Data Link Control Protocols

Raj Jain Washington University Saint Louis, MO 63131 Jain@cse.wustl.edu

These slides are available on-line at:

http://www.cse.wustl.edu/~jain/cse473-05/

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- **q** Flow Control
- q Effect of propagation delay, speed, frame size
- q Error Recovery
- q HDLC

Flow Control



- **q** Flow Control Goals:
 - 1. Sender does not flood the receiver,
 - 2. Maximize throughput
- q Sender throttled until receiver grants permission

Space-Time Diagrams



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Utilization: Examples

- q Satellite Link: Propagation Delay $t_{prop} = 270 \text{ ms}$ Frame Size = 500 Bytes = 4 kb Data rate = 56 kbps $\Rightarrow t_{frame} = 4/56 = 71 \text{ ms}$ $\alpha = t_{prop}/t_{frame} = 270/71 = 3.8$ $U = 1/(2\alpha+1) = 0.12$
- q Short Link: 1 km = 5 μ s, Rate=10 Mbps, Frame=500 bytes $\Rightarrow t_{\text{frame}} = 4k/10M = 400 \ \mu$ s $\alpha = t_{\text{prop}}/t_{\text{frame}} = 5/400 = 0.012 \Rightarrow U = 1/(2\alpha + 1) = 0.98$

Note: The textbook uses B for t_{prop} and L for t_{frame}

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Efficiency Principle

q For all protocols, the maximum utilization (efficiency) is a *non-increasing* function of α .



Sliding Window Protocols

- **q** Window = Set of sequence numbers to send/receive
- q Sender window
 - q Sender window increases when ack received
 - Packets in sender window must be buffered at source

Sliding Window Diagram



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Error Control

- q Error Control = Deliver frames without error, in the proper
 order to network layer
- q Error control Mechanisms:
 - q Ack/Nack: Provide sender some feedback about other end
 - q Time-out: for the case when entire packet or ack is lost
 - q Sequence numbers: to distinguish retransmissions from originals







Go-back-N (Cont)

All possible scenarios are handled:

- 1. Damaged Frame:
 - q Frame received with error
 - q Frame lost
 - q Last frame lost
- 2. Damaged Ack:
 - q One ack lost, next one makes it
 - q All acks lost
- 3. Damaged Nack:
- q Maximum Window = $2^n 1$ with *n*-bit sequence numbers





Performance: Maximum Utilization

- **q** Stop and Wait Flow Control: $U = 1/(1+2\alpha)$
- q Window Flow Control:

$$U = \begin{cases} 1 & W \ge 2\alpha + 1 \\ W/(2\alpha + 1) & W < 2\alpha + 1 \end{cases}$$

- q Stop and Wait ARQ: $U = (1-P)/(1+2\alpha)$
- q Go-back-N ARQ:

$$U = \begin{cases} (1-P)/(1+2\alpha P) & W \ge 2\alpha+1 \\ W(1-P)/[(2\alpha+1)(1-P+wP)] & W < 2\alpha+1 \end{cases}$$

q Selective Reject ARQ:

$$U = \begin{cases} (1-P) & W \ge 2\alpha + 1 \\ W(1-P)/(2\alpha + 1) & W < 2\alpha + 1 \end{cases}$$



HDLC Family

- q Synchronous Data Link Control (SDLC): IBM
- q High-Level Data Link Control (HDLC): ISO
- q Link Access Procedure-Balanced (LAPB): X.25
- q Link Access Procedure for the D channel (LAPD): ISDN
- q Link Access Procedure for modems (LAPM): V.42
- q Link Access Procedure for half-duplex links (LAPX): Teletex
- q Point-to-Point Protocol (PPP): Internet
- q Logical Link Control (LLC): IEEE
- q Advanced Data Communications Control Procedures (ADCCP): ANSI
- q V.120 and Frame relay also use HDLC





Bit Stuffing

- **q** HDLC Flag = 01111110
- q Every where else in the frame: Replace 11111 with 111110

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Original Pattern
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111111111111011111101111110



1111101111101101111101011111010

7-25



HDLC Frames

Information Frames: User data C q Piggybacked Acks: Next frame expected q Poll/Final = Command/Response **q** Supervisory Frames: Flow and error control q Go back N and Selective Reject q Final \Rightarrow No more data to send **Unnumbered Frames: Control** q q Mode setting commands and responses q Information transfer commands and responses q Recovery commands and responses q Miscellaneous commands and responses

HDLC Commands and Responses Name Function Description

Information (I)	C/R	Exchange user data
Supervisory (S)		
Receive Ready (RR)	C/R	Positive Acknowledgement; ready to receive I-frame
Receive Not Ready (RNR)	C/R	Positive acknowledgement; not ready to receive
Reject (REJ)	C/R	Negative acknowledgement; go back N
Selective Reject (SREJ)	C/R	Negative acknowledgement;
Unnumbered (U)		selective reject
Set Normal Response /	С	Set mode;extended=two-octet control
Extended Mode (SNRM / SNRME)		field
Set Asynchronous Response / Extended Mode (SARM / SARME)	С	Set mode;extended=two-octet control
Set Asynchronous Balanced /	C	Set mode extended=two-octet control
Extended Mode (SABM / SABME)	C	field
Set Initialization Mode (SIM)	С	Initialize link control functions in addressed station
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HDLC Commands and Responses (cont)

Name Function Description

Disconnect (DISC)	С	Terminate logical link connection
Unnumbered Acknowledgement (UA)	R	Acknowledges acceptance of one of the above set-mode commands
Disconnect Mode (DM)	R	Secondary is logically disconnected
Request Disconnect (RD)	R	Request for DISC command
Request Initialization Mode (RIM)	R	Initialization needed; request for SIM command
Unnumbered Information (UI)	C/R	Used to exchange control information
Unnumbered Poll (UP)	С	Used to solicit control information
Reset (RSET)	С	Used for recovery; resets $N(R)$, $N(S)$
Exchange Identification (XID)	C/R	Used to request/report identity and status
Test (TEST)	C/R	Exchange identical information fields for testing
Frame Reject (FRMR)	R	Reports receipt of unacceptable frame
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- q Flow Control: Stop and Wait, Sliding window
- q Effect of propagation delay, speed, frame size
- q Piggybacking
- q Error Control: Stop and wait ARQ, Go-back-N, Selective Reject
- q HDLC: Primary and secondary stations, NRM, ABM, ARM
- q HDLC Frames: Flag, Bit stuffing, I-Frame, RR, RNR

Reading Assignment

- q Read Chapter 7 and Appendix 7A of 7th edition of Stallings.
- q Do the following Exercise from the textbook:7.8 (maximum link utilizations)
- q There is no need to submit the answers.Next Monday is the first mid-term.