# IP Next Generation (IPv6)

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- □ Limitations of current Internet Protocol (IP)
- □ How many addresses do we need?
- □ IPv6 Addressing
- □ IPv6 header format
- □ IPv6 features

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#### **IP Addresses**

- **Example**: 164.107.134.5
  - $= 1010\ 0100:0110\ 1011:1000\ 0110:0000\ 0101$
  - = A4:6B:86:05 (32 bits)
- Maximum number of address =  $2^{32}$  = 4 Billion
- Class A Networks: 15 Million nodes
- □ Class B Networks: 64,000 nodes or less
- Class C Networks: 250 nodes or less

#### **IP Address**

Class A:

0 Network Local 1 7 24 bits

□ Class B:

10NetworkLocal21416 bits

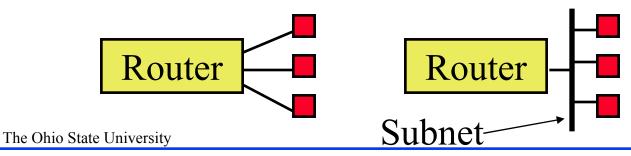
Class C:

110NetworkLocal3218 bits

Class D:

1110 Host Group (Multicast)
4 28 bits

□ Local = Subnet + Host (Variable length)



#### **IP Address Format**

- Three all-zero network numbers are reserved
- □ 127 Class A + 16,381 Class B + 2,097,151 Class C networks = 2,113,659 networks total
- Class B is most popular.
- □ 20% of Class B were assigned by 7/90 and doubling every 14 months ⇒ Will exhaust by 3/94
- Question: Estimate how big will you become?
  Answer: More than 256!
  Class C is too small. Class B is just right.

# **How Many Addresses?**

- □ 10 Billion people by 2020
- □ Each person will be served by more than one computer
- □ Assuming 100 computers per person  $\Rightarrow$  10<sup>12</sup> computers
- More addresses may be required since
  - □ Multiple interfaces per node
  - □ Multiple addresses per interface
- $\Box$  Some believe 2<sup>6</sup> to 2<sup>8</sup> addresses per host
- $\square$  Safety margin  $\Rightarrow 10^{15}$  addresses
- □ IPng Requirements  $\Rightarrow 10^{12}$  end systems and  $10^{9}$  networks. Desirable  $10^{12}$  to  $10^{15}$  networks

#### **Address Size**

- $\square$  H Ratio =  $\log_{10}$ (number of objects)/available bits
- $\square$  2<sup>n</sup> objects with n bits: H-Ratio =  $\log_{10} 2 = 0.30103$
- □ French telephone moved from 8 to 9 digits at  $10^7$  households  $\Rightarrow$  H = 0.26 (assuming 3.3 bits/digit)
- □ US telephone expanded area codes with  $10^8$  subscribers  $\Rightarrow$  H = 0.24
- □ SITA expanded 7-character address at 64k nodes  $\Rightarrow$  H = 0.14 (assuming 5 bits/char)
- □ Physics/space science net stopped at 15000 nodes using 16-bit addresses  $\Rightarrow$  H = 0.26
- □ 3 Million Internet hosts currently using 32-bit addresses  $\Rightarrow$  H = 0.20  $\Rightarrow$  A few more years to go

#### **IPv6 Addresses**

- □ 128-bit long. Fixed size
- □  $2^{128} = 3.4 \times 10^{38}$  addresses ⇒  $665 \times 10^{21}$  addresses per sq. m of earth surface
- $\Box$  If assigned at the rate of  $10^6/\mu s$ , it would take 20 years
- Expected to support  $8 \times 10^{17}$  to  $2 \times 10^{33}$  addresses  $8 \times 10^{17} \Rightarrow 1,564$  address per sq. m
- Allows multiple interfaces per host.
- Allows multiple addresses per interface
- Allows unicast, multicast, anycast
- Allows provider based, site-local, link-local
- 85% of the space is unassigned

#### **Colon-Hex Notation**

- **Dot-Decimal**: 127.23.45.88
- **□** Colon-Hex:

FEDC:0000:0000:0000:3243:0000:0000:ABCD

- □ Can skip leading zeros of each word
- □ Can skip <u>one</u> sequence of zero words, e.g.,

FEDC::3243:0000:0000:ABCD

::3243:0000:0000:ABCD

- □ Can leave the last 32 bits in dot-decimal, e.g.,
  - ::127.23.45.88
- □ Can specify a prefix by /length, e.g.,

2345:BA23:7::/40

#### **Initial IPv6 Prefix Allocation**

Allocation	Prefix	Allocation	Prefix
Reserved	0000 0000	Unassigned	101
Unassigned	0000 0001	Unassigned	110
NSAP	0000 001	Unassigned	1110
IPX		Unassigned	1111 0
Unassigned	0000 011	Unassigned	1111 10
Unassigned	0000 1	Unassigned	1111 110
Unassigned	0001	Unassigned	1111 1110
Unassigned	001	Unassigned	1111 1110 0
Provider-based	010	Link-Local	1111 1110 10
Unassigned	011	Site-Local	1111 1110 11
Geographic	100	Multicast	1111 1111

#### **Local-Use Addresses**

□ Link Local: Not forwarded outside the link,

FE:80::xxx

10 bits	n bits	l 118-n
1111 1110 10	0	Interface ID

□ Site Local: Not forwarded outside the site,

FE:C0::xxx

∟ 10 bits	n bits	m bits	118-n-m bits
1111 1110 11	0	Subnet ID	Interface ID

Provides plug and play

#### **Multicast Addresses**

8 bits	4 bits	4 bits	_ 112 bits _
1111 1111	Flags	Scope	Group ID
	000T		

- □  $T = 0 \Rightarrow$  Permanent (well-known) multicast address,  $1 \Rightarrow$  Transient
- Scope: 1 Node-local, 2 Link-local, 5 Site-local,8 Organization-local, E Global
- □ Predefined:  $1 \Rightarrow \text{All nodes}, 2 \Rightarrow \text{Routers},$  $1:0 \Rightarrow \text{DHCP servers}$

## **Multicast Addresses (Cont)**

- $\square$  Example: 43  $\Rightarrow$  Network Time Protocol Servers
  - $\Box$  FF01::43  $\Rightarrow$  All NTP servers on this node
  - $\Box$  FF02::43  $\Rightarrow$  All NTP servers on this link
  - $\Box$  FF05::43  $\Rightarrow$  All NTP servers in this site
  - $\Box$  FF08::43  $\Rightarrow$  All NTP servers in this organization
  - $\Box$  FF0F::43  $\Rightarrow$  All NTP servers in the Internet

#### Header

□ IPv6:

Version Priority		Flow Labe	el
Payload Let	ngth	Next Header	Hop Limit
- - -	Source	Address	- - -
Destination Address		<del>-</del>	

□ IPv4:

Version IHL   Type of Service		Total Length
Identification	Flags	Fragment Offset
Time to Live Protocol		ader Checksum
Source	e Address	$\mathbf{S}$
Destinati	ion Addre	ess
Options		Padding

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# **Protocol and Header Types**

Decimal	Keyword	Header Type
	HBH	Hop-by-hop (IPv6)
1	ICMP	Internet Control Message (IPv4)
2	IGMP	Internet Group Management (IPv4)
2	ICMP	Internet Control Message (IPv6)
3	GGP	Gateway-to-Gateway
4	IP	IP in IP (IPv4 Encaptulation)
5	ST	Stream
6	TCP	
17	UDP	
29	ISO-TP4	
43	RH	Routing Header (IPv6)
44	FS	Fragmentation Header (IPv6)
45	IDRP	Interdomain Routing
51	AH	Authentication header (IPv6)
52	ESP	Encrypted Security Payload
59	Null	No next header
60	ISO-IP	CLNP
88	IGRP	
89	OSPF	Open Shortest Path First

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#### IPv6 vs IPv4

- □ 1995 vs 1975
- □ IPv6 only twice the size of IPv4 header
- Only version number has the same position and meaning as in IPv4
- Removed: header length, type of service, identification, flags, fragment offset, header checksum
- Datagram length replaced by payload length
- Protocol type replaced by next header
- □ Time to live replaced by hop limit
- □ Added: Priority and flow label
- □ All fixed size fields.

- □ No optional fields. Replaced by extension headers.
- 8-bit hop limit = 255 hops max (Limits looping)
- $\square$  Next Header = 6 (TCP), 17 (UDP),

#### **Extension Headers**

Base	Extension	Extension	Doto
Header	Header 1	Header <i>n</i>	Data

- Most extension headers are examined only at destination
- Routing: Loose or tight source routing
- Fragmentation: All IPv6 routers can carry 536 Byte payload
- Authentication
- Security Encaptulation: Confidentiality
- Hop-by-Hop Option
- Destination Options:

# **Extension Header (Cont)**

Only Base Header:

Base Header TCP
Next = TCP Segment

Only Base Header and One Extension Header:

Base Header Route Header TCP
Next = TCP Next = TCP Segment

Only Base Header and Two Extension Headers:

Base HeaderRoute HeaderAuth HeaderTCPNext = TCPNext = AuthNext = TCPSegment

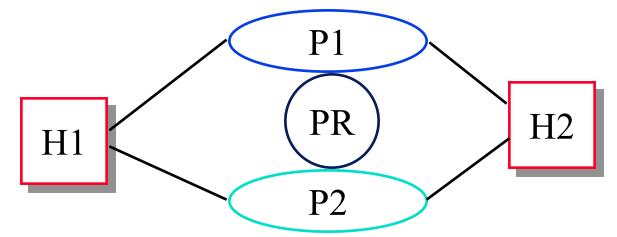
# **Routing Header**

Next Header	Routing Type Num. Address Next Address		
Reserved	Strict/Loose bit mask		
-	Address 1		
Address 2			
	Address n		

- □ Strict ⇒ Discard if Address[Next-Address] ≠ neighbor
- $\square$  Type = 0  $\Rightarrow$  Current source routing
- □ Type  $> 0 \Rightarrow$  Policy based routing (later)
- New Functionality: Provider selection, Host mobility, Auto-readdressing (route to new address)

#### **Provider Selection**

- Possible using routing extension header
- Source specified intermediate systems
- □ No preference: H1, H2
- □ P1 Preferred: H1, P1, H2
- □ H1 becomes Mobile: H1, PR, P1, H2



#### **IPv6 Features**

- Larger addresses
- □ Flexible header format
- Improved options
- Support for resource allocation
- Provision for protocol extension
- Built-in Security:
   Both authentication and confidentiality

# **Address Autoconfiguration**

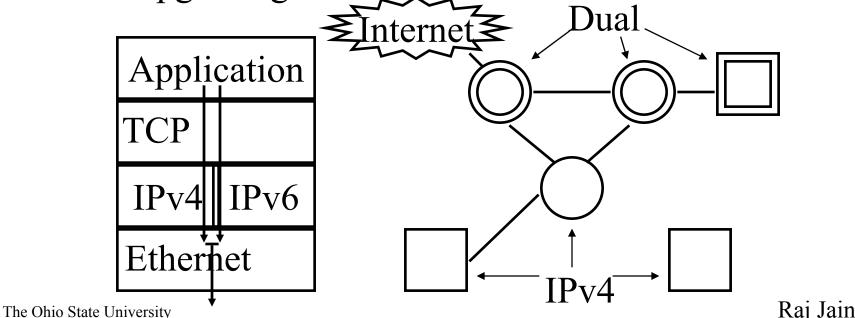
- Allows plug and play
- BOOTP and DHCP are used in IPv4
- □ DHCPng will be used with IPv6
- Two Methods: Stateless and Stateful
- Stateless:
  - □ A system uses link-local address as source and multicasts to "All routers on this link"
  - □ Router replies and provides all the needed prefix info

# **Address Autoconfiguration (Cont)**

- □ All prefixes have a associated lifetime
- System can use link-local address permanently if no router
- Stateful:
  - □ Problem w stateless: Anyone can connect
  - □ Routers ask the new system to go DHCP server (by setting managed configuration bit)
  - □ System multicasts to "All DHCP servers"
  - □ DHCP server assigns an address

#### **Transition Mechanisms**

- Dual-IP Hosts, Routers, Name servers
- Tunneling IPv6 over IPv4
- Hosts and Routers can be gradually upgraded to IPv6
- ☐ It is better (though not required) to upgrade routers before upgrading hosts



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## **Application Issues**

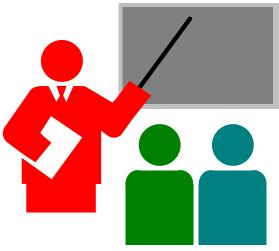
- Most application protocols will have to be upgraded: FTP, SMTP, Telnet, Rlogin
- □ 27 of 51 Full Internet standards, 6 of 20 draft standards, 25 of 130 proposed standards will be revised for IPv6
- No checksum ⇒ checksum at upper layer is mandatory, even in UDP
- non-IETF standards: X-Open, Kerberos, ... will be updated
- □ Should be able to request and receive new DNS records

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# **Implementations**

- □ 4.4-lite BSD by US Naval Research Laboratory
- UNIX, OPEN-VMS by DEC
- DOS/WINDOWS by FTP Software
- □ HP-UX SICS (Swedish Institute of Comp. Science)
- Linux
- NetBSD by INRIA Rocquencourt
- □ Solaris 2 by Sun
- Streams by Mentat

# Summary



- ☐ IPv6 uses 128-bit addresses
- Allows provider-based, site-local, link-local, multicast, anycast addresses
- Fixed header size. Extension headers instead of options.
   Extension headers for provider selection, security
- Allows auto-configuration
- Dual-IP router and host implementations for transition The Ohio State University