Traffic Shaping in ATM Networks

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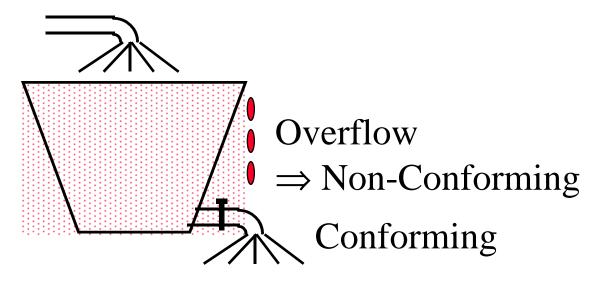


- Leaky bucket
- Generic Cell Rate Algorithm
- □ GCRA Implementations:
 - Virtual Scheduling Algorithm
 - Leaky bucket algorithm
- Examples

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Leaky Bucket



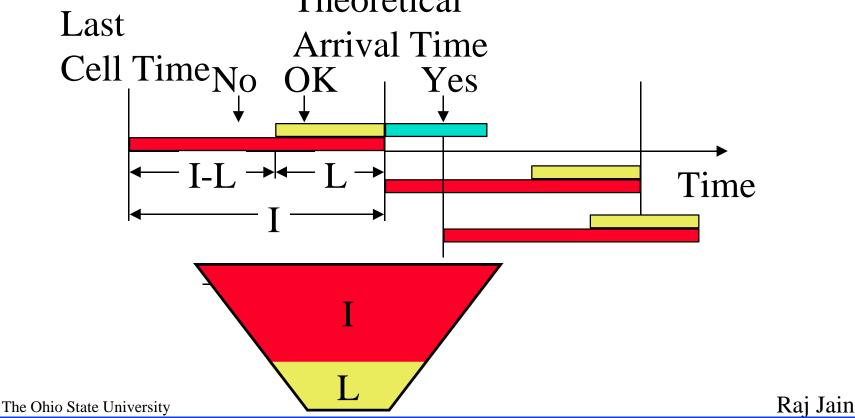
- Provides traffic shaping:Input bursty. Output rate controlled.
- □ Provides traffic policing: Ensure that users are sending traffic within specified limits
 Excess traffic discarded or admitted with CLP = 1

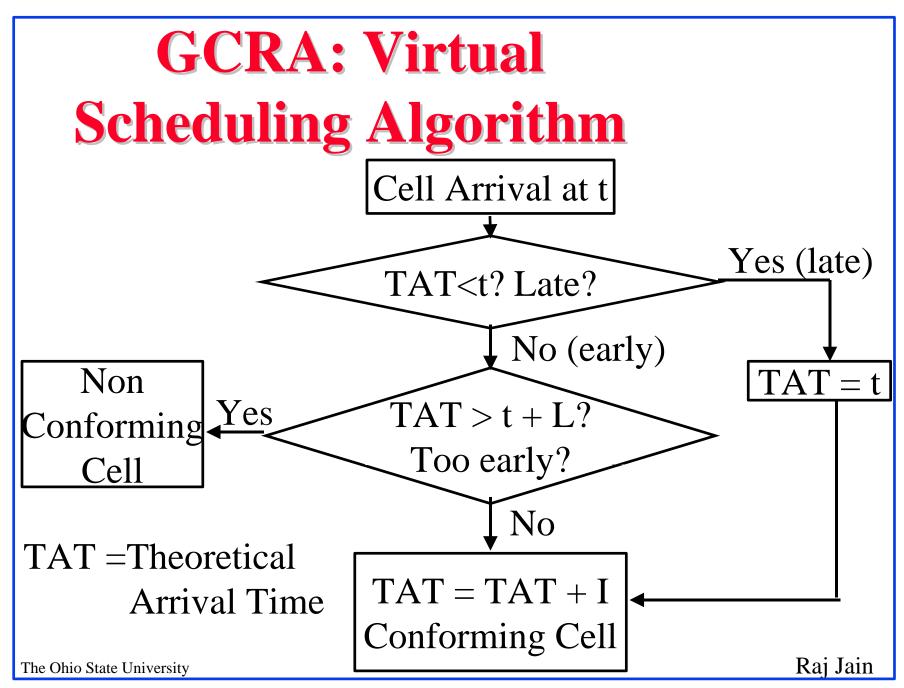
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Generic Cell Rate Algorithm: GCRA(I, L)

- □ I = Increment = Inter-cell Time = Cell size/PCR
- □ $L = Limit \Rightarrow Leaky bucket of size I + L and rate 1$ Theoretical





GCRA: Leaky Bucket Algorithm

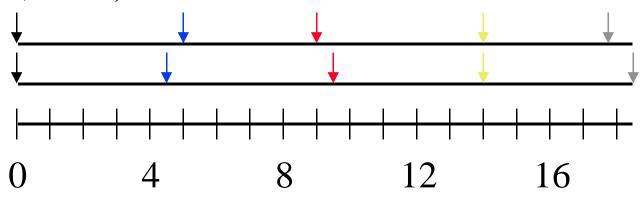
F = X-(t-LCT)Yes F < 0? No Non-Conforming Yes F > L? Cell No LCT = Last Compliance Time X = Bucket contents at LCT X = F + I; LCT = tF = Bucket contents now**Conforming Cell** Raj Jain The Ohio State University

GCRA Examples

 δ = cell time = 2.73 μ s at 155 Mbps

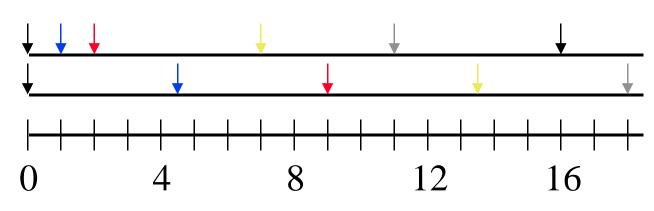
q GCRA(4.5δ , 0.5δ):

Arrivals TAT Time



q GCRA($4.5 \delta, 7 \delta$):

Arrivals
TAT
Time



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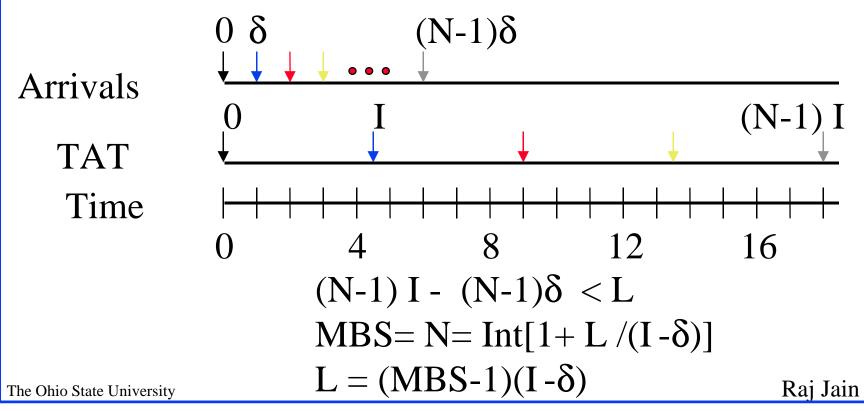
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Maximum Burst Size

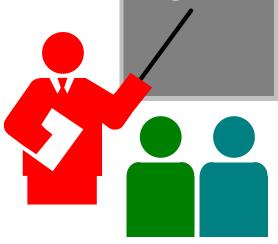
 δ = cell time at PCR, I = cell time at SCR, L=Limit

N = Maximum burst size

GCRA(I, L):



Summary



- □ Leaky bucket is used to smooth bursty arrivals
- □ GCRA requires increment (inter-cell arrival time) and limit (on earlyness)
- □ Two implementations: Virtual scheduling and leaky bucket

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Homework

- Read pages 240-243 of Black's Emerging
 Technologies book 2nd edition.
 (Or Read pages 505-513 of Stallings' ISDN and Broadband ISDN with Frame Relay and ATM)
- Conduct Lab exercise 1

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