# **IP Next Generation (IPv6)** Raj Jain **Raj Jain is now at** Washington University in Saint Louis Jain@cse.wustl.edu http://www.cse.wustl.edu/~jain/ is677-98/ Raj Jain The Ohio State University



- □ Limitations of current Internet Protocol (IP)
- □ How many addresses do we need?
- □ IPv6 Addressing
- □ IPv6 header format

#### **IP Addresses**

- Example: 164.107.134.5
  = 1010 0100 : 0110 1011 : 1000 0110 : 0000 0101
  = A4:6B:86:05 (32 bits)
- □ Maximum number of address =  $2^{32} = 4$  Billion
- Class A Networks: 15 Million nodes
- □ Class B Networks: 64,000 nodes or less
- Class C Networks: 254 nodes or less

#### **IP Address Format**

- □ Three all-zero network numbers are reserved
- 127 Class A + 16,381 Class B + 2,097,151 Class C networks = 2,113,659 networks total
- □ Class B is most popular.
- □ 20% of Class B were assigned by 7/90 and doubling every 14 months ⇒ Will exhaust by 3/94
- Question: Estimate how big will you become? Answer: More than 256! Class C is too small. Class B is just right.

## **How Many Addresses?**

- □ 10 Billion people by 2020
- □ Each person will be served by more than one computer
- Assuming 100 computers per person  $\Rightarrow 10^{12}$  computers
- □ More addresses may be required since
  - Multiple interfaces per node
  - In Multiple addresses per interface
- □ Some believe  $2^6$  to  $2^8$  addresses per host
- □ Safety margin  $\Rightarrow 10^{15}$  addresses
- □ IPng Requirements  $\Rightarrow 10^{12}$  end systems and  $10^{9}$  networks. Desirable  $10^{12}$  to  $10^{15}$  networks

#### **Address Size**

- $\Box$  H Ratio = log<sub>10</sub>(number of objects)/available bits
- $\square$  2<sup>n</sup> objects with n bits: H-Ratio =  $\log_{10} 2 = 0.30103$
- □ French telephone moved from 8 to 9 digits at  $10^7$  households  $\Rightarrow$  H = 0.26 (assuming 3.3 bits/digit)
- □ US telephone expanded area codes with  $10^8$  subscribers  $\Rightarrow$  H = 0.24
- □ SITA expanded 7-character address at 64k nodes  $\Rightarrow$  H = 0.14 (assuming 5 bits/char)
- □ Physics/space science net stopped at 15000 nodes using 16-bit addresses  $\Rightarrow$  H = 0.26
- □ 3 Million Internet hosts currently using 32-bit addresses  $\Rightarrow$  H = 0.20  $\Rightarrow$  A few more years to go

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#### **IPv6 Addresses**

- □ 128-bit long. Fixed size
- $2^{128} = 3.4 \times 10^{38}$  addresses
  - $\Rightarrow$  665×10<sup>21</sup> addresses per sq. m of earth surface
- □ If assigned at the rate of  $10^{6}/\mu$ s, it would take 20 years
- □ Expected to support  $8 \times 10^{17}$  to  $2 \times 10^{33}$  addresses  $8 \times 10^{17} \Rightarrow 1,564$  address per sq. m
- Allows multiple interfaces per host.
- □ Allows multiple addresses per interface
- □ Allows unicast, multicast, anycast
- Allows provider based, site-local, link-local

#### **Colon-Hex Notation**

- **Dot-Decimal**: 127.23.45.88
- **Colon-Hex:** 
  - FEDC:0000:0000:0000:3243:0000:0000:ABCD
  - □ Can skip leading zeros of each word
  - Can skip <u>one</u> sequence of zero words, e.g., FEDC::3243:0000:0000:ABCD ::3243:0000:0000:ABCD
  - □ Can leave the last 32 bits in dot-decimal, e.g., ::127.23.45.88
  - Can specify a prefix by /length, e.g.,
    2345:BA23:7::/40

#### **Initial IPv6 Prefix Allocation**

Allocation	Prefix	Allogation	Drofin
Allocation		Allocation	
Reserved	0000 0000	Unassigned	101
Unassigned	0000 0001	Unassigned	110
NSAP		Unassigned	
IPX	0000 010	Unassigned	1111 0
Unassigned	0000 011	Unassigned	1111 10
Unassigned	0000 1	Unassigned	1111 110
Unassigned	0001	Unassigned	1111 1110
Unassigned	001	Unassigned	1111 1110 0
Provider-based	010	Link-Local	1111 1110 10
Unassigned	011	Site-Local	1111 1110 11
Geographic	100	Multicast	1111 1111

#### **Local-Use Addresses**

Link Local: Not forwarded outside the link, FE:80::xxx

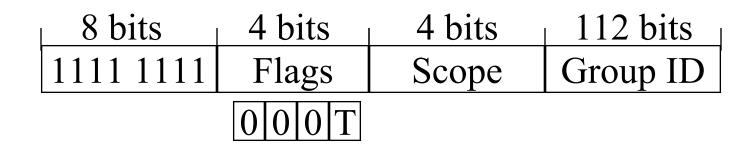
10 bits	n bits	118-n
1111 1110 10	0	Interface ID

Site Local: Not forwarded outside the site, FE:C0::xxx

10 bits	n bits	m bits	118-n-m bits 1
1111 1110 11	0	Subnet ID	Interface ID

Provides plug and play

#### **Multicast Addresses**



- □  $T = 0 \Rightarrow$  Permanent (well-known) multicast address, 1  $\Rightarrow$  Transient
- Scope: 1 Node-local, 2 Link-local, 5 Site-local, 8 Organization-local, E Global
- □ Predefined: 1  $\Rightarrow$  All nodes, 2  $\Rightarrow$  Routers, 1:0  $\Rightarrow$  DHCP servers

#### **Multicast Addresses (Cont)**

□ Example: 43 ⇒ Network Time Protocol Servers
 □ FF01::43 ⇒ All NTP servers on this node
 □ FF02::43 ⇒ All NTP servers on this link
 □ FF05::43 ⇒ All NTP servers in this site
 □ FF08::43 ⇒ All NTP servers in this organization
 □ FF0F::43 ⇒ All NTP servers in the Internet

#### Header

#### □ IPv6:

Version PriorityFlow LabelPayload LengthNext HeaderHop Limit

Source Address

**Destination Address** 

#### $\Box$ IPv4:

Version IHL	Type of Serv	vice	Total Le	ength
		Flags	s Fragment Offset	
Time to Live	Protocol	He	eader Chec	
Source Address				
	Destinatio			
Options Padding			adding	
	*			

#### **Protocol and Header Types**

				-
	Decimal	Keyword	Header Type	_
		HBH	Hop-by-hop (IPv6)	
	1	ICMP	Internet Control Message (IPv4)	
	2	IGMP	Internet Group Management (IPv4)	
	2	ICMP	Internet Control Message (IPv6)	
	3	GGP	Gateway-to-Gateway	
	4	IP	IP in IP (IPv4 Encaptulation)	
	5	ST	Stream	
	6	ТСР		
	17	UDP		
	29	ISO-TP4		
	43	RH	Routing Header (IPv6)	
	44	FS	Fragmentation Header (IPv6)	
	45	IDRP	Interdomain Routing	
	51	AH	Authentication header (IPv6)	
	52	ESP	Encrypted Security Payload	
	59	Null	No next header	
	60	ISO-IP	CLNP	
	88	IGRP		
	89	OSPF	Open Shortest Path First	
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#### IPv6 vs IPv4

- **1995** vs 1975
- □ IPv6 only twice the size of IPv4 header
- Only version number has the same position and meaning as in IPv4
- Removed: header length, type of service, identification, flags, fragment offset, header checksum
- Datagram length replaced by payload length
- Protocol type replaced by next header
- □ Time to live replaced by hop limit
- □ Added: Priority and flow label
- □ All fixed size fields.

- □ No optional fields. Replaced by extension headers.
- □ 8-bit hop limit = 255 hops max (Limits looping)
- □ Next Header = 6 (TCP), 17 (UDP),

#### **Extension Headers**

Base	Extension	Extension	Data
Header	Header 1	Header <i>n</i>	Data

- Most extension headers are examined only at destination
- Routing: Loose or tight source routing
- Fragmentation: All IPv6 routers can carry 536 Byte payload
- Authentication
- □ Security Encaptulation: Confidentiality
- Hop-by-Hop Option
- Destination Options:

## **Extension Header (Cont)**

Only Base Header:

Base Header	ТСР	
Next = TCP	Segment	

Only Base Header and One Extension Header:

Base Header	Route Header	ТСР
Next = TCP	Next = TCP	Segment

□ Only Base Header and Two Extension Headers:

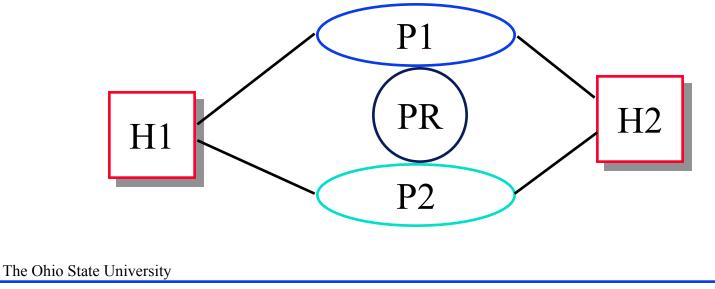
Base Header	Route Header	Auth Header	ТСР
Next = TCP	Next = Auth	Next = TCP	Segment

<b>Routing Header</b>		
Next Header Routing Type Num. Address Next Address		
Reserved  Strict/Loose bit mask		
Address 1		
Address 2		
Address n		
□ Strict $\Rightarrow$ Discard if Address[Next-Address] $\neq$ neighbor □ Type = 0 $\Rightarrow$ Current source routing		

- $\Box \text{ Type} = 0 \Rightarrow \text{Current source routing}$
- □ Type > 0  $\Rightarrow$  Policy based routing (later)
- New Functionality: Provider selection, Host mobility, Auto-readdressing (route to new address)

#### **Provider Selection**

- Possible using routing extension header
- Source specified intermediate systems
- □ No preference: H1, H2
- □ P1 Preferred: H1, P1, H2
- □ H1 becomes Mobile: H1, PR, P1, H2



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□ IPv6 uses 128-bit addresses

- Allows provider-based, site-local, link-local, multicast, anycast addresses
- Fixed header size. Extension headers instead of options. Extension headers for provider selection, security

#### Homework

- □ Read Section 16.5 of Stallings
- □ Submit answer to Exercise 16.19
- Due Date: Next Class