Computer Networks and the Internet

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Audio/Video recordings of this lecture are available on-line at:

http://www.cse.wustl.edu/~jain/cse473-10/

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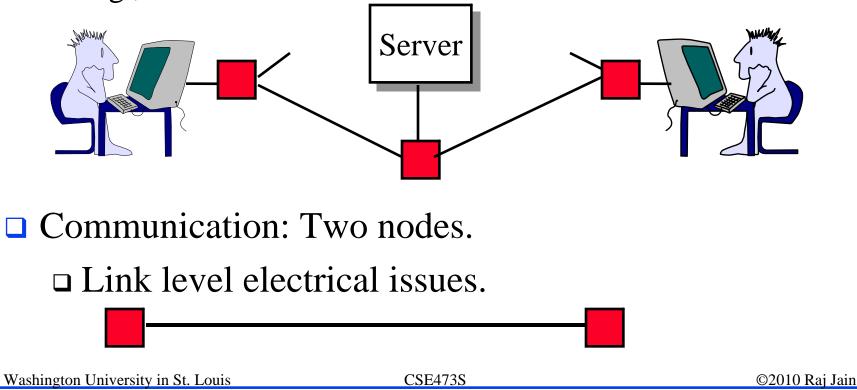
- 1. Physical Media
- 2. Switching: Circuit vs. Packet
- 3. Internet:Edge, Core
- 4. Network Performance Measures: Delay, Loss, Throughput
- 5. Protocol Layers
- 6. Network Security
- 7. History

Note: This class lecture is based on Chapter 1 of the textbook (Kurose and Ross) and the slides provided by the authors.

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What is a Network?

- Network: Enables data transfer among nodes
 Generally heterogeneous nodes
 - □ More than 2 nodes
 - □ E.g., Your home or office network





- End Systems: Systems that are sinks or sources of data, e.g., Desktops, Laptops, Servers, Printers, Cell Phones, etc.
- □ Intermediate Systems: Systems that forward/switch data from one link to another, e.g., routers, switches
- □ **Hosts**: End Systems
- **Gateways:** Routers
- Servers: End Systems that provide service, e.g., print server, storage server, Mail server, etc.
- **Clients**: End systems that request service
- Links: Connect the systems.
 Characterized by transmission rate, propagation delay

Transmission Media

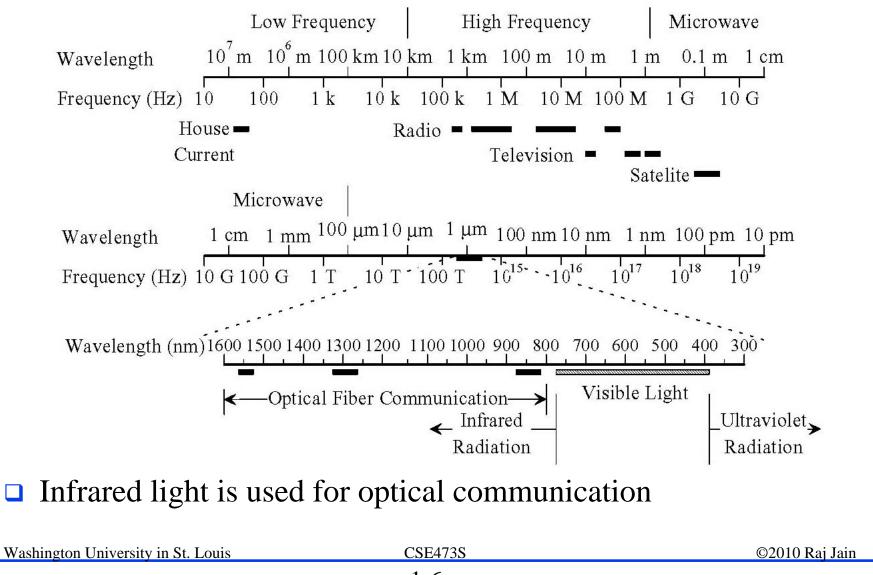
Guided:

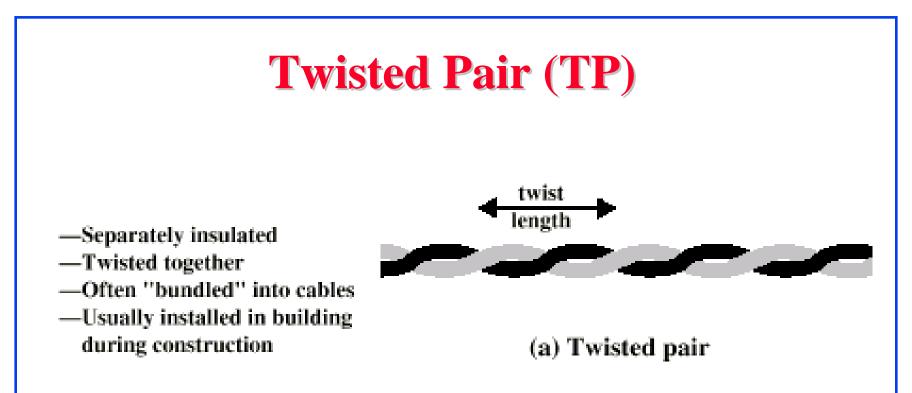
- □ Twisted Pair
- □ Coaxial cable
- □ Optical fiber

Unguided:

- □ Microwave
- □ Satellite
- □ Wireless

Electromagnetic Spectrum





□ Twists decrease the cross-talk

- Neighboring pairs have different twist length
- Most of telephone and network wiring in homes and offices is TP.

Shielded and Unshielded TP

- □ Shielded Twisted Pair (STP)
 - □ Metal braid or sheathing that reduces interference
 - □ More expensive
 - □ Harder to handle (thick, heavy)
 - □ Used in token rings
- □ Unshielded Twisted Pair (UTP)
 - □ Ordinary telephone wire
 - □ Cheap, Flexible
 - \Rightarrow Easiest to install
 - □ No shielding
 - \Rightarrow Suffers from external interference
 - □ Used in Telephone and Ethernet

UTP Categories

Cat 3

□ Up to 16MHz

□ Voice grade found in most offices

□ Twist length of 7.5 cm to 10 cm

Cat 4

□ Up to 20 MHz. Not used much in practice.

Cat 5

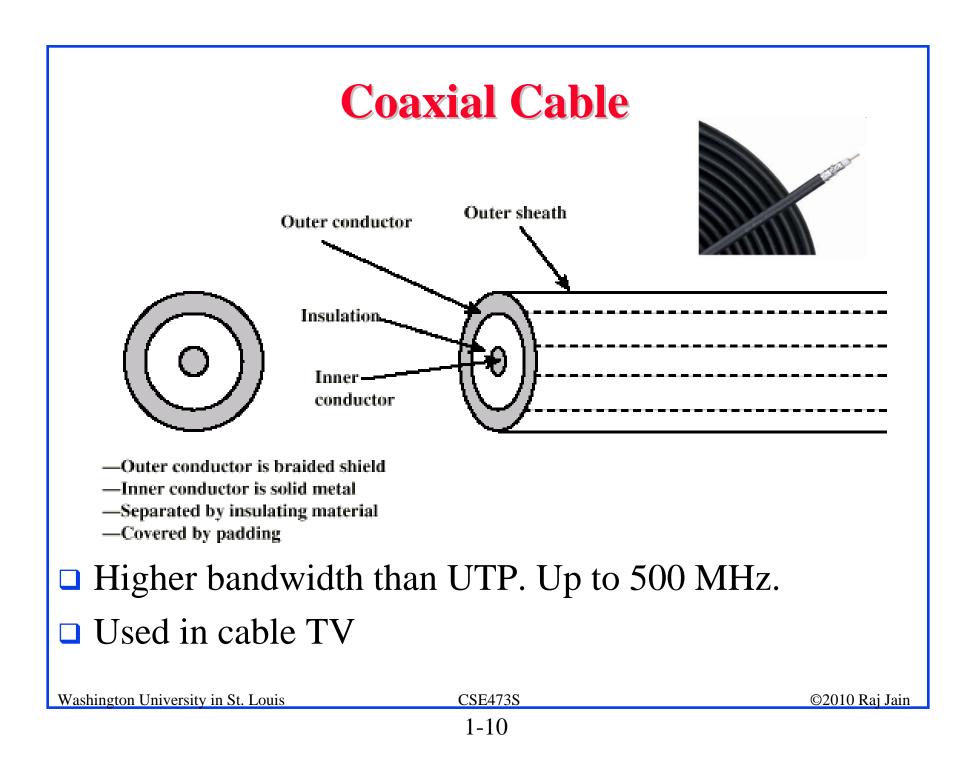
□ Up to 100MHz

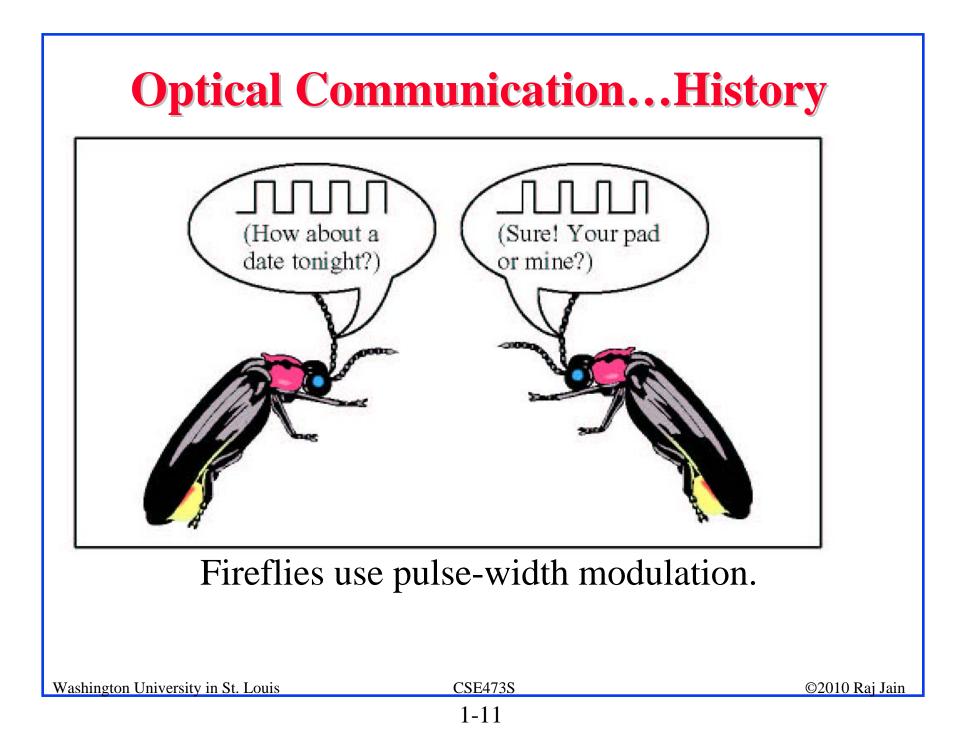
□ Used in 10 Mbps and 100 Mbps Ethernet

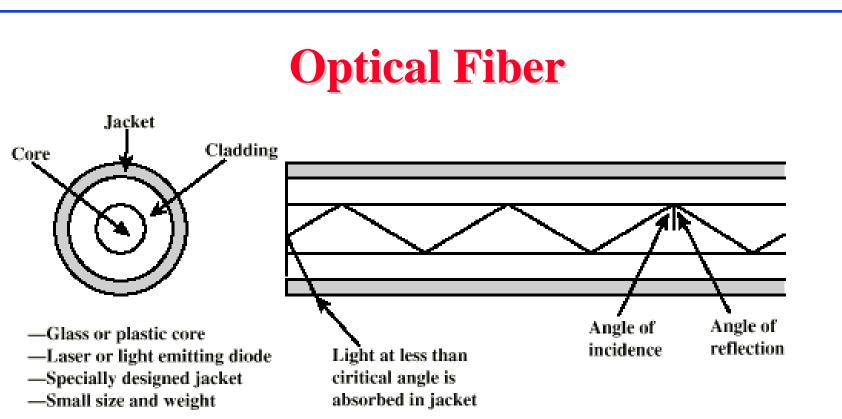
□ Twist length 0.6 cm to 0.85 cm

□ Cat 5E (Enhanced), Cat 6, Cat 7, ...









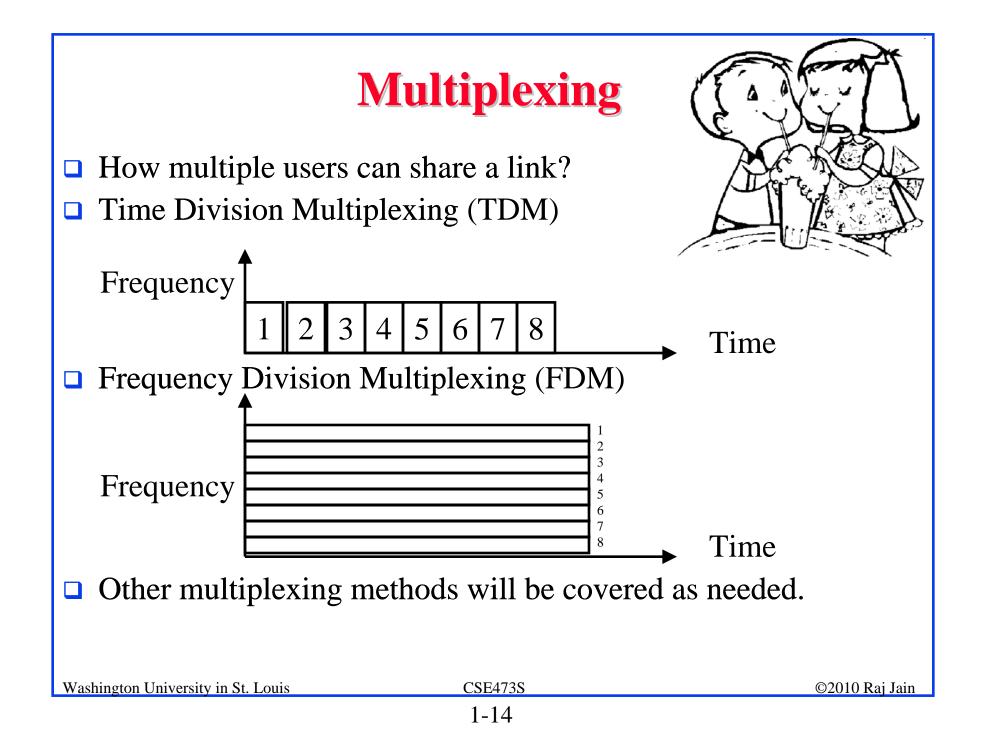
- □ A cylindrical mirror is formed by the cladding
- □ The light wave propagate by continuous reflection in the fiber
- □ Not affected by external interference \Rightarrow low bit error rate
- □ Fiber is used in all long-haul or high-speed communication
- □ Infrared light is used in communication

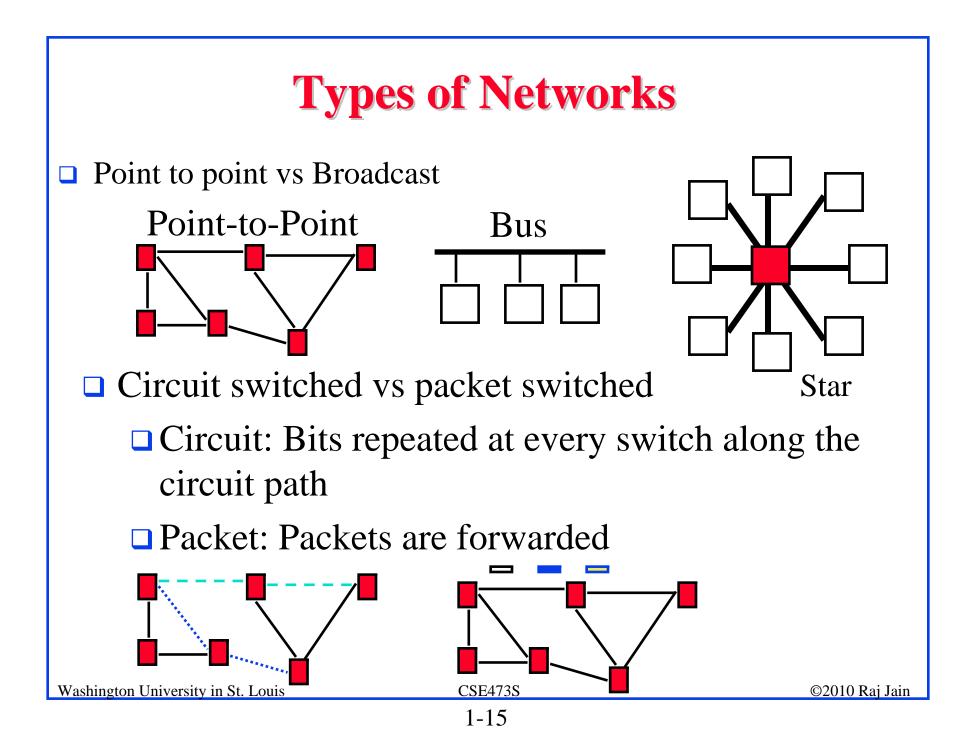
Wireless Transmission Frequencies

\Box 2GHz to 60GHz

Terrestrial Microwave, Satellite Microwave

- □ Highly directional
- □ Point to point
- □ 30MHz to 1GHz
 - □ Omni-directional
 - Broadcast radio
- **3** x 10¹¹ to 2 x 10¹⁴
 - □ Infrared
 - □ Short distance





Circuit vs. Packet Switching

	Circuit Switching	Packet Switching
Call setup	Required	Optional
Overhead during call	Minimal	Per packet overhead
State	Stateful	No state
Resource Reservation	Easy	Difficult
Quality of Service	Easy	Difficult
Sharing	By overbooking	Easy

Myth: Circuits require dedicated resources
 No sharing
 True only for constant bit rate (CBR) circuits

Types of Networks (Cont)

- Enterprise vs Telecom Networks
 Ethernet is the most common interface in Enterprise
 Frame relay and ATM are common in Telecom Networks
- Local Area Networks (LAN) 0-2 km, Single Ownership Metropolitan Area Networks (MAN) 2-50 km, Wide Area Networks (WAN) 50+ km

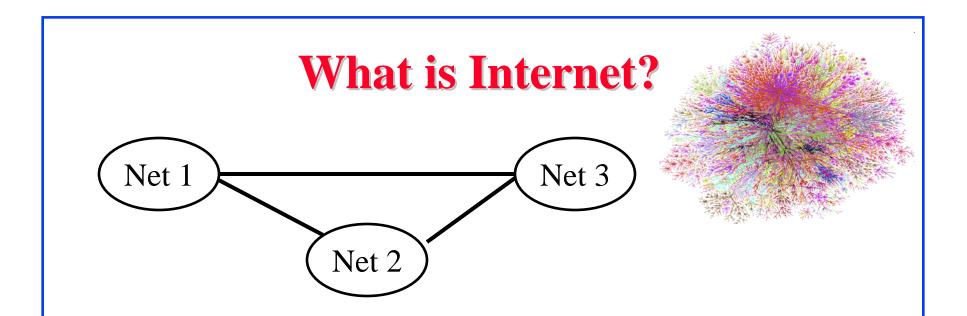
□ Originally LAN/MAN/WAN technologies were different

 \Box Now they are all same

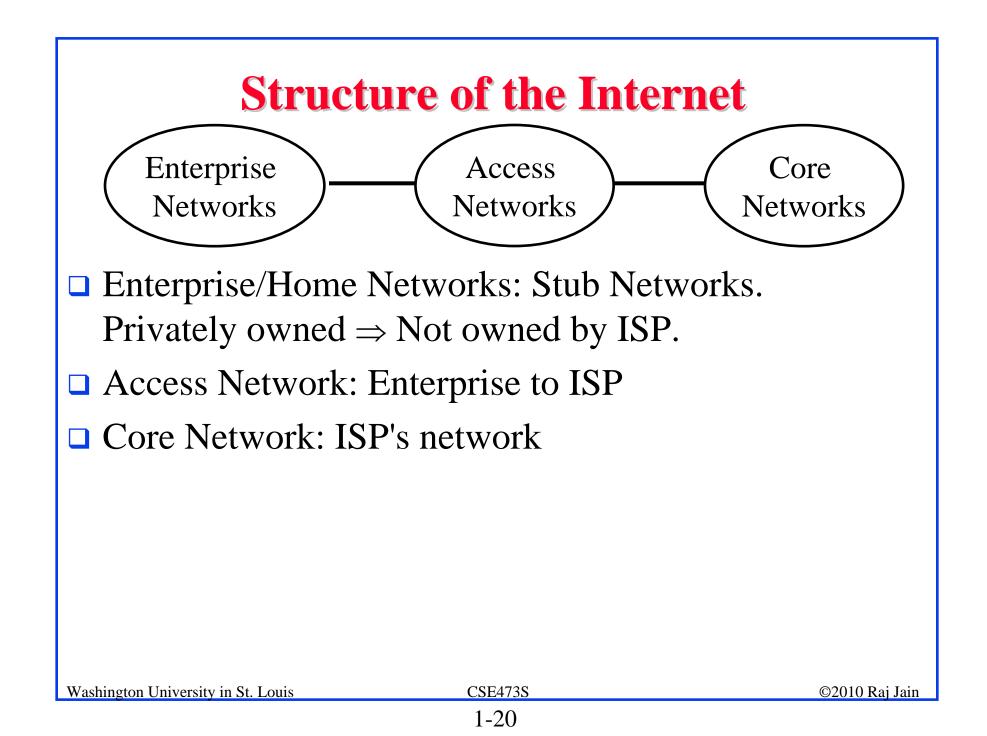
- **Telecom** Networks:
 - □ Access: Between subscriber and the service provider
 - □ Metro: Covering a city
 - □ Core: Between cities

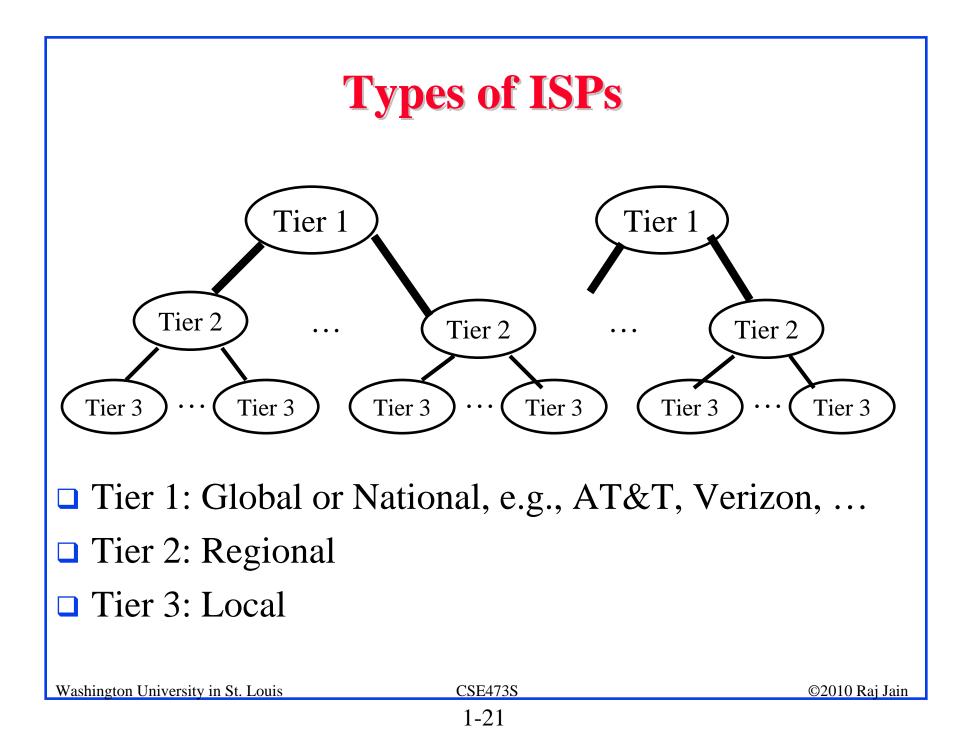
Homework 1A

- Which networking media will you use for the following applications and why?
- 1. Very large file transfer at home
- 2. High-speed multiple channel video transmission at office
- 3. News reading while traveling in a car



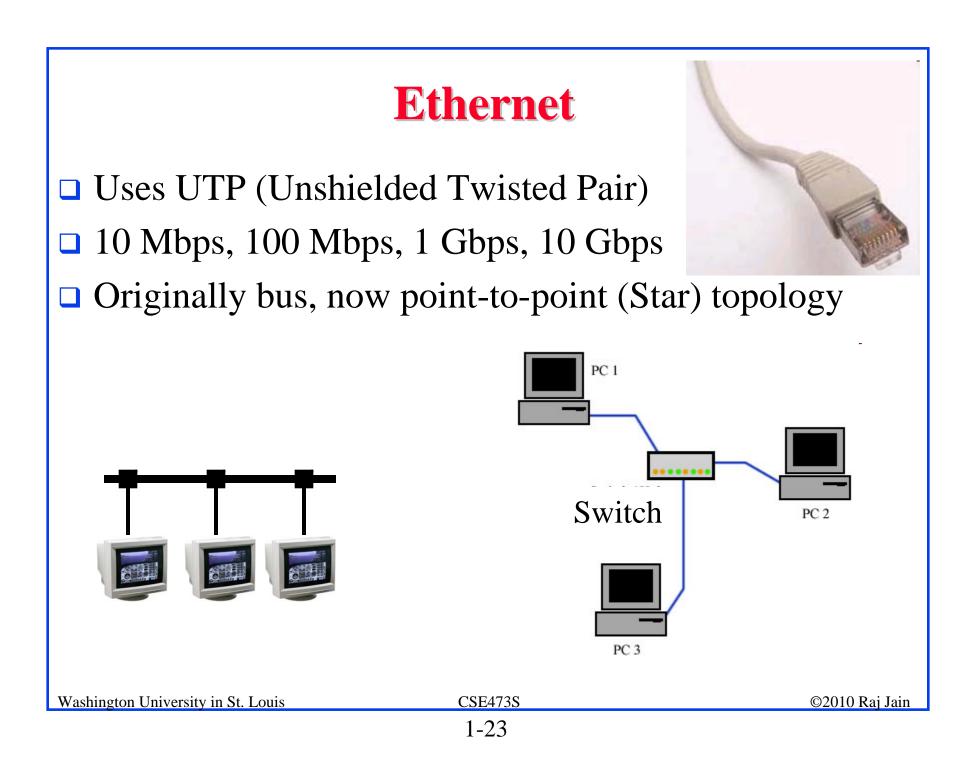
- □ Internet = Network connecting networks
- Approximately 600 million hosts on Internet in July 2008.
- □ ISP: Internet Service Provider.
 - □ Provide access to Internet.
 - Telecommunications (Telephone) Companies, AT&T, Verizon, Comcast, ...
 - □ Coffee Shops (Wi-Fi)

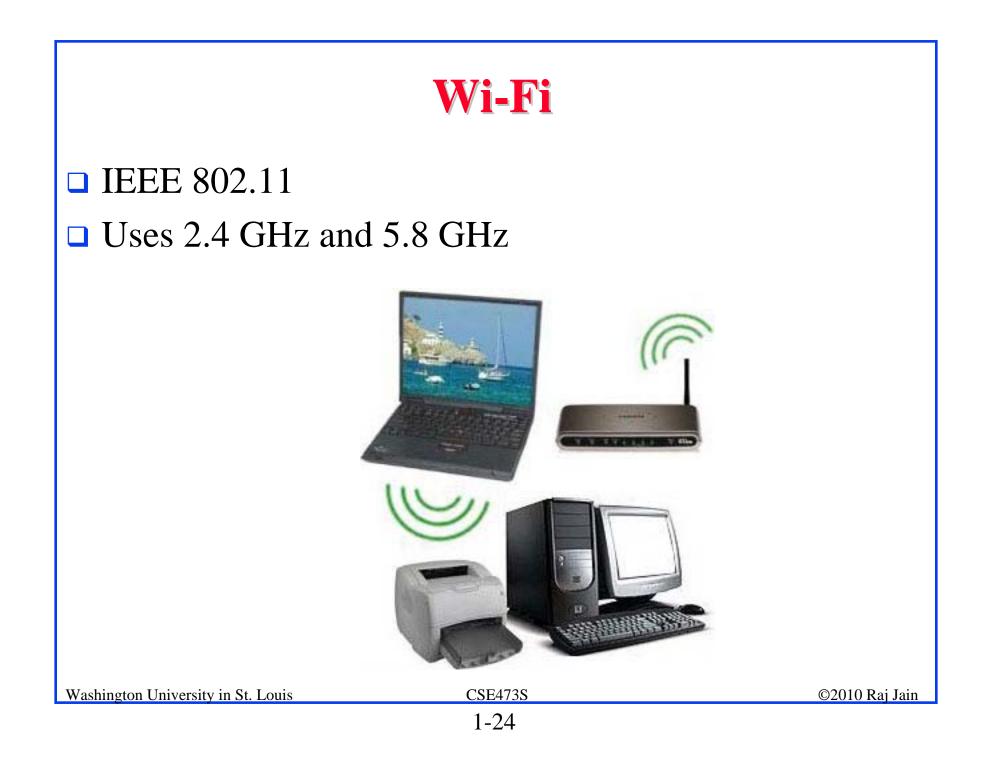




Network Edge: Enterprise Networks

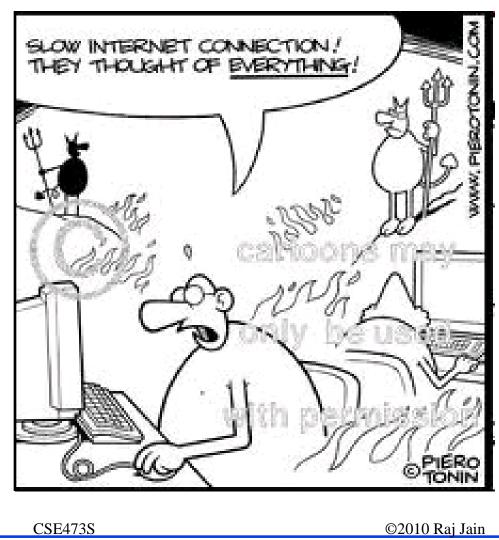
- 1. Ethernet
- 2. Wi-Fi



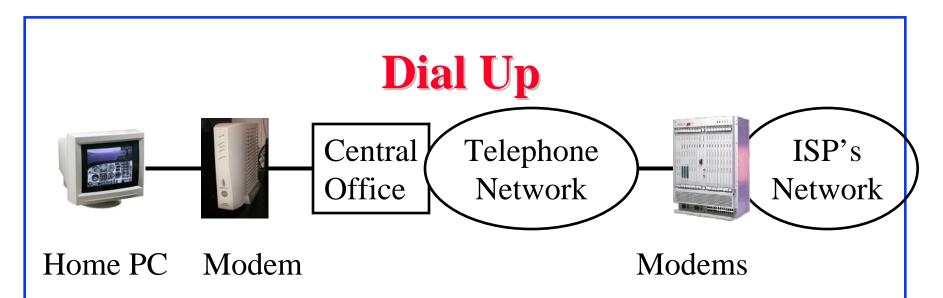


Access Networks

- 1. Dial Up
- DSL 2.
- Cable 3.
- Fiber-To-The-Home 4.
- Wi-Fi 5.
- 6. WiMAX



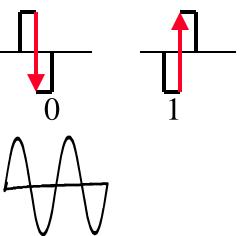




- Modem (Modulator/Demodulator) convert electrical bits to sound waveforms for transmission over telephone network
- □ Telephone network designed to carry 4 kHz voice
- Up to 56 kbps
- Does not need much help from the phone company

Bits, Hertz, and Baud

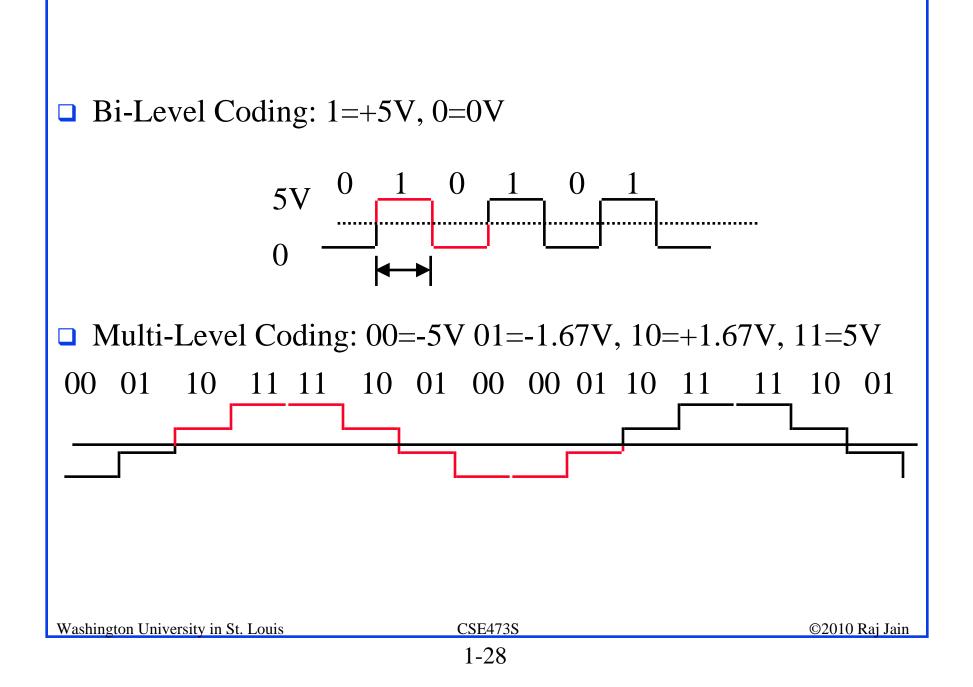
- **Bits:** Unit of information. Binary 0 or 1
- □ Bits are transmitted as pulses: E.g., Manchester encoding

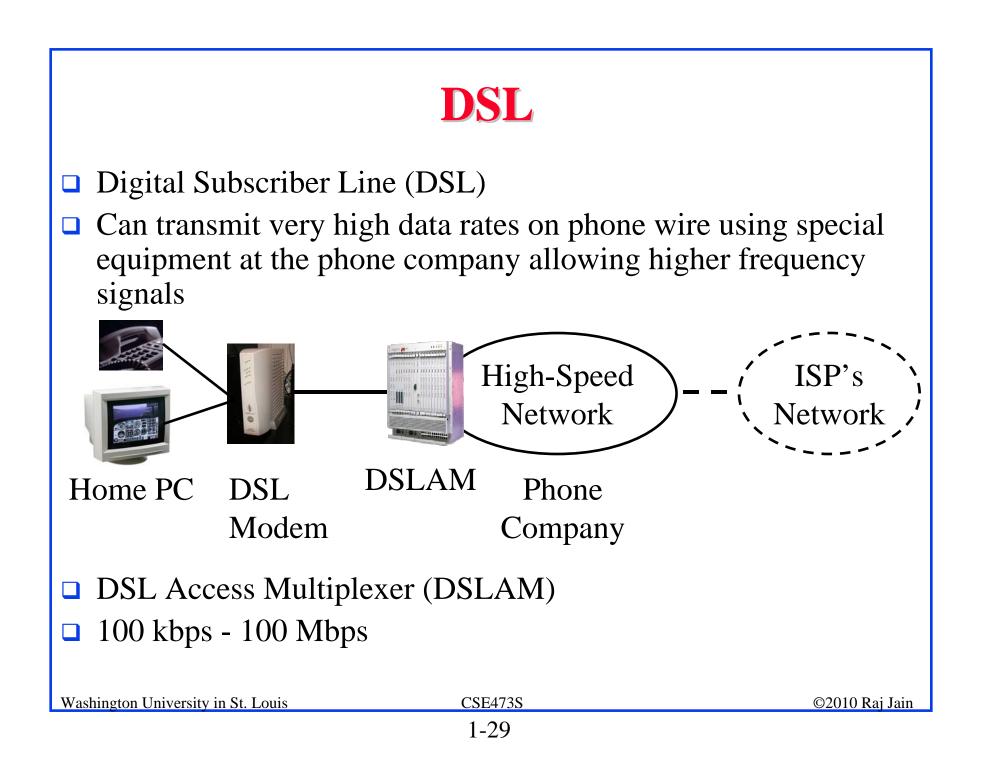


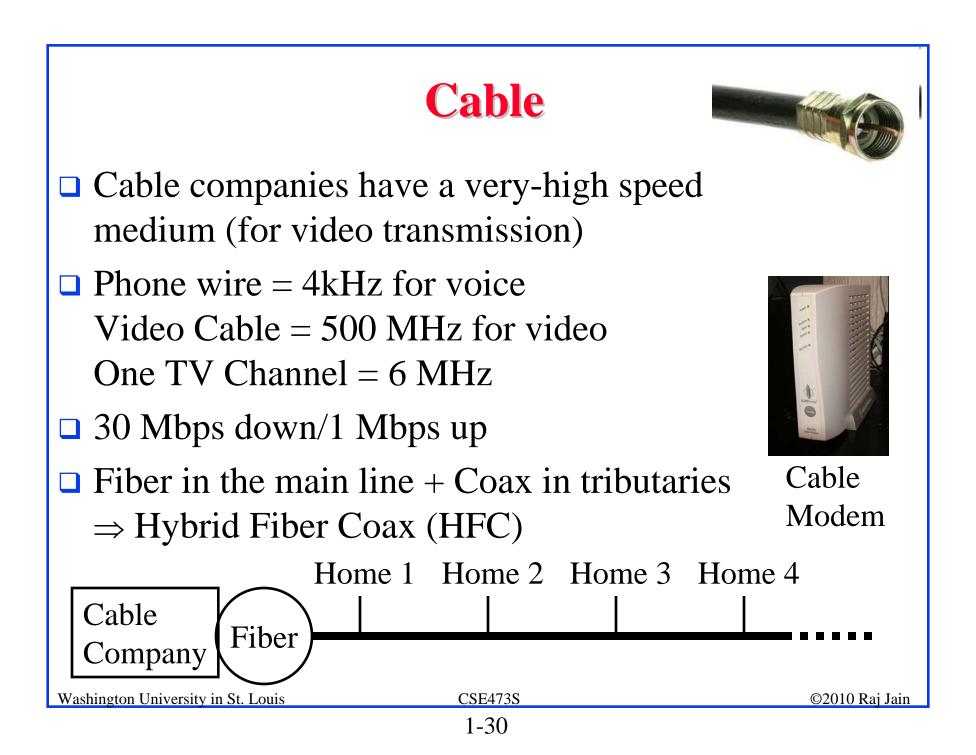
0=High-to-low transition 1=Low-to-high transition

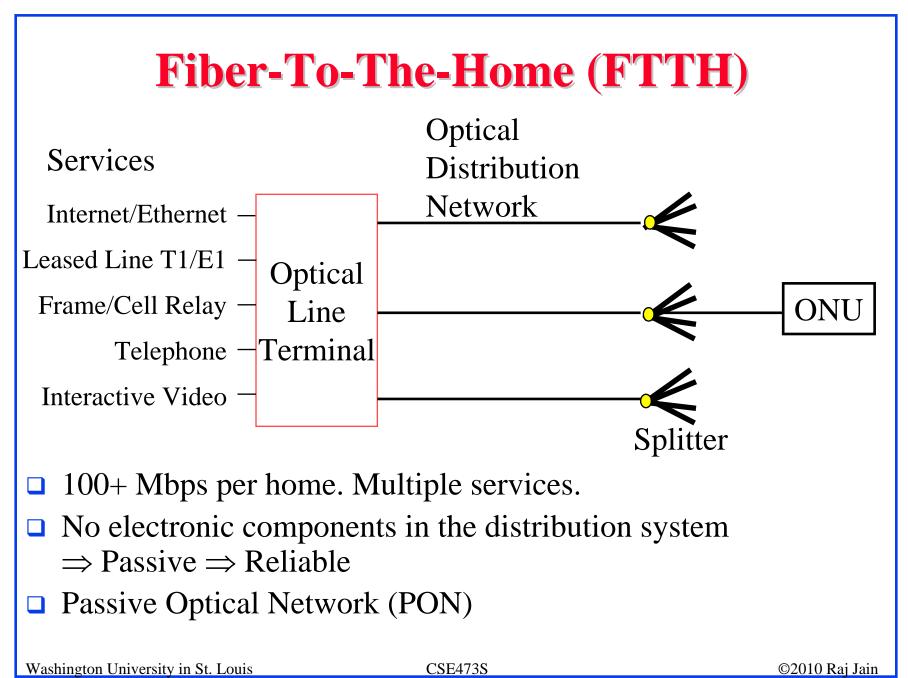
- Receiver design depends on the duration of smallest pulse 1kbps ⇒ One bit per millisecond ⇒ Each pulse is ½ ms ⇒ 2 kBaud
- □ The pulses become a mixture of sine waves on the medium

 $\square Wires allow only certain frequencies \Rightarrow Hertz = cycles/second$ Washington University in St. Louis CSE473S \bigcirc 2010 Raj Jain





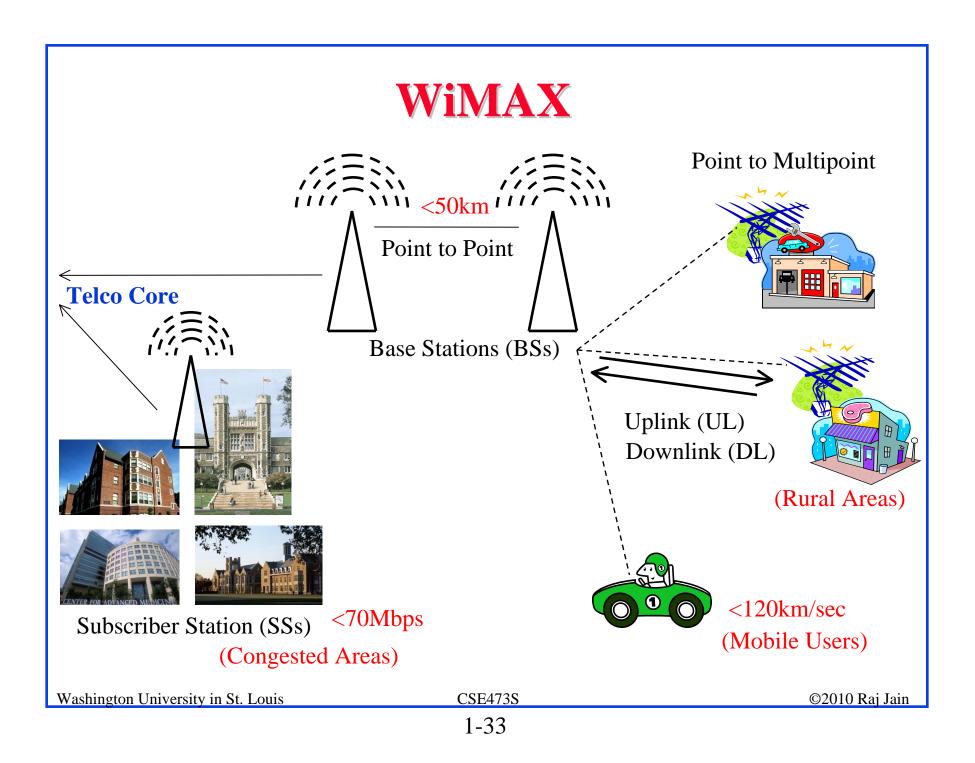




Wireless Access Networks

- □ Wi-Fi hot spots
- Cellular access
- □ WiMAX

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Network Performance Measures

- Delay
- □ Throughput
- Loss Rate

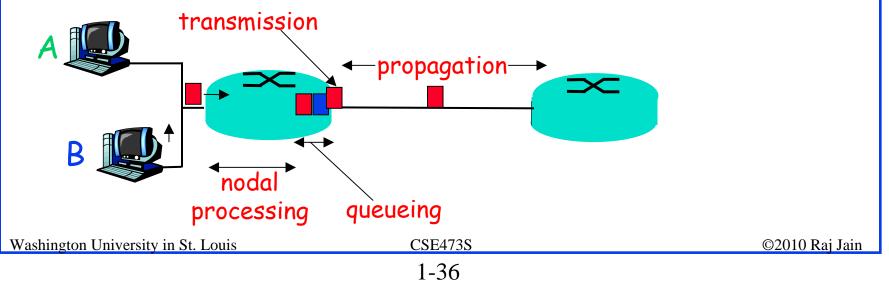
Delay Example (CBR Circuits)

- How long would it take to send a file of 640,000 bits from host A to host B over a circuit-switched network?
 - □ All links are 1.536 Mbps
 - □ Each link is shared by 24 users
 - □ 500 ms to establish end-to-end circuit
- \Box Per User Rate = 1536/24 = 64 kbps
- $\Box \text{ Time to transfer} = 640 \text{kb}/64 \text{kb} = 10 \text{ s}$
- **\Box** Total time = .5 s + 10 s = 10.5 s

Packet Switching Delay

- 1. Processing Delay: Check packets, decide where to send, etc.
- 2. Queuing Delay: Wait behind other packets
- 3. Transmission Delay: First-bit out to last-bit out on the wire= Packet Length/bit rate
- 4. Propagation Delay: Time for a bit to travel from in to out= Distance/speed of signal

Light speed = $3x10^8$ m/s in vacuum, $2x10^8$ m/s in fiber

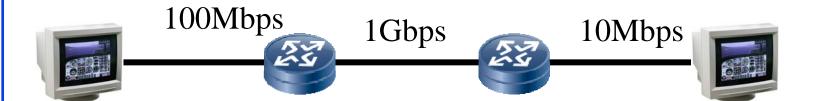


Packet Switching Delay: Example

- □ 1500 Byte packets on 10 Mbps Ethernet, 1km segment
- □ Transmission Delay = $1500x8/10x10^6 = 1200 \ \mu s = 1.2ms$
- □ Propagation delay = $1000 \text{ m}/2 \times 10^8 = 5 \mu \text{s}$

Throughput

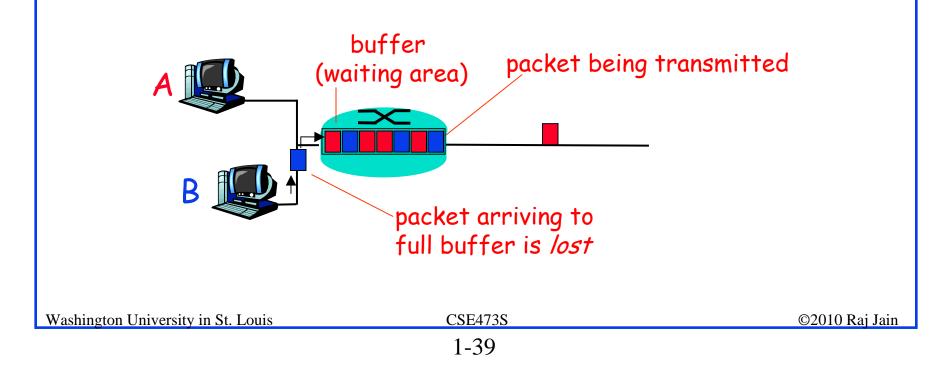
- Measured in Bits/Sec
- **Capacity: Nominal Throughput**
- **Throughput: Realistic**
- □ Bottleneck determines the end-to-end throughput



Net end-to-end capacity = 10 Mbps Actual throughput will be less due to sharing and overhead.

Loss Rate

- $\Box Queuing \Rightarrow Buffer overflow$
- □ Bit Error Rate on the link
- Lost packets are retransmitted by the previous node or the source

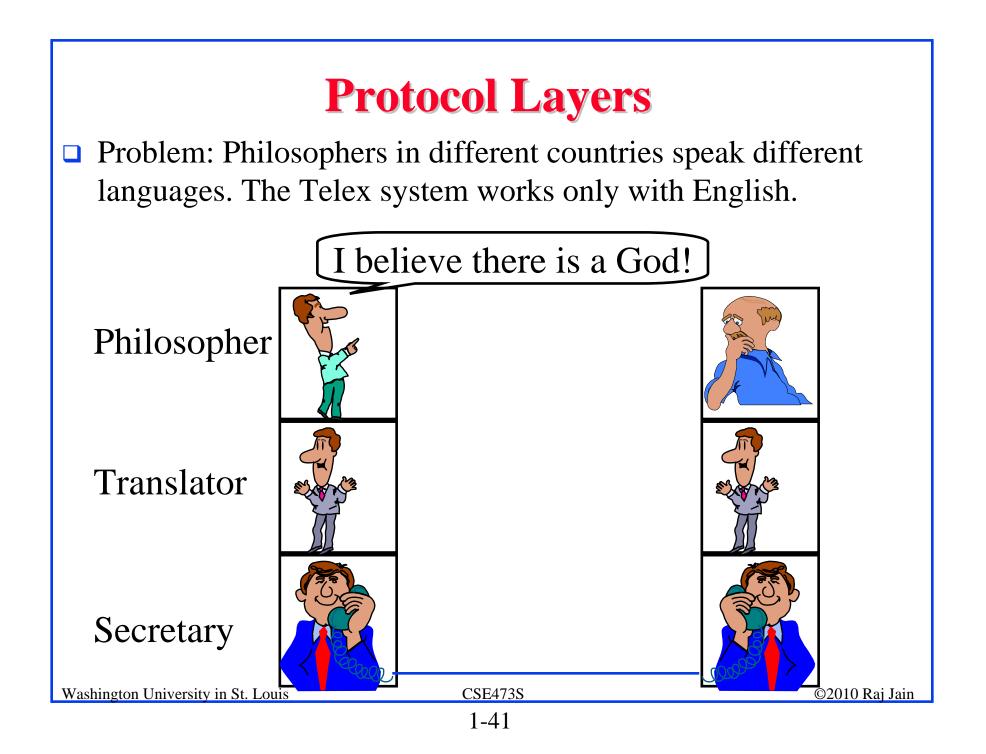


Homework 1B

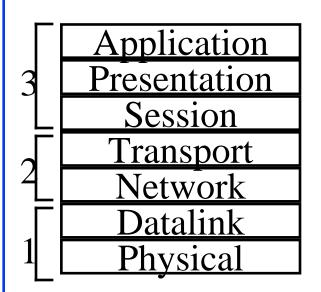
- P5: Consider two hosts, A and B, connected by a single link of rate R bps. Suppose that the two hosts are separated by m meters, and suppose the propagation speed along the link is s meters/sec. Host A is to send a packet of size L bits to Host B.
- A. Express the propagation delay, d_{prop} in terms of *m* and *s*
- B. Determine the transmission time of the packet d_{trans} in terms of L and R.
- C. Ignoring processing queuing delays, obtain an expression for the end-to-end delay
- D. Suppose Host A begins to transmit the packet at time t=0. At time $t=d_{trans}$ where is the last bit of the packet?
- E. Suppose d_{prop} is greater than d_{trans} . At time $t=d_{trans}$, where is the first bit of the packet?
- F. Suppose d_{prop} is less than d_{trans} , at time $t=d_{trans}$, where is the first bit of the packet
- G. Suppose $s=2.5 \times 10^8$ m/s, L=240 bits, and R=56 kbps,. Find the distance *m* so that d_{prop} equals d_{trans} .

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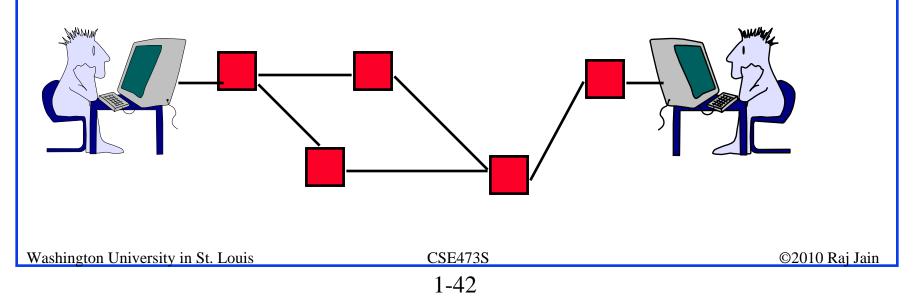
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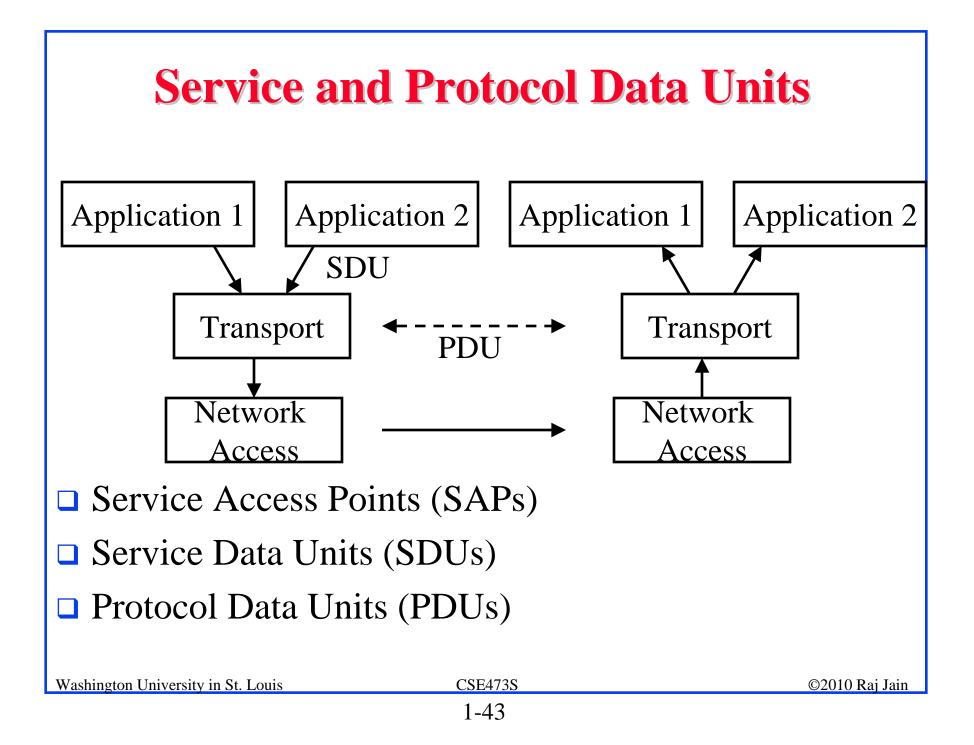


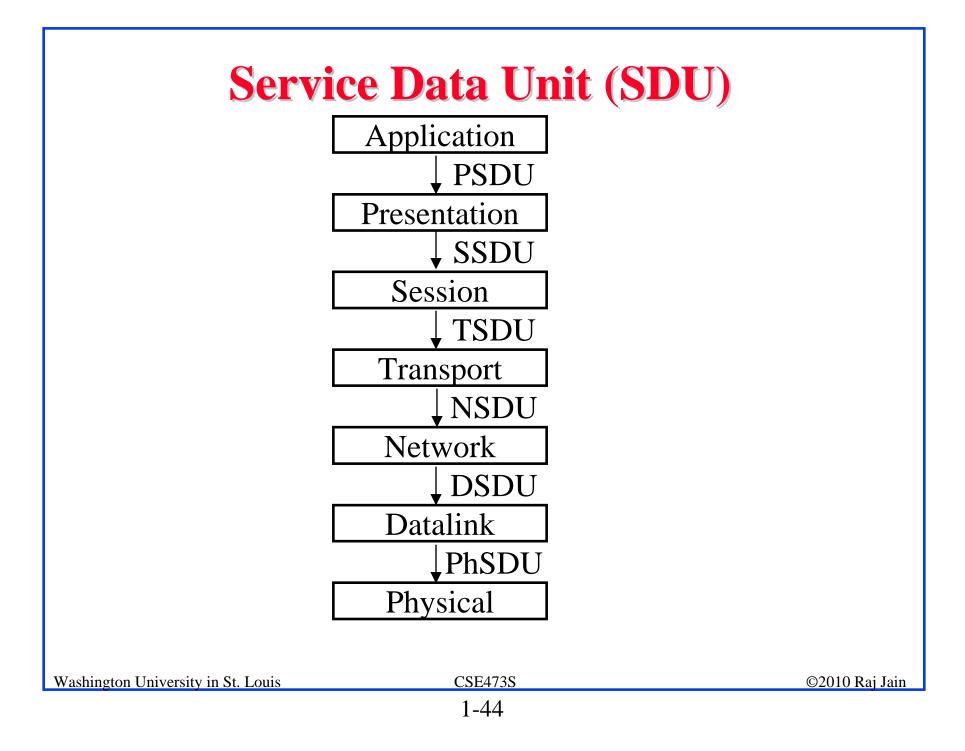
ISO/OSI Reference Model



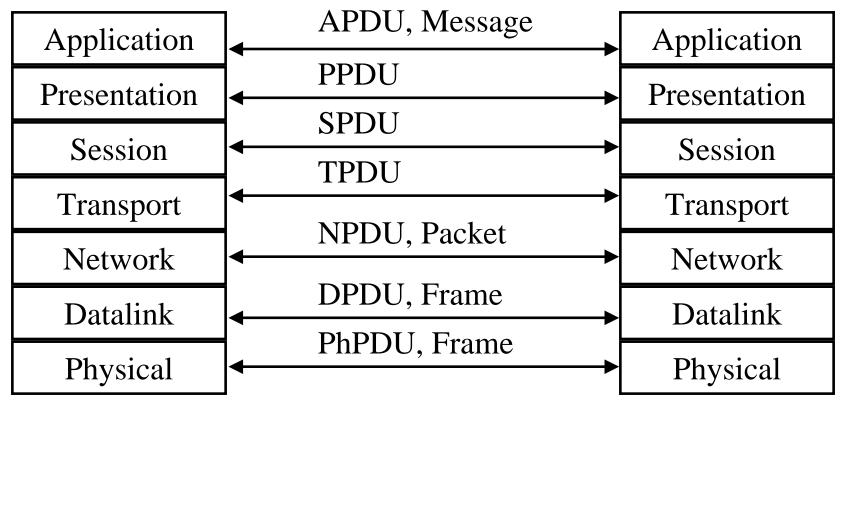
File transfer, Email, Remote Login ASCII Text, Sound Establish/manage connection End-to-end communication: TCP Routing, Addressing: IP Two party communication: Ethernet How to transmit signal: Coding



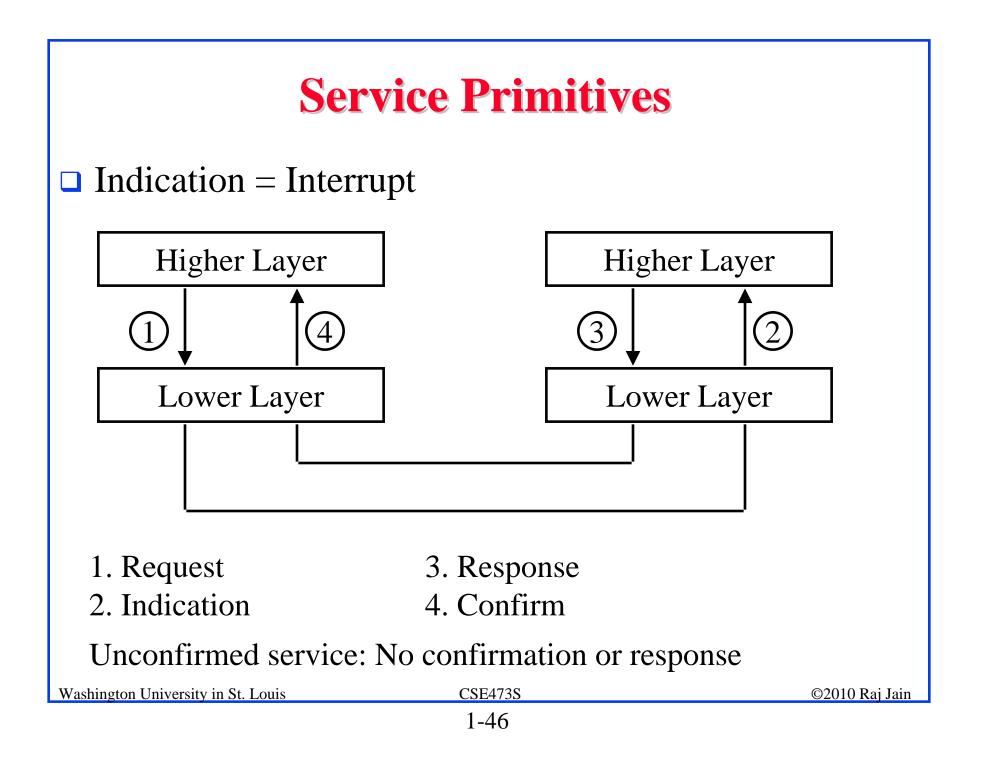


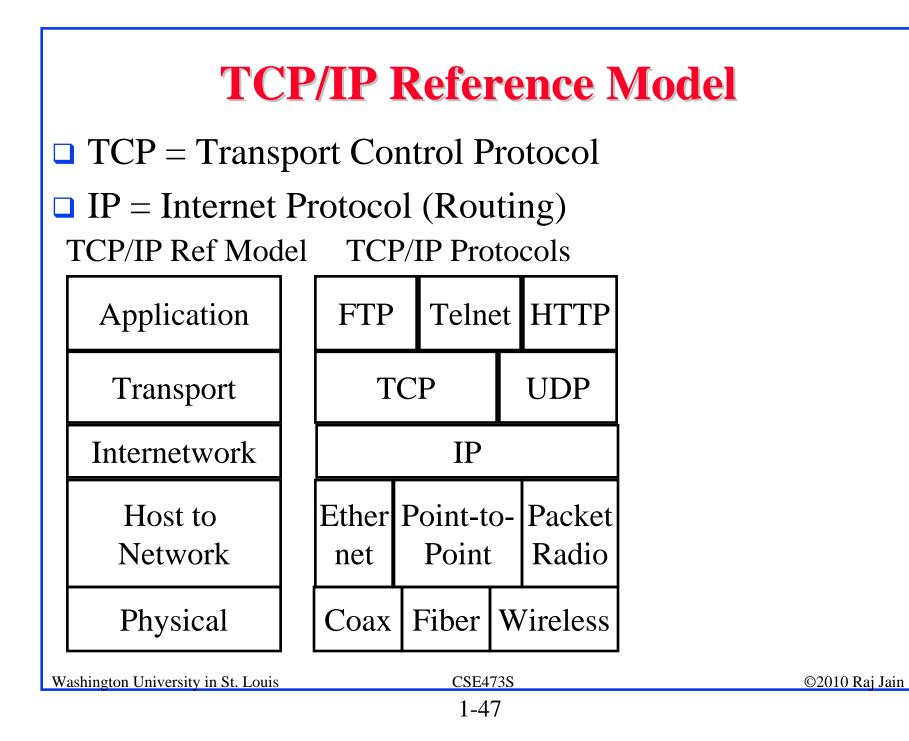


Protocol Data Unit (PDU)



1-45





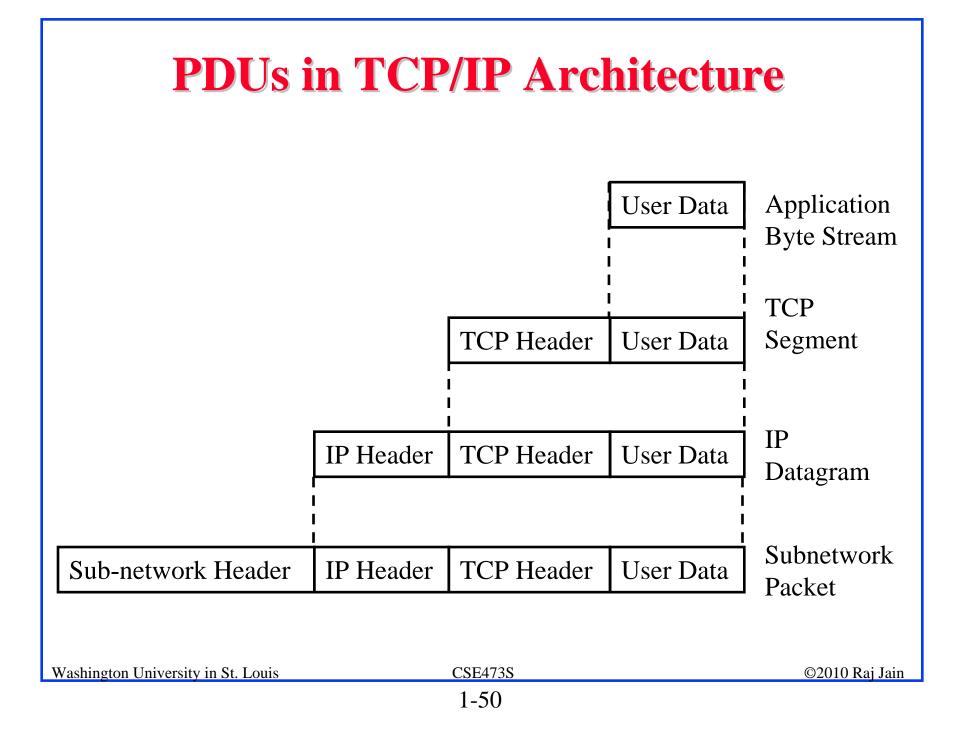
OSI vs TCP/IP

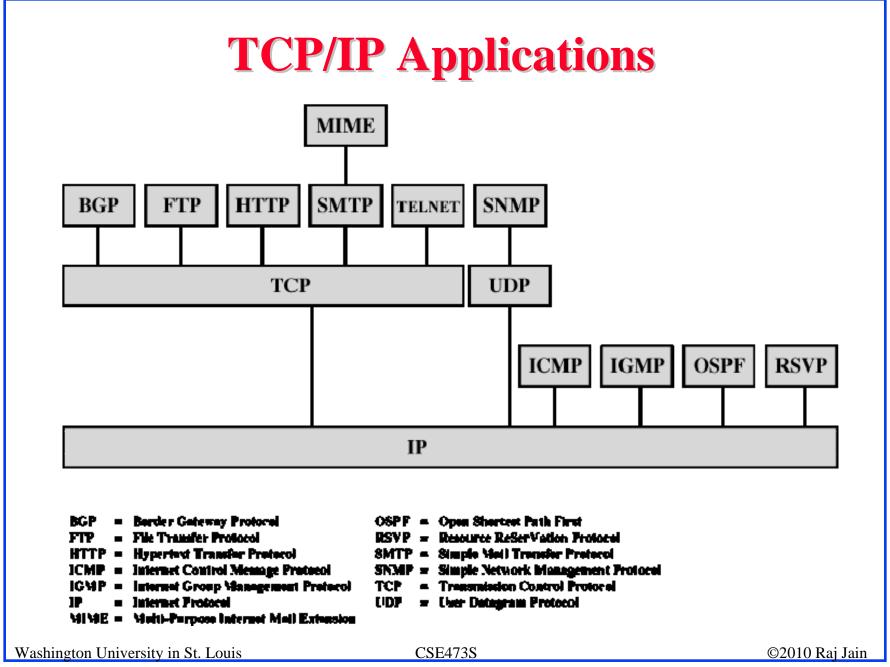
081	тср/пр
Application	
Presentation	Application
Session	
	Transport
Transport	Transport (host-to-host)
Network	Internet
	Network
Data Link	Access
Physical	Physical
CSE473S	

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OSI vs TCP Reference Models

- OSI introduced concept of services, interface, protocols. These were force-fitted to TCP later ⇒ It is not easy to replace protocols in TCP.
- In OSI, reference model was done before protocols.
 In TCP, protocols were done before the model
- OSI: Standardize first, build later TCP: Build first, standardize later
- OSI took too long to standardize. TCP/IP was already in wide use by the time.
- □ OSI became too complex.
- □ TCP/IP is not general. Ad hoc.





Network Security

- Security Components
- **Types of Malware**
- **Types of Attacks**
- Buffer Overflows
- Distributed DoS Attacks

Security Components

- Confidentiality: Need access control, Cryptography, Existence of data
- □ Integrity: No change, content, source, prevention mechanisms, detection mechanisms
- □ Availability: Denial of service attacks,
- □ Confidentiality, Integrity and Availability (CIA)



Types of Malware

- □ Viruses: Code that *attaches* itself to programs, disks, or memory to propagate itself.
- □ Worms: Installs copies of itself on other machines on a network, e.g., by finding user names and passwords
- □ **Trojan horses**: Pretend to be a utility. Convince users to install on PC.
- **Spyware**: Collect personal information
- □ Hoax: Use emotion to propagate, e.g., child's last wish.
- **Trap Door**: Undocumented entry point for debugging purposes
- Logic Bomb: Instructions that trigger on some event in the future
- Zombie: Malicious instructions that can be triggered remotely. The attacks seem to come from other victims.

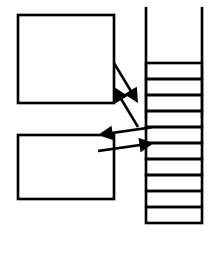
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Types of Attacks

- Denial of Service (DoS): Flooding with traffic/requests
- Buffer Overflows: Error in system programs. Allows hacker to insert his code in to a program.
- □ Malware
- **Brute Force**: Try all passwords.
- □ Port Scanning:
 - \Rightarrow Disable unnecessary services and close ports
- Network Mapping

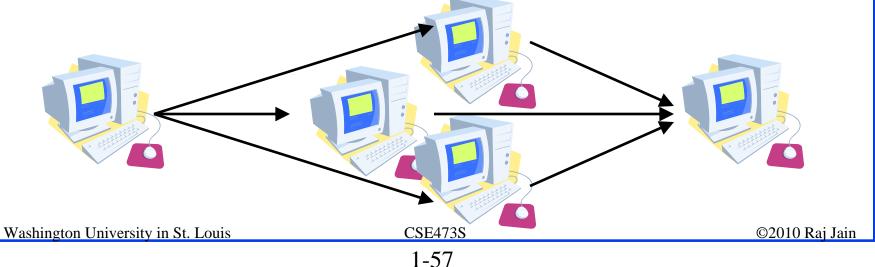
Buffer Overflows

- □ Return address are saved on the top of stack.
- □ Parameters are then saved on the stack.
- □ Writing data on stack causes stack overflow.
- Return the program control to a code segment written by the hacker.



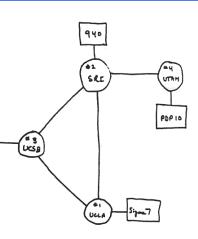
Distributed DoS Attacks

- Tribe Flood Network (TFN) clients are installed on compromised hosts.
- All clients start a simultaneous DoS attack on a victim on a trigger from the attacker.
- □ **Trinoo** attack works similarly. Use UDP packets. Trinoo client report to Trinoo master when the system comes up.
- Stacheldraht uses handlers on compromised hosts to receive encrypted commands from the attacker.



History of Internet

1961: Kleinrock developed queueing theory. Showed effectiveness of packet-switching

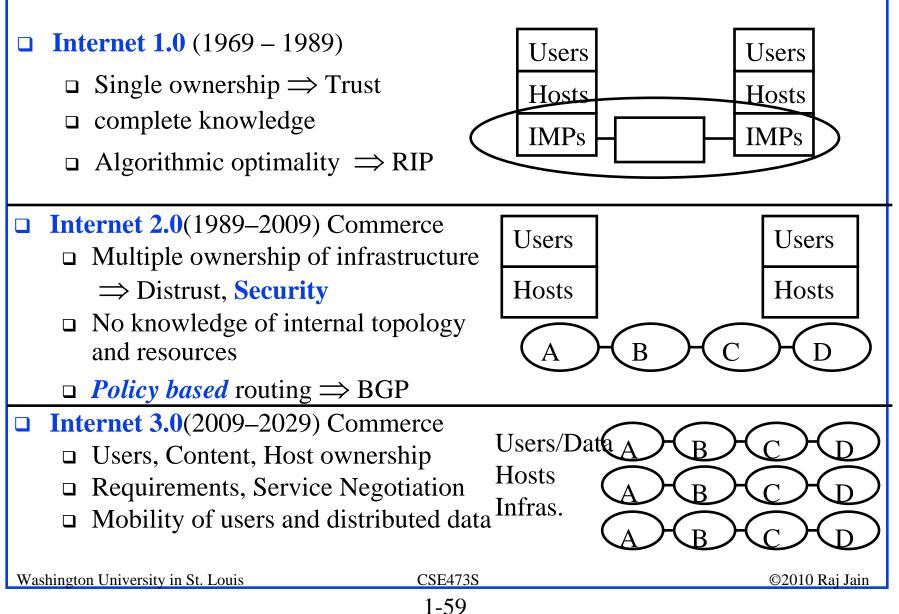


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THE ARPA NETWORK

- 1964: Baran's report on packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- Image: 1969: First ARPAnet node operational First Request for Comment (RFC) <u>www.ietf.org</u>

Internet Generations



History of Internet (Cont)

- □ Early 1990s: HTML, HTTP: Berners-Lee
- □ 1994: Mosaic, later Netscape

2007:

- □ ~500 million hosts
- □ Voice, Video over IP
- P2P applications: BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- Video applications: YouTube, gamingWireless, Mobility

Key Concepts

- □ Internet Protocol (IP): Protocol
- Address: All systems have an IP address, for example, 125.36.47.23
- Name: All systems have a human readable name, e.g., scorpio.cec.wustl.edu, ibm.com.
- Technically called DNS (domain name systems) name. Details will be introduced later.
- IETF: Internet Engineering Task Force. Make standards for Internet. IETF.org
- RFC: Request for comments. Documents that describe Internet protocols.

Homework 1C

- 1. Find the IP address of your computer
- 2. Find the IP address of <u>www.google.com</u> (different from google.com)
- 3. Measure delay from your computer to <u>www.google.com</u>

For all cases submit the screen snapshot showing the command used and the output. (Use Alt-Print-screen to capture a window to clipboard and then paste to word)



- 1. Most common medium is UTP, wireless, fiber
- 2. Internet is a network of networks
- 3. Enterprise, access, and core networks
- 4. Performance Measures: Delay, Throughput, Loss Rate
- 5. Protocol Layers: ISO and TCP/IP reference models



Solution to Homework 1A

- Which networking media will you use for the following applications and why?
- 1. Very large file transfer at home: CAT-6 UTP for gigabit Ethernet
- 2. High-speed multiple channel video transmission at office: Fiber or Coax
- 3. News reading while traveling in a car: wireless

Solution to Homework 1B

- P5: Consider two hosts, A and B, connected by a single link of rate R bps. Suppose that the two hosts are separated by m meters, and suppose the propagation speed along the link is s meters/sec. Host A is to send a packet of size L bits to Host B.
- A. Express the propagation delay, d_{prop} in terms of *m* and *s*

$$d_{prop} = m/s$$

B. Determine the transmission time of the packet d_{trans} in terms of L and R.

$$d_{trans} = L/R$$

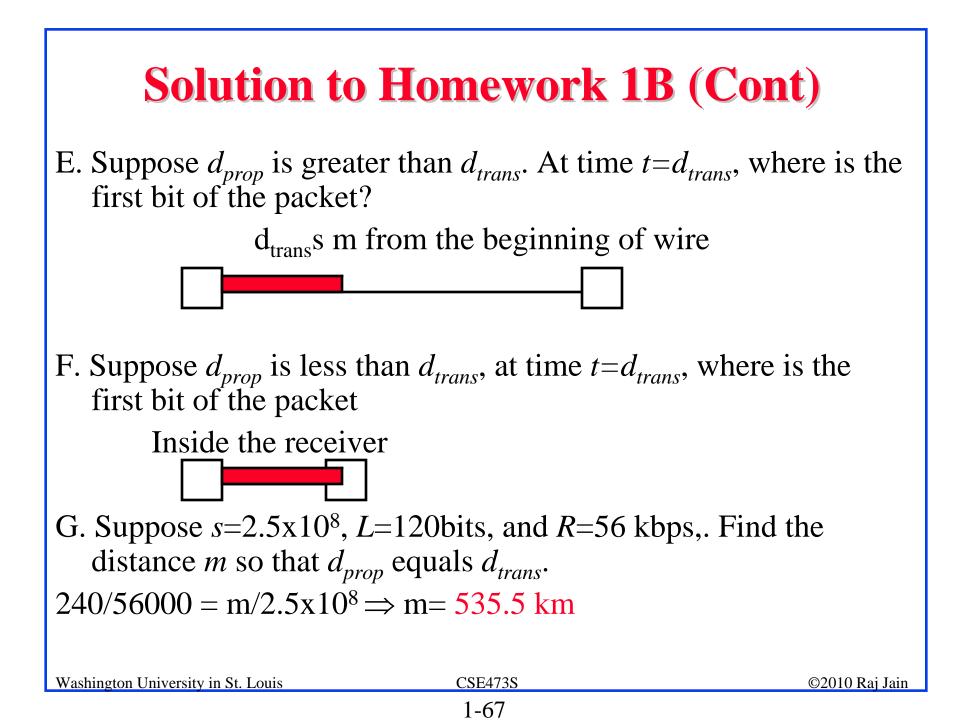
C. Ignoring processing queuing delays, obtain an expression for the end-to-end delay

$$Delay = d_{trans} + d_{prop} = m/s + L/R$$

D. Suppose Host A begins to transmit the packet at time t=0. At time $t=d_{trans}$ where is the last bit of the packet?

Just entering the wire

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Solution to Homework 1C

- □ Find the IP address of your computer
 - □ Ipconfig
 - **u** 192.168.0.108
- □ Find the IP address of <u>www.google.com</u>
 - □ Ping <u>www.google.com</u>
 - **□** 74.125.95.105
- Measure delay from your computer to <u>www.google.com</u>

For all cases show the command used and the output.

Solution to Homework 1C (Cont)

D:\exp>tracert www.google.com

Tracing route to www.l.google.com [74.125.95.105]

over a maximum of 30 hops:

- 1 <1 ms <1 ms <1 ms 192.168.0.1
- 2 9 ms 9 ms 9 ms bras4-10.stlsmo.sbcglobal.net [151.164.182.113]
- 3 8 ms 7 ms 7 ms dist2-vlan50.stlsmo.sbcglobal.net [151.164.14.131]
- 4 7 ms 7 ms 7 ms 151.164.93.224
- 5 15 ms 15 ms 15 ms 151.164.241.189
- 6 21 ms 21 ms 22 ms 72.14.197.85
- 7 21 ms 22 ms 21 ms 209.85.254.120
- 8 32 ms 26 ms 42 ms 209.85.241.22
- 9 32 ms 26 ms 26 ms 209.85.241.35
- 10 28 ms 34 ms 37 ms 72.14.239.193
- 11 32 ms 32 ms 33 ms iw-in-f105.google.com [74.125.95.105]

Trace complete.

