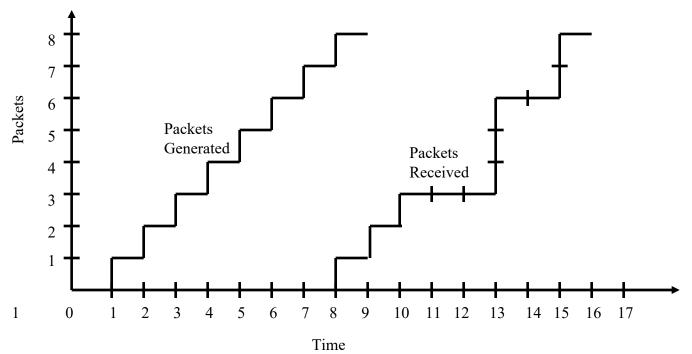
- Example: Packets sent every 20 ms during talk spurt.
- First packet received at time r
- If playout begins at p, 4th packet will arrive too late
- If playout begins at p', all packets can be played on time



Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

#### Homework 9



- Consider the packet generation and reception sequence shown below. The first packet is generated at t=1 and is received at t=8.
- A. If Playout delay is zero and playout begins at t=8, which of the packets will not arrive in time?
- B. What is the minimum playout delay at the receiver that result in all of the first eight packets arriving in time for their playout? Washington University in St. Louis <a href="http://www.cse.wustl.edu/~jain/cse473-19/">http://www.cse.wustl.edu/~jain/cse473-19/</a>

## **Adaptive Playout Delay**

- Estimate network delay, adjust playout delay at beginning of each talk spurt
- Silent periods compressed and elongated
- □ Chunks still played out every 20 ms during talk spurt
- Adaptively estimate packet delay: Similar to TCP RTT estimate

# **Adaptive Playout Delay**

- $\Box$   $t_i$ =Sending time
- $ightharpoonup r_i$ = Receiving time
- Measured delay sample =  $r_i$ - $t_i$
- $\Box$   $d_i$ = Average network delay

$$d_i = (1-a)d_{i-1} + a(r_i - t_i)$$

 $\mathbf{v}_i$ = Variation of the delay

$$v_i = (1-a)v_{i-1} + a|r_i - t_i - d_i|$$

 $\Box$   $p_i$ = Playout time

$$p_i = t_i + d_i + Kv_i$$

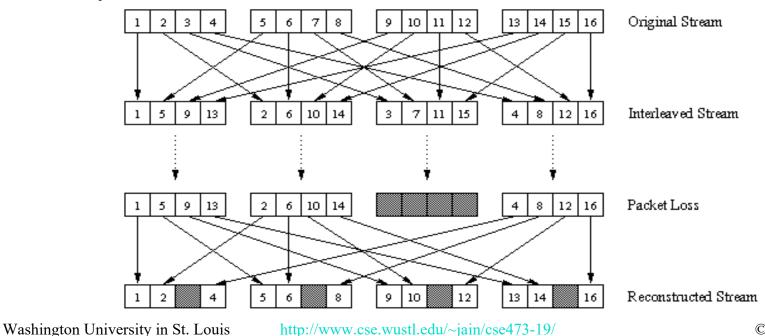
- $\square$  Here K is a constant, say 4.
- Sequence numbers and timestamps used to determine talk

spurts and silence

http://www.cse.wustl.edu/~jain/cse473-19/

#### **Recovering From Packet Loss**

- Forward Error Correction
- Send n+1 packets in place of n packets
- Send a lower-resolution stream in addition
- Play out the old syllable
- Busty Loss ⇒ Interleave audio/video frames

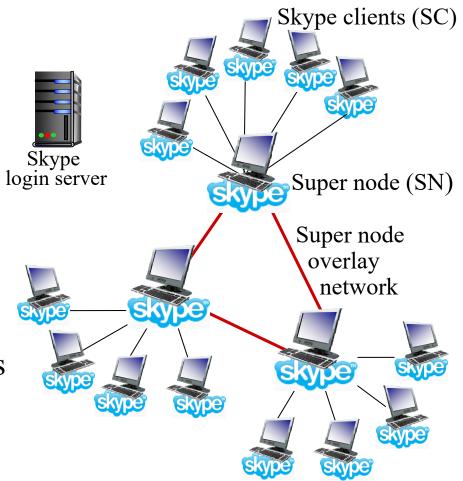


http://www.cse.wustl.edu/~jain/cse473-19/

# Voice-over-IP: Skype

- Proprietary

   application-layer protocol
   (inferred via reverse engineering)
- **Encrypted** messages
- **P2P**: Media does not go through a central server
- Clients: Skype peers connect directly to each other for VoIP call
- Super Nodes (SN): Skype peers with special functions
- Overlay Network: Among SNs to locate clients
- □ Login server



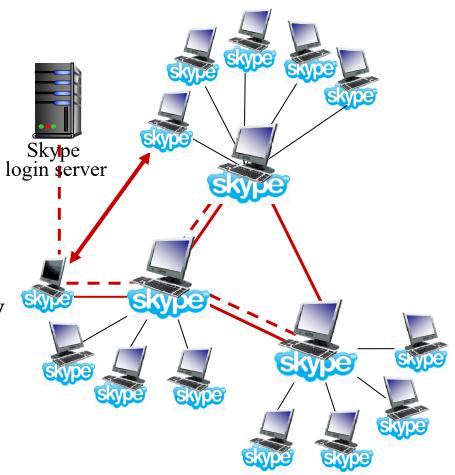
Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

## P2P voice-over-IP: Skype

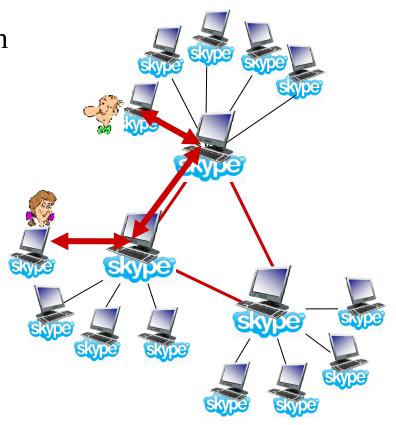
#### Skype client operation:

- 1. Joins Skype network by contacting SN (IP address cached) using TCP
- Logs-in: Username, password to centralized Skype login server
- 3. Obtains IP address for callee from SN, SN overlay
- 4. Initiate call directly to callee via SN



#### Skype: Super Nodes as Relays

- **Problem**: both Alice, Bob are behind "NATs"
  - NAT prevents outside peer from initiating connection to inside peer
  - > Inside peer *can* initiate connection to outside
- Relay solution: Alice, Bob maintain open connection to their SNs
  - Alice signals her SN to connect to Bob
  - Alice's SN connects to Bob's SN
  - Bob's SN connects to Bob over open connection Bob initiated



Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

#### Review



#### Voice over IP

- 1. Talk spurts are delayed to overcome jitter
- 2. Playout delay is estimated adaptively using mean and standard deviation
- 3. Forward error correction and interleaving is used to overcome losses and burst errors
- 4. Skype uses super nodes to help connect peers. A login server is used for authentication.
- 5. Skype nodes maintain an outgoing connection with the super nodes. These connections are used for incoming VoIP packets.

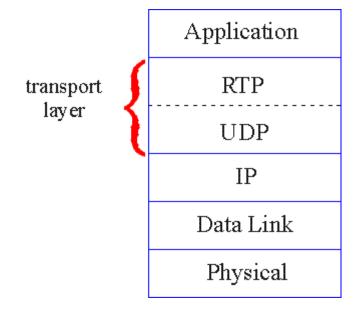


# Overview | Protocols for Real-Time Applications

- 1. Real-Time Transport Protocol (RTP)
- 2. Session Initiation Protocol (SIP)
- 3. H.323 Protocols

#### **Real-Time Transport Protocol (RTP)**

- Common sublayer between applications and UDP
- □ Provides sequence numbers, timestamps, and other facilities
- Supports both unicast and multicast



#### **RTP Packet Format**

Paylcad	Sequence	Timestamp	Syncrhron zation	Miscellaneous
Type	Number		Source Identifer	Fields
8b	16b	32h	32h	

□ SSRC = Synchronization Source = Stream #

Payload	Coding	Rate	
<b>Type</b>			
0	PCM mu-law	64 kbps	
3	GSM	13 kbps	
7	LPC	2.4 kbps	
26	Motion JPEG		
31	H.261		
33	MPEG2 video		

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

#### **Session Initiation Protocol (SIP)**

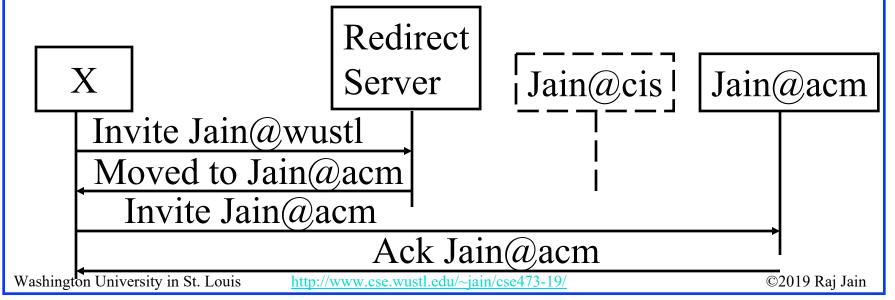
- Application level signaling protocol for voice and video conferencing over Internet
- □ Allows creating, modifying, terminating sessions with one or more participants
- □ Carries session descriptions (media types) for user capabilities negotiation
- □ Supports user location, call setup, call transfers
- Supports mobility by proxying and redirection

#### SIP (Cont)

- □ SIP Uniform Resource Identifiers (URIs):
  Similar to email URLs
  sip:jain@wustl.edu
  sip:+1-614-292-3989:123@wustl.edu?subject=lecture
- □ SIP can use UDP or TCP
- □ SIP messages are sent to SIP servers:
  - > Registrar: Clients register and tell their location to it
  - Location: Given name, returns possible addresses for a user. Like Directory service or DNS.
  - > Redirect: Returns current address to requesters
  - > Proxy: Intermediary. Acts like a server to internal client and like a client to external server

#### Locating using SIP

- □ Allows locating a callee at different locations
- □ Callee registers different locations with Registrar
- □ SIP Messages: Ack, Bye, Invite, Register, Redirection, .



# **SIP Proxy**







SIP Invite	l l			
SIP Trying	SIP Invite			
SIP 180 Ringing	SIP 180 Ringing			
SIP 200 OK	SIP 200 OK			
SIP Ack	SIP Ack			
Conversation using RTP				
SIP Bye	SIP Bye			
SIP OK	SIP OK			

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

#### **H.323 Protocols**

- □ Multimedia over LANs, V1 (June 96), V2(Feb 98)
- Provides component descriptions, signaling procedures, call control, system control, audio/video codecs, data protocols

Video	Audio	Control and Management		Data		
H.261 H.263	G.711, G.722, G.723.1, G.728, G.729	RTCP	H.225.0 RAS	H.225.0 Signaling	H.245 Control	T.124
RTP		X.224 Class 0			T.125	
UDP			TCP			T.123
Network (IP)						1.123
Datalink (IEEE 802.3)						

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

#### Review



#### **Protocols for Real-Time Applications**

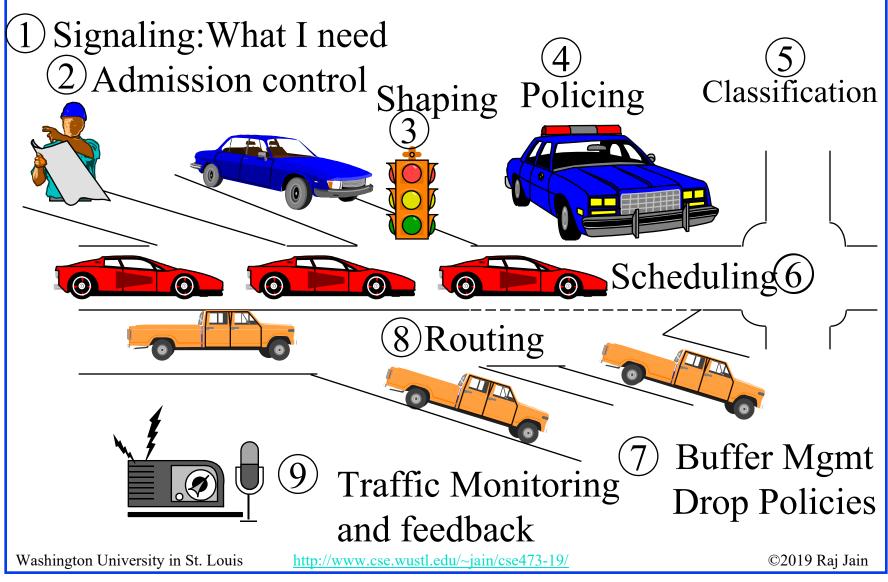
- 1. RTP is used to transmit multimedia over UDP
- 2. SIP is a signaling (control) protocol to establish multimedia connections
- 3. H.323 is a framework for a group of protocols used for multimedia



# **Networking Support for Multimedia**

- 1. QoS Components
- 2. Traffic Shaping
- 3. Token Bucket Shaper
- 4. Traffic Policing
- 5. Differentiated Services

# **QoS Components**



#### **QoS Components (Cont)**

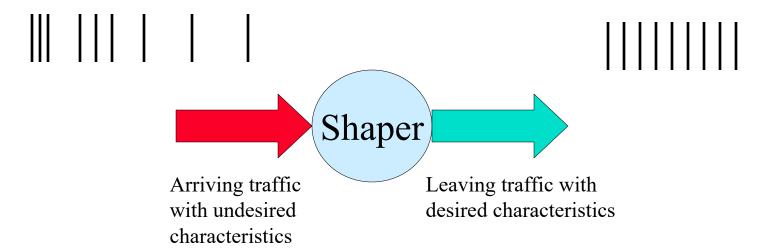
- 1. Signaling: Users need to tell/negotiate their QoS requirements with the network
- 2. Admission Control: Network can deny requests that it can not meet
- 3. Shaping: Traffic is smoothed out so that it is easier to handle
- 4. Policing: Ensuring that the users are sending at the rate they agreed to
- 5. Marking/Classification: Packets are classified based on the source, destination, TCP ports (application)
- 6. Scheduling: Different flows get appropriate treatment
- 7. Drop Policies: Low priority packets are dropped.
- 8. Routing: Packets are sent over paths that can meet the QoS
- 9. Traffic Management: Sources may be asked to reduce their rates to meet the loss rate and delay guarantees

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

# **Traffic Shaping**

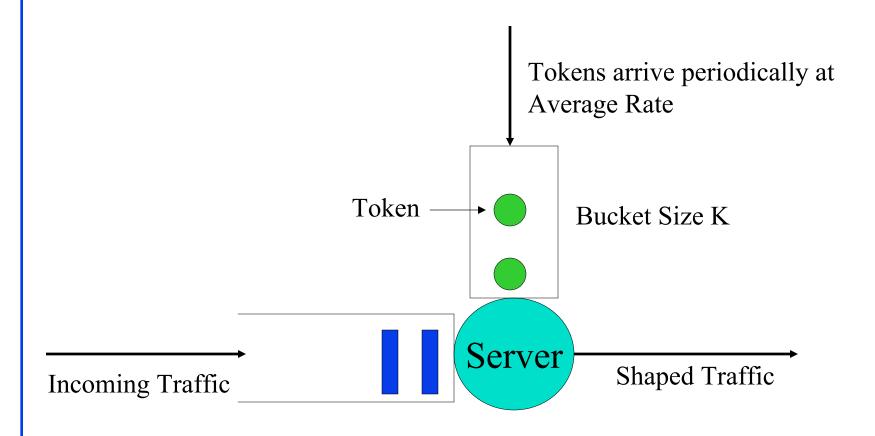
- Altering the traffic characteristics of a given flow is called traffic shaping
- □ The source must shape its traffic prior to sending it to network so it does not violate traffic contract



Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

## **Token Bucket Shaper**

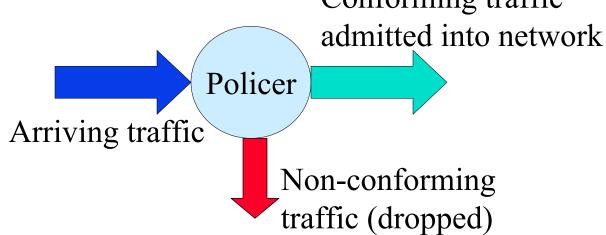


Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

# **Traffic Policing**

- ☐ Users violating the traffic contract can jeopardise the QoS of other connections
- ☐ The network must protect well behaving users against such traffic violations
- Policing functions are deployed at the edge (entry) of the network
   Conforming traffic



Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

# Peak Rate Policing with Leaky Bucket

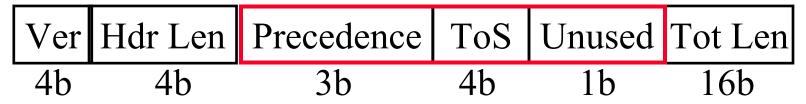
- Enforces sustained rate and maximum burst size
- □ Requires only one counter
  - > counter is decremented, to a minimum of zero, at the avg rate
  - counter is incremented by one,
     to a maximum of a limiting
     value, for each packet arrival
- □ An arriving packet is nonconforming if counter is at its limit

Incoming Packets



Accepted

### **Differentiated Services**



- □ IPv4: 3-bit precedence + 4-bit ToS
- □ OSPF and integrated IS-IS can compute paths for each ToS
- Many vendors use IP precedence bits but the service varies ⇒ Need a standard ⇒ Differentiated Services
- $\square$  Edge routers can mark the packets  $\Rightarrow$  Set ToS field
- ☐ Core routers use ToS field to provide "Per-Hop-Behavior"

Washington University in St. Louis

## **Per-hop Behaviors**

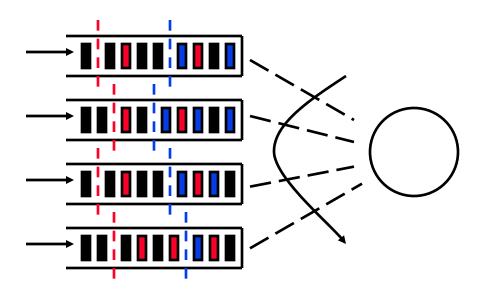


- Externally Observable Forwarding Behavior
- □ x% of link bandwidth
- ☐ Minimum x% and fair share of excess bandwidth
- □ Priority relative to other PHBs

# **Expedited Forwarding**

- □ Also known as "Premium Service"
- □ Virtual leased line
- □ Guaranteed minimum service rate
- □ Policed: Arrival rate < Minimum Service Rate
- □ Not affected by other data PHBs
  - ⇒ Highest data priority (if priority queueing)
- □ Code point: 101 110

# **Assured Forwarding**



- □ PHB Group
- □ Four Classes: No particular ordering
- ☐ Three drop preference per class

## **Assured Forwarding (Cont)**

- □ DS nodes SHOULD implement all 4 classes and MUST accept all 3 drop preferences. Can implement 2 drop preferences.
- □ Similar to nrt-VBR/ABR/GFR
- Code Points:

Drop Prec.	Class 1	Class 2	Class 3	Class 4
Low	010 000	011 000	100 000	101 000
Medium	010 010	011 010	100 010	101 010
High	010 100	011 100	100 100	101 100

□ Avoids 11x000 (used for network control)

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

### **Network Support for Multimedia**

Approach	Granularity	Guarantee	Mechanisms	Complex	Deployed?
Making best	All traffic	None or	No network	low	everywhere
of best effort	treated	soft	support (all at		
service	equally		application)		
Differentiated	Traffic	None of	Packet market,	med	some
service	"class"	soft	scheduling,		
			policing.		
Per-	Per-	Soft or hard	Packet market,	high	little to
connection	connection	after flow	scheduling,		none
QoS	flow	admitted	policing, call		
			admission		

#### Review



### **Network Support for Multimedia**

- 1. QoS is obtained using several components including shaping, policing, differentiated services
- 2. Shaping is done by a token bucket
- 3. Policing is done using a leaky bucket
- 4. Differentiated services specifies per-hop behaviors
  - 1. Expedited Forwarding: min service rate
  - 2. Assured Forwarding: 4 classes, 3 drop precedence's

## Summary



- 1. Multimedia applications require bounded delay, delay jitter, and minimum throughput
- 2. RTP allows sequencing and time stamping
- 3. SIP allows parameter negotiation and location
- 4. QoS requires shaping, policing, scheduling, etc.
- 5. Diffserv allows different packets to get different service

### Acronyms

□ ABR Available Bit Rate

CBR Constant Bit Rate

CD Compact Disk

DNS Domain Name System

□ DS DiffServe

□ GFR Guaranteed Frame Rate

□ HTTP HyperText Transfer Protocol

□ IEEE Institution of Electrical and Electronics Engineers

□ IP Internet Protocol

□ IPTV Internet Protocol Television

□ IPv4 Internet Protocol Version 4

□ IS Integrated Services

□ LAN Local Area Network

□ NAT Network Address Translator

OSPF Open Shortest Path First

PHB Per-Hop Behavior

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

## Acronyms (Cont)

QoS Quality of Service

□ RAS Registration, Admission, and Status

■ RTCP Real-Time Transport Protocol Control Protocol

RTP Real-Time Transport Protocol

□ RTSP Real-Time Streaming Protocol

□ RTT Round Trip Time

□ SC Skype Clients

□ SIP Session Initiation Protocol

□ SN Super Node

□ SSRC Synchronization Source

□ TCP Transmission Control Protocol

□ ToS Type of Service

UDP User Datagram Protocol

□ URI Uniform Resource Identifiers

URL Uniform Resource Locator

□ VBR Variable Bit Rate

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

# Acronyms (Cont)

□ VoIP Voice over IP

### Scan This to Download These Slides





Raj Jain <a href="http://rajjain.com">http://rajjain.com</a>

http://www.cse.wustl.edu/~jain/cse473-19/i\_9mmn.htm

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/

### **Related Modules**



CSE 567: The Art of Computer Systems Performance Analysis

https://www.youtube.com/playlist?list=PLjGG94etKypJEKjNAa1n\_1X0bWWNyZcof

CSE473S: Introduction to Computer Networks (Fall 2011),

https://www.youtube.com/playlist?list=PLjGG94etKypJWOSPMh8Azcgy5e\_10TiDw





CSE 570: Recent Advances in Networking (Spring 2013)

https://www.youtube.com/playlist?list=PLjGG94etKypLHyBN8mOgwJLHD2FFIMGq5

CSE571S: Network Security (Spring 2011),

https://www.youtube.com/playlist?list=PLjGG94etKypKvzfVtutHcPFJXumyyg93u





Video Podcasts of Prof. Raj Jain's Lectures,

 $\underline{https://www.youtube.com/channel/UCN4-5wzNP9-ruOzQMs-8NUw}$ 

Washington University in St. Louis

http://www.cse.wustl.edu/~jain/cse473-19/