Computer Networks and the Internet

Raj Jain

Washington University in Saint Louis Saint Louis, MO 63130 Jain@wustl.edu

Audio/Video recordings of this lecture are available on-line at:

http://www.cse.wustl.edu/~jain/cse473-11/

Washington University in St. Louis



- 1. Physical Media
- 2. Switching: Circuit vs. Packet
- 3. Internet:Edge, Core
- 4. Network Performance Measures: Delay, Loss, Throughput
- 5. Protocol Layers
- 6. Network Security
- 7. History

Note: This class lecture is based on Chapter 1 of the textbook (Kurose and Ross) and the slides provided by the authors.

What is a Network?

Network: Enables data transfer among nodes

- □ Generally heterogeneous nodes
- □ More than 2 nodes
- □ E.g., Your home or office network



□ Communication: Two nodes.

□ Link level electrical issues.

1 - 3



- End Systems: Systems that are sinks or sources of data, e.g., Desktops, Laptops, Servers, Printers, Cell Phones, etc.
- □ Intermediate Systems: Systems that forward/switch data from one link to another, e.g., routers, switches
- **Hosts**: End Systems
- **Gateways**: Routers
- Servers: End Systems that provide service, e.g., print server, storage server, Mail server, etc.
- **Clients**: End systems that request service
- Links: Connect the systems.
 Characterized by transmission rate, propagation delay

Transmission Media

Guided:

- □ Twisted Pair
- Coaxial cable
- □ Optical fiber

Unguided:

- □ Microwave
- □ Satellite
- □ Wireless

Electromagnetic Spectrum



Twisted Pair (TP)

twist length

- -Separately insulated
- —Twisted together
- -Often "bundled" into cables
- Usually installed in building during construction



- **Twists decrease the cross-talk**
- Neighboring pairs have different twist length
- Most of telephone and network wiring in homes and offices is TP.

Shielded and Unshielded TP

- □ Shielded Twisted Pair (STP)
 - □ Metal braid or sheathing that reduces interference
 - □ More expensive
 - □ Harder to handle (thick, heavy)
 - □ Used in token rings
- Unshielded Twisted Pair (UTP)
 - □ Ordinary telephone wire
 - □ Cheap, Flexible
 - \Rightarrow Easiest to install
 - □ No shielding
 - \Rightarrow Suffers from external interference
 - □ Used in Telephone and Ethernet



1-8

UTP Categories

Cat 3

- □ Up to 16MHz
- □ Voice grade found in most offices
- □ Twist length of 7.5 cm to 10 cm

Cat 4

□ Up to 20 MHz. Not used much in practice.

Cat 5

- □ Up to 100MHz
- □ Used in 10 Mbps and 100 Mbps Ethernet
- □ Twist length 0.6 cm to 0.85 cm
- □ Cat 5E (Enhanced), Cat 6, Cat 7, ...





1-9



- -Outer conductor is braided shield
- -Inner conductor is solid metal
- -Separated by insulating material
- -Covered by padding

□ Higher bandwidth than UTP. Up to 500 MHz.

□ Used in cable TV



Optical Fiber



- □ A cylindrical mirror is formed by the cladding
- □ The light wave propagate by continuous reflection in the fiber
- □ Not affected by external interference \Rightarrow low bit error rate
- □ Fiber is used in all long-haul or high-speed communication
- □ Infrared light is used in communication

Wireless Transmission Frequencies

\Box 2GHz to 60GHz

- □ Terrestrial Microwave, Satellite Microwave
- Highly directional
- □ Point to point
- □ 30MHz to 1GHz
 - D Omni-directional
 - Broadcast radio
- **a** $3 \ge 10^{11}$ to $2 \ge 10^{14}$
 - □ Infrared
 - □ Short distance

Multiplexing

Time

How multiple users can share a link?
Time Division Multiplexing (TDM)



• Other multiplexing methods will be covered as needed.

Types of Networks



¹⁻¹⁵

Circuit vs. Packet Switching

| | Circuit Switching | Packet Switching |
|----------------------|--------------------------|---------------------|
| Call setup | Required | Optional |
| Overhead during call | Minimal | Per packet overhead |
| State | Stateful | No state |
| Resource Reservation | Easy | Difficult |
| Quality of Service | Easy | Difficult |
| Sharing | By overbooking | Easy |

Myth: Circuits require dedicated resources
 No sharing
 True only for constant bit rate (CBR) circuits

Types of Networks (Cont)

- Enterprise vs Telecom Networks
 Ethernet is the most common interface in Enterprise
 Frame relay and ATM are common in Telecom Networks
- Local Area Networks (LAN) 0-2 km, Single Ownership Metropolitan Area Networks (MAN) 2-50 km, Wide Area Networks (WAN) 50+ km
 - Originally LAN/MAN/WAN technologies were different
 - □ Now they are all same
- **Telecom** Networks:
 - □ Access: Between subscriber and the service provider
 - □ Metro: Covering a city
 - □ Core: Between cities

Homework 1A

- Which networking media will you use for the following applications and why?
- 1. Very large file transfer at home
- 2. High-speed multiple channel video transmission at office
- 3. News reading while traveling in a car



- □ Internet = Network connecting networks
- □ Approximately 600 million hosts on Internet in July 2008.
- □ ISP: Internet Service Provider.
 - □ Provide access to Internet.
 - Telecommunications (Telephone) Companies, AT&T, Verizon, Comcast, ...
 - □ Coffee Shops (Wi-Fi)



□ Core Network: ISP's network



Network Edge: Enterprise Networks

- 1. Ethernet
- 2. Wi-Fi

Ethernet

- □ Uses UTP (Unshielded Twisted Pair)
- **1** 10 Mbps, 100 Mbps, 1 Gbps, 10 Gbps
- □ Originally bus, now point-to-point (Star) topology





Wi-Fi

IEEE 802.11 Uses 2.4 GHz and 5.8 GHz



CSE473S

1-24



Access Networks

- 1. Dial Up
- 2. DSL
- 3. Cable
- 4. Fiber-To-The-Home
- 5. Wi-Fi
- 6. WiMAX



1-25



- Modem (Modulator/Demodulator) convert electrical bits to sound waveforms for transmission over telephone network
- □ Telephone network designed to carry 4 kHz voice
- □ Up to 56 kbps
- Does not need much help from the phone company

Bits, Hertz, and Baud

- Bits: Unit of information. Binary 0 or 1
- Bits are transmitted as pulses: E.g., Manchester encoding



0=High-to-low transition 1=Low-to-high transition

- Receiver design depends on the duration of smallest pulse 1kbps ⇒ One bit per millisecond ⇒ Each pulse is ½ ms ⇒ 2 kBaud
- □ The pulses become a mixture of sine waves on the medium
- $\square Wires allow only certain frequencies \Rightarrow Hertz = cycles/second$ $Washington University in St. Louis CSE473S <math>\bigcirc$ 2011 Raj Jain



DSL

- Digital Subscriber Line (DSL)
- Can transmit very high data rates on phone wire using special equipment at the phone company allowing higher frequency signals



Cable

- Cable companies have a very-high speed medium (for video transmission)
- Phone wire = 4kHz for voice
 Video Cable = 500 MHz for video
 One TV Channel = 6 MHz
- □ 30 Mbps down/1 Mbps up
- □ Fiber in the main line + Coax in tributaries ⇒ Hybrid Fiber Coax (HFC)





Cable

Modem





- □ No electronic components in the distribution system \Rightarrow Passive \Rightarrow Reliable
- Passive Optical Network (PON)

Wireless Access Networks

- □ Wi-Fi hot spots
- Cellular access
- □ WiMAX

| Washington | University | in St. Louis |
|--------------|------------|---------------|
| w asinington | University | III St. Louis |



Network Performance Measures

- Delay
- □ Throughput
- Loss Rate

| Washington | University | in St. Louis | |
|-----------------|------------|------------------|--|
| 11 WOITING VOIT | 0111,01010 | 111 0 01 10 0110 | |

Delay Example (CBR Circuits)

- How long would it take to send a file of 640,000 bits from host A to host B over a circuit-switched network?
 - □ All links are 1.536 Mbps
 - □ Each link is shared by 24 users
 - □ 500 ms to establish end-to-end circuit
- \Box Per User Rate = 1536/24 = 64 kbps
- $\Box \text{ Time to transfer} = 640 \text{kb}/64 \text{kb} = 10 \text{ s}$
- **Total time** = .5 s + 10 s = 10.5 s

Packet Switching Delay

- 1. Processing Delay: Check packets, decide where to send, etc.
- 2. Queuing Delay: Wait behind other packets
- 3. Transmission Delay: First-bit out to last-bit out on the wire= Packet Length/bit rate
- 4. Propagation Delay: Time for a bit to travel from in to out = Distance/speed of signal

Light speed = $3x10^8$ m/s in vacuum, $2x10^8$ m/s in fiber



Packet Switching Delay: Example

- □ 1500 Byte packets on 10 Mbps Ethernet, 1km segment
- □ Transmission Delay = $1500x8/10x10^6 = 1200 \ \mu s = 1.2ms$
- Propagation delay = $1000 \text{ m}/2x10^8 = 5 \mu \text{s}$

Throughput

- □ Measured in Bits/Sec
- **Capacity:** Nominal Throughput
- **Throughput: Realistic**
- Bottleneck determines the end-to-end throughput



Net end-to-end capacity = 10 Mbps

Actual throughput will be less due to sharing and overhead.

1 - 38

Loss Rate

- $\Box Queuing \Rightarrow Buffer overflow$
- **Bit** Error Rate on the link
- Lost packets are retransmitted by the previous node or the source



Homework 1B

- P5: Consider two hosts, A and B, connected by a single link of rate R bps. Suppose that the two hosts are separated by m meters, and suppose the propagation speed along the link is s meters/sec. Host A is to send a packet of size L bits to Host B.
- A. Express the propagation delay, d_{prop} in terms of *m* and *s*
- B. Determine the transmission time of the packet d_{trans} in terms of L and R.
- C. Ignoring processing queuing delays, obtain an expression for the end-to-end delay
- D. Suppose Host A begins to transmit the packet at time t=0. At time $t=d_{trans}$ where is the last bit of the packet?
- E. Suppose d_{prop} is greater than d_{trans} . At time $t=d_{trans}$, where is the first bit of the packet?
- F. Suppose d_{prop} is less than d_{trans} , at time $t=d_{trans}$, where is the first bit of the packet
- G. Suppose $s=2.5 \times 10^8$ m/s, L=240 bits, and R=56 kbps,. Find the distance *m* so that d_{prop} equals d_{trans} .

Protocol Layers

 Problem: Philosophers in different countries speak different languages. The Telex system works only with English.



ISO/OSI Reference Model



File transfer, Email, Remote Login ASCII Text, Sound Establish/manage connection End-to-end communication: TCP Routing, Addressing: IP Two party communication: Ethernet How to transmit signal: Coding







Protocol Data Unit (PDU)

| Application | APDU, Message | Application |
|--------------|---|--------------|
| Presentation | ← PPDU | Presentation |
| Session | SPDU | Session |
| 56551011 | TPDU | 56551011 |
| Transport | NIDDLI Dealtat | Transport |
| Network | | Network |
| Datalink | DPDU, Frame | Datalink |
| Physical | And | Physical |



2. Indication

Unconfirmed service: No confirmation or response

Washington University in St. Louis

CSE473S 1-46

4. Confirm

TCP/IP Reference Model

- □ TCP = Transport Control Protocol
- □ IP = Internet Protocol (Routing)
- TCP/IP Ref Model TCP/IP Protocols

| Application | FTP | | Teln | et | HTTP |
|--------------------|--------------|---|------------------|-----|-----------------|
| Transport | TCP | | | UDP | |
| Internetwork | IP | | | | |
| Host to Network | Ether net | Р | oint-te Point | 0- | Packet Radio |
| Physical | Coax |] | Fiber | W | vireless |

Washington University in St. Louis

CSE473S

OSI vs TCP/IP

| OSI | TCP/IP | | | |
|--------------|----------------|--|--|--|
| Application | | | | |
| Presentation | Application | | | |
| Session | | | | |
| | Transport | | | |
| Transport | (host-to-host) | | | |
| Network | Internet | | | |
| Network | | | | |
| Data Link | Network | | | |
| Data Lilik | Access | | | |
| Physical | Physical | | | |
| CSE473S | | | | |

©2011 Raj Jain

Washington University in St. Louis

OSI vs TCP Reference Models

- OSI introduced concept of services, interface, protocols. These were force-fitted to TCP later ⇒ It is not easy to replace protocols in TCP.
- In OSI, reference model was done before protocols. In TCP, protocols were done before the model
- OSI: Standardize first, build later TCP: Build first, standardize later
- OSI took too long to standardize.
 TCP/IP was already in wide use by the time.
- □ OSI became too complex.
- □ TCP/IP is not general. Ad hoc.



¹⁻⁵⁰

TCP/IP Applications



| BGP = Border Gateway Protocol FTP = File Transfer Protocol HTTP = Hypertext Transfer Protocol ICMP = Internet Control Message Protocol IGMP = Internet Group Management Protocol IP = Internet Protocol | OSPF RSVP SMTP SNMP TCP UDP | | Open Shortest Path First Resource ReSerVation Protocol Simple Mail Transfer Protocol Simple Network Management Protocol Transmission Control Protocol User Datagram Protocol | |
|--|--|----|---|----------------|
| IP = Internet Protocol MIME = Multi-Purpose Internet Mail Extension | UDP | = | User Datagram Protocol | |
| Washington University in St. Louis | | CS | E473S | ©2011 Raj Jain |

1-51

Network Security

- Security Components
- □ Types of Malware
- □ Types of Attacks
- Buffer Overflows
- Distributed DoS Attacks

Security Components

- Confidentiality: Need access control, Cryptography, Existence of data
- Integrity: No change, content, source, prevention mechanisms, detection mechanisms
- □ Availability: Denial of service attacks,
- □ Confidentiality, Integrity and Availability (CIA)



1-53

Types of Malware

- □ Viruses: Code that *attaches* itself to programs, disks, or memory to propagate itself.
- □ Worms: Installs copies of itself on other machines on a network, e.g., by finding user names and passwords
- Trojan horses: Pretend to be a utility. Convince users to install on PC.
- **Spyware**: Collect personal information
- **Hoax**: Use emotion to propagate, e.g., child's last wish.
- **Trap Door**: Undocumented entry point for debugging purposes
- Logic Bomb: Instructions that trigger on some event in the future
- Zombie: Malicious instructions that can be triggered remotely. The attacks seem to come from other victims.

Types of Attacks

- Denial of Service (DoS): Flooding with traffic/requests
- Buffer Overflows: Error in system programs. Allows hacker to insert his code in to a program.
- □ Malware
- **Brute Force**: Try all passwords.
- □ Port Scanning:
 - \Rightarrow Disable unnecessary services and close ports
- Network Mapping

Buffer Overflows

- □ Return address are saved on the top of stack.
- Parameters are then saved on the stack.
- □ Writing data on stack causes stack overflow.
- Return the program control to a code segment written by the hacker.



Distributed DoS Attacks

- Tribe Flood Network (TFN) clients are installed on compromised hosts.
- All clients start a simultaneous DoS attack on a victim on a trigger from the attacker.
- **Trinoo** attack works similarly. Use UDP packets. Trinoo client report to Trinoo master when the system comes up.
- Stacheldraht uses handlers on compromised hosts to receive encrypted commands from the attacker.



History of Internet

1961: Kleinrock developed queueing theory. Showed effectiveness of packet-switching



THE ARPA NETWORK

- 1964: Baran's report on packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: First ARPAnet node operational First Request for Comment (RFC) www.ietf.org

Internet Generations



Washington University in St. Louis

1-59

©2011 Rai Jain

History of Internet (Cont)

- □ Early 1990s: HTML, HTTP: Berners-Lee
- 1994: Mosaic, later Netscape

2007:

- □ ~500 million hosts
- □ Voice, Video over IP
- P2P applications: BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- □ Video applications: YouTube, gaming
- Wireless, Mobility

Key Concepts

- □ Internet Protocol (IP): Protocol
- Address: All systems have an IP address, for example, 125.36.47.23
- Name: All systems have a human readable name, e.g., scorpio.cec.wustl.edu, ibm.com.
- Technically called DNS (domain name systems) name. Details will be introduced later.
- IETF: Internet Engineering Task Force. Make standards for Internet. IETF.org
- RFC: Request for comments. Documents that describe Internet protocols.

Homework 1C

- 1. Find the IP address of your computer
- 2. Find the IP address of <u>www.google.com</u> (different from google.com)
- 3. Measure delay from your computer to <u>www.google.com</u>

For all cases submit the screen snapshot showing the command used and the output. (Use Alt-Print-screen to capture a window to clipboard and then paste to word)



- 1. Most common medium is UTP, wireless, fiber
- 2. Internet is a network of networks
- 3. Enterprise, access, and core networks
- 4. Performance Measures: Delay, Throughput, Loss Rate
- 5. Protocol Layers: ISO and TCP/IP reference models