Multimedia Networking

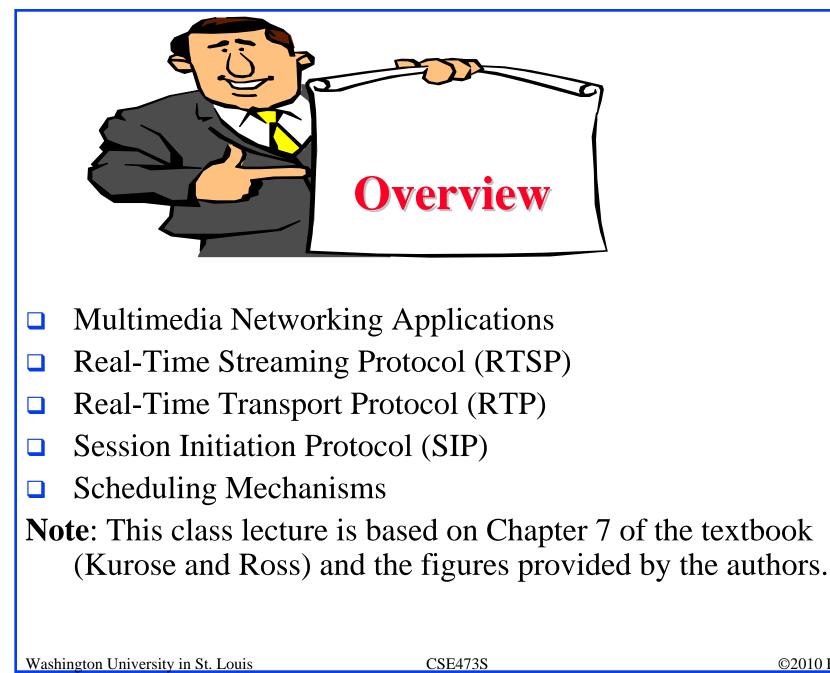
Raj Jain

Washington University in Saint Louis Saint Louis, MO 63130 Jain@wustl.edu

Audio/Video recordings of this lecture are available on-line at:

http://www.cse.wustl.edu/~jain/cse473-10/

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Multimedia Networking Applications

□ Streaming Stored Audio and Video

□ Stored Media: Fast rewind, pause, fast forward

- □ Streaming: simultaneous play out and download
- Continuous play out: Delay jitter smoothed by playout buffer
- Streaming Live Audio and Video: IPTV and Internet Radio
 No fast-forward
- High data rate to large number of users \Rightarrow multicast or P2P,

□ delay jitter controlled by caching,

Real-Time Interactive Audio and Video: Internet Telephone, Video Conferencing

 \Box Delay<400 ms.

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Multimedia on Internet

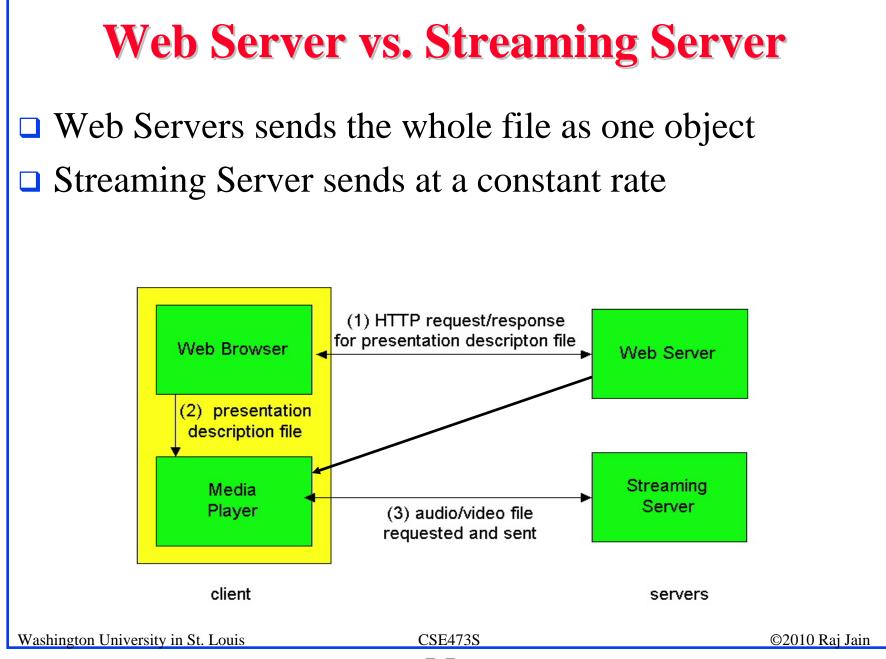
- Best Effort Service
- **TCP** not used due to retransmission delays
- □ Limited packet loss tolerated
- Packet jitter smoothed by buffering
- □ Hard Guarantee: Min Throughput, Max Delay, Max delay jitter
- □ Soft Guarantee: Quality of service with a high probability
- □ Protocol for Bandwidth Reservation and Traffic Description
- □ Scheduling to honor bandwidth reservation
- High Bandwidth
- Content Distribution Networks: Akamai

Audio Compression Standards

- □ 4kHz audio \Rightarrow Audio sampled at 8000 samples per second
- □ 256 levels per sample \Rightarrow 8 bits/sample \Rightarrow 64 kbps
- □ Pulse Code Modulation (PCM)
- CD's use 44.1 kSamples/s, 16 b/sample ⇒ 705.6 kbps (mono) or 1.411 Mbps (Stereo)
- GSM Cell phones: 13 kbps
- **G**.711: 64 kbps
- **G**.729: 8 kbps
- **G**.723.3: 6.4 and 5.3 kbps
- □ MPEG 1 Layer 3 (MP3): 96 kbps, 128 kbps, or 160 kbps

Video Compression Standards

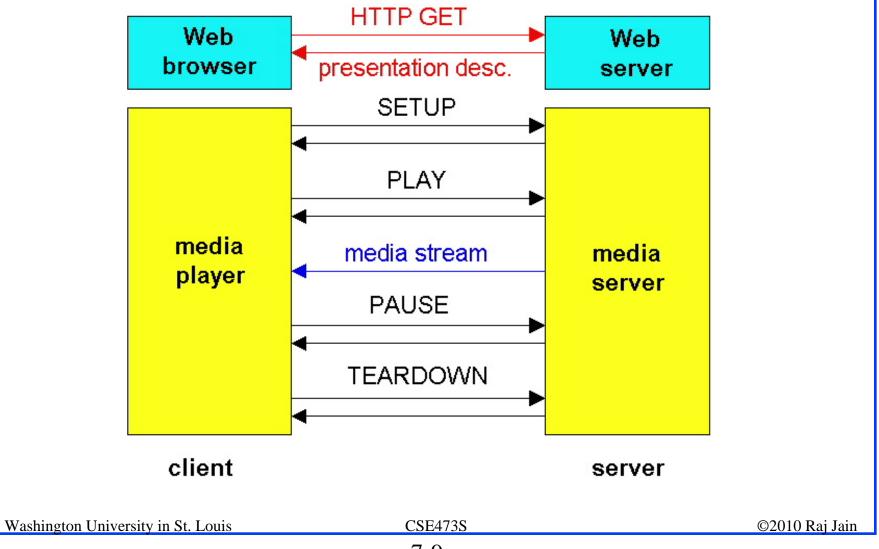
- □ Moving Pictures Expert Group (MPEG)
- □ MPEG 1: CD quality video (1.5 Mbps)
- □ MPEG 2: DVD quality Video 3-6 Mbps
- MPEG 4: Low-rate high-quality video (.divx or .mp4)
 H.261



Real-Time Streaming Protocol (RTSP)

- □ Protocol to control streaming media
- □ Allows start, stop, pause, fast forward, rewinding a stream
- Data and control channels
- □ All commands are sent on control channel (Port 544)
- Specified as a URL in web pages: rtsp://www.cse.wustl.edu/~jain/cse473-09/ftp/i_7mmn0.rm

RTSP Operation



RTSP Exchange Example

- C: SETUP rtsp://audio.example.com/twister/audio RTSP/1.0 Transport: rtp/udp; compression; port=3056; mode=PLAY
- S: RTSP/1.0 200 1 OK

Session 4231

C: PLAY rtsp://audio.example.com/twister/audio.en/lofi RTSP/1.0 Session: 4231

Range: npt=0-

C: PAUSE rtsp://audio.example.com/twister/audio.en/lofi RTSP/1.0 Session: 4231

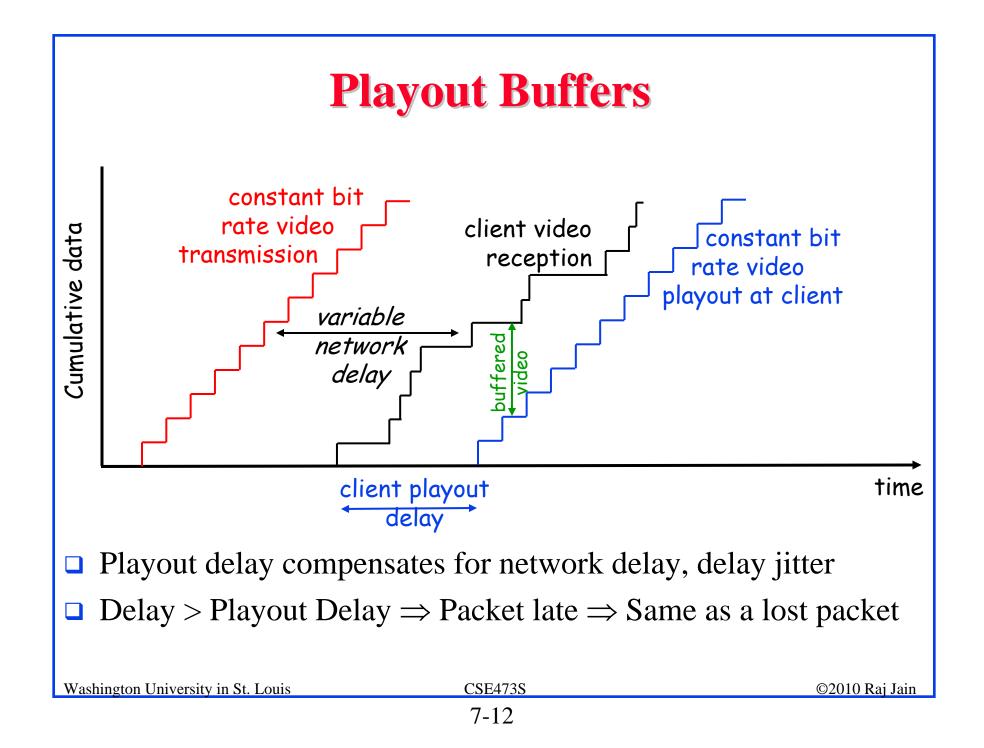
Range: npt=37

C: TEARDOWN rtsp://audio.example.com/twister/audio.en/lofi RTSP/1.0 Session: 4231

S: 200 3 OK

Multimedia with Best Effort Service

- □ High Compression \Rightarrow Low Rate \Rightarrow Low loss
- □ 1% to 20% loss can be concealed
- □ Forward Error Correction (FEC) can be used to overcome loss.
- □ End-to-end delay limited to 400 ms
- □ Jitter overcome by play out buffer
- □ Large jitter \Rightarrow Packets arrive too late \Rightarrow same as Lost
- □ Each chunk comes with a sequence number and timestamp
- Play out delay can be adaptively adjusted according to measured delay variation



Adaptive Playout Delay

- \Box t_i=Sending time
- \Box r_i= Receiving time
- $\square Measured delay sample = r_i t_i$

 \Box d_i= Average network delay

 $d_i = (1-a)d_{i-1} + a(r_i - t_i)$

 \Box v_i= Variation of the delay

$$v_i = (1-a)v_{i-1} + a|r_i - t_i - d_i|$$

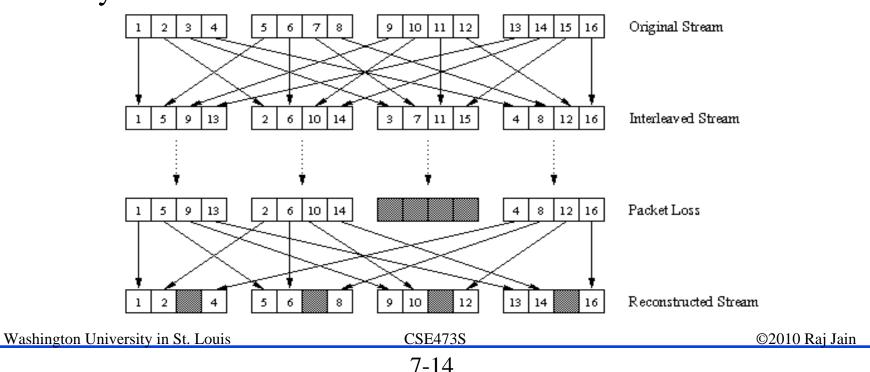
 \Box p_i= Playout time

$$p_i = t_i + d_i + Kv_i$$

□ Here K is a constant, say 4.

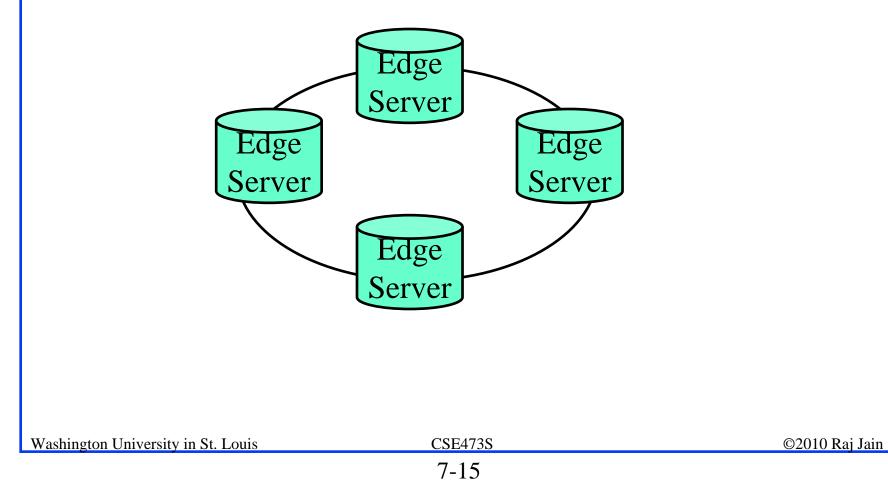
Recovering From Packet Loss

- Forward Error Correction
- □ Send n+1 packets in place of n packets
- □ Send a lower-resolution stream in addition
- □ Play out the old syllable
- □ Busty Loss \Rightarrow Interleave audio/video frames



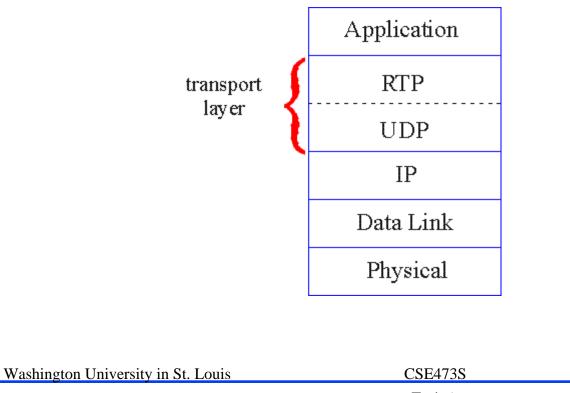
Content Distribution Networks

Authoritative DNS server resolves the server address according to the requester's IP address



Real-Time Transport Protocol (RTP)

- Common sublayer between applications and UDP
- Provides sequence numbers, timestamps, and other facilities
- Supports both unicast and multicast



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RTP Packet Format												
Paylcad Type			Timestamp		hron zation ce Identifer	Miscellaneous Fields						
8b	8b		32b	32b								
□ SSRC = Synchronization Source Identifier = Stream #												
Pa	ayload	Coding		Rate								
	ype	U										
	0	PCM mu-law		64 kbps								
	3		GSM		5							
	7		LPC		S							
	26		Motion JPEG									
	31		H.261									
	33		MPEG2 video									

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RTP Control Protocol (RTCP)

- □ Used to send report about reception quality back to sender
- □ Also used by sender to report stream information
- Can be used to adjust the transmission speed, quality, or for diagnosis
- □ SSRC
- □ Fraction of packets lost
- □ Last sequence number received
- □ Inter-arrival jitter
- □ Receiver report rate is adjusted inversly to number of receivers
- Sender report rate is adjusted inversly to number of senders
- □ Total RTCP traffic < 5% of media datarate

Session Initiation Protocol (SIP)

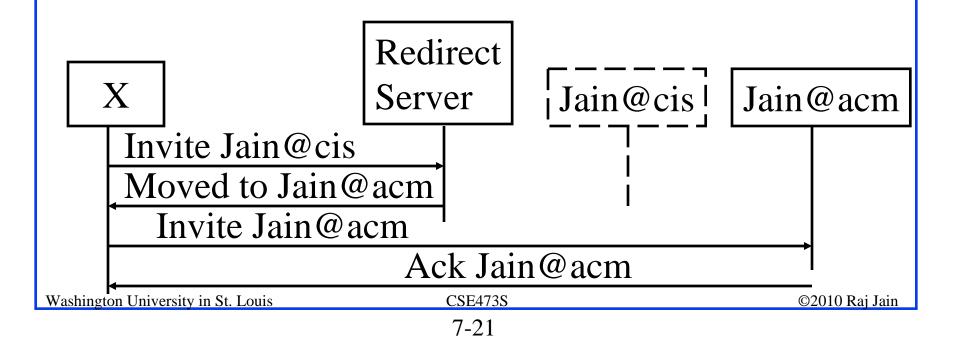
- Application level signaling protocol for voice and video conferencing over Internet
- Allows creating, modifying, terminating sessions with one or more participants
- Carries session descriptions (media types) for user capabilities negotiation
- □ Supports user location, call setup, call transfers
- □ Supports mobility by proxying and redirection

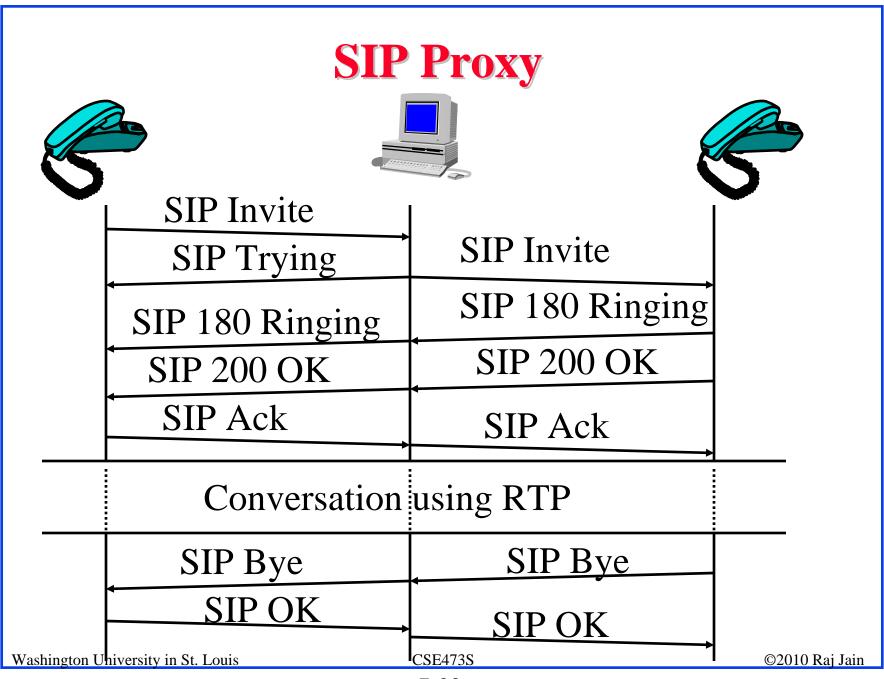
SIP (Cont)

- SIP Uniform Resource Identifiers (URIs): Similar to email URLs sip:jain@cis.ohio-state.edu sip:+1-614-292-3989:123@osu.edu?subject=lecture
- □ SIP can use UDP or TCP
- □ SIP messages are sent to SIP servers:
 - Registrar: Clients register and tell their location to it
 Location: Given name, returns possible addresses for a user. Like Directory service or DNS.
 - □ Redirect: Returns current address to requesters
 - Proxy: Intermediary. Acts like a server to internal client and like a client to external server

Locating using SIP

- □ Allows locating a callee at different locations
- □ Callee registers different locations with Registrar
- □ SIP Messages: Ack, Bye, Invite, Register, Redirection, ...





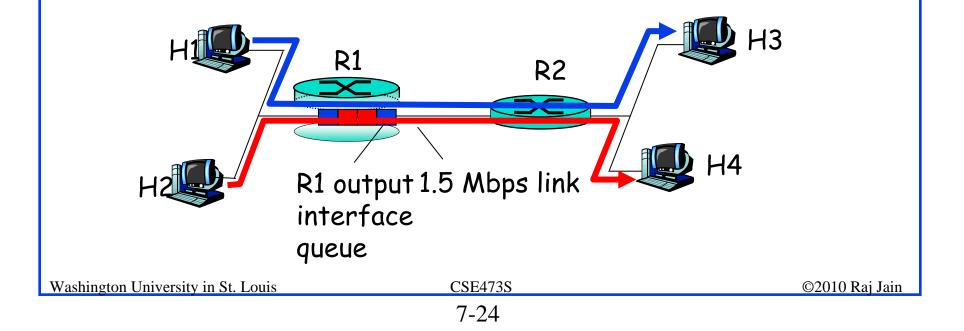
H.323 Protocols

- □ Multimedia over LANs, V1 (June 96), V2(Feb 98)
- Provides component descriptions, signaling procedures, call control, system control, audio/video codecs, data protocols

Vide	0	Audio	(Data							
H.26 H.26	13	G.711, G.722, G.723.1, G.728, G.729	RTCP		H.225.0 Signaling	H.245 Control	T.124				
	RTP			X.224 Class 0			T.125				
		UDP			T.123						
	Network (IP)										
Datalink (IEEE 802.3)											
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Multiple Classes of Service

- Flow Classification: Based on Source IP, Dest IP, Source Port, Dest Port, Type of Service, ...
- Differentiation: Routers can provide different service to different traffic
- □ Isolation: One class cannot affect other classes severly



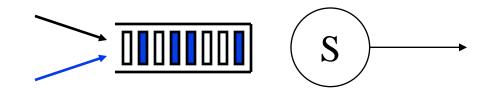
Scheduling Mechanisms

How to service multiple flows?

- First Come First Served Scheduling
- Priority Queueing
- Round Robin Scheduling
- Generalized Processor Sharing
- □ Fair Queueing
- □ Weighted Fair Queueing (WFQ)
- **Desired Properties:**
 - 🗆 Fair

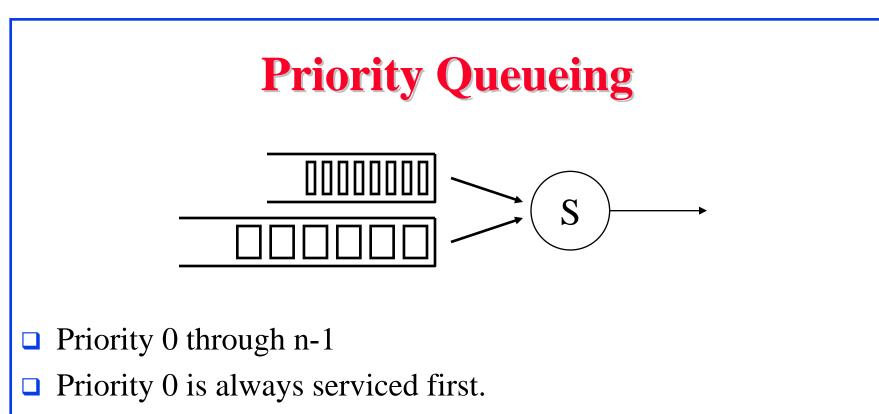
Work-Conserving: Do not waste resources if there is no traffic

First Come First Served Scheduling

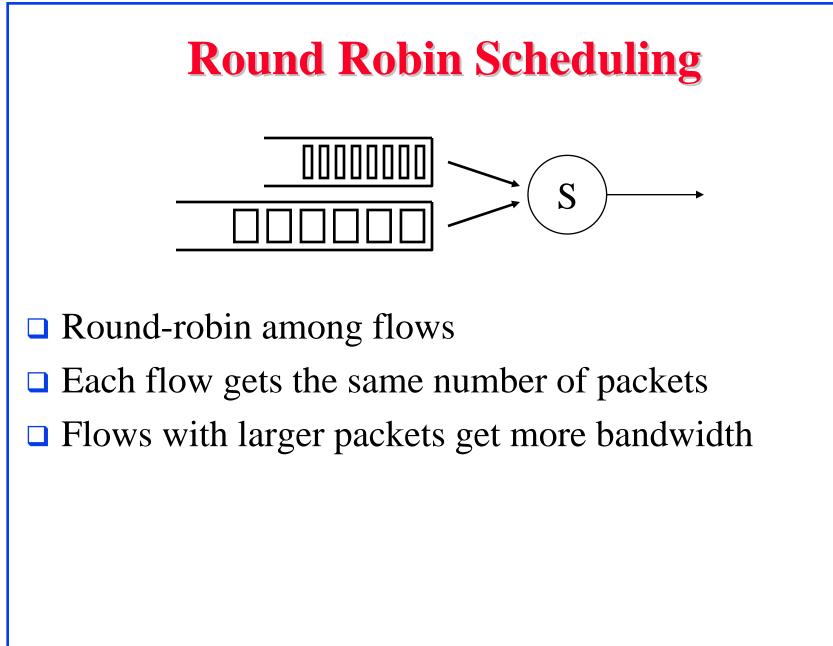


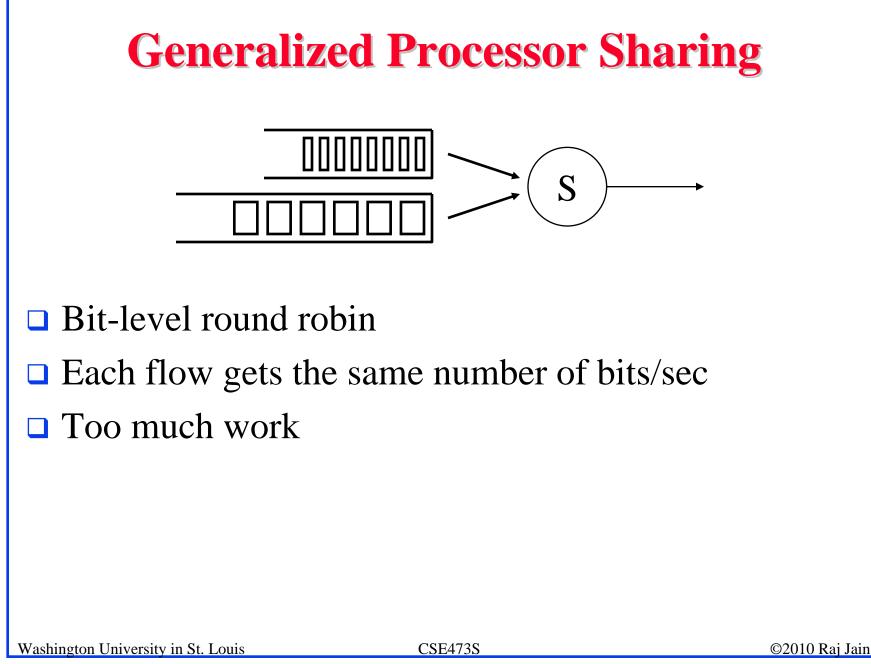
□ Unfair: Overloading flows get more service

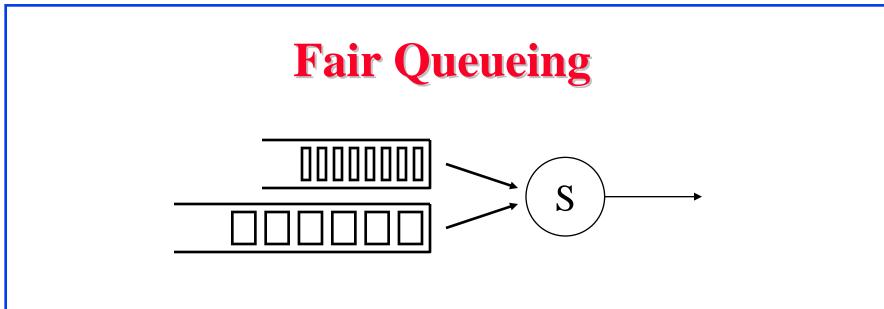
□ No isolation among users



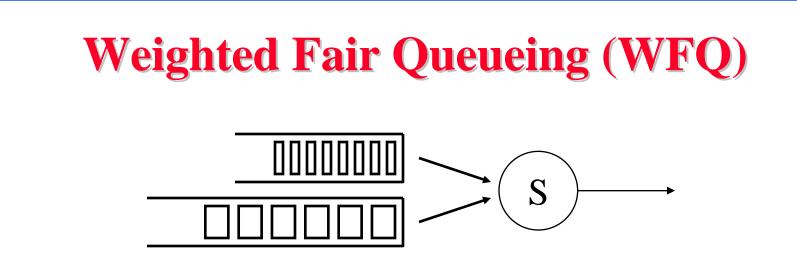
- □ Priority i is serviced only if 0 through i-1 are empty
- Highest priority has the lowest delay, highest throughput, lowest loss
- Lower priority classes may be starved if higher priority are overloaded







- □ Bit-level round robin but packet level scheduling
- Count the packet size and determine which packet would finish first. Serve that packet.
- □ Each flow gets the same number of bits/sec



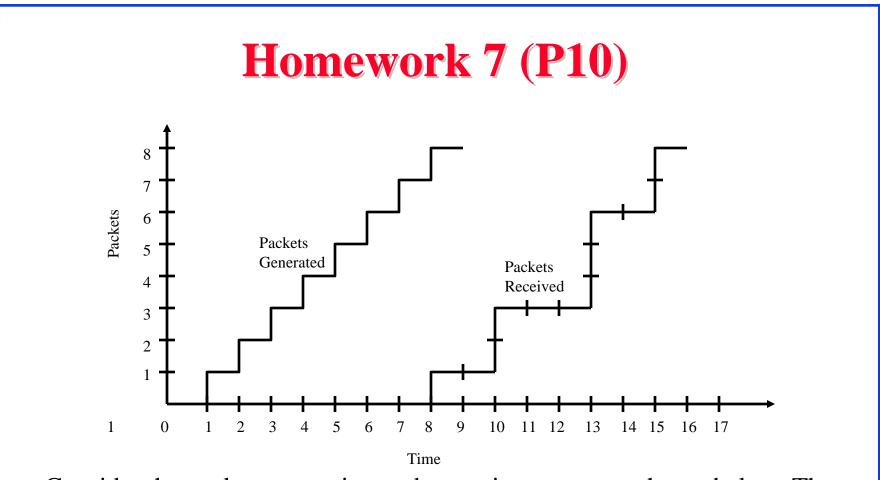
- □ Fair queueing with different weight for each queue
- □ Flow 1 gets x bit/sec
- □ Flow 2 gets y bit/sec
- □ Flow n gets z bit/sec
- □ Here, x, y, z are weights



- Multimedia applications require bounded delay, delay jitter, and minimum throughput
- Three Approaches: Service guarantees, Simple priority type service, Increase Capacity
- □ RTSP allows streaming controls like pause, forward, ...
- **RTP** allows sequencing and timestamping
- □ SIP allows parameter negotiation and location
- □ Weighted fair queueing allows packet based fair scheduling

Review Exercises

- □ Read Pages 597-657 of the textbook.
- □ Review Exercises R1-R15
- □ Problems P2-P4,P9, P11, P16, P19-P22



- □ Consider the packet generation and reception sequence shown below. The first packet is generated at t=1 and is received at t=8.
- A. If Playout delay is zero and playout begins at t=8, which of the packets will not arrive in time?
- B. What is the minimum playout delay at the receiver that result in all of the first eight packets arriving in time for their playout?
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