

Lecture 12, Review Question 2:

- Assume that a TCP sender A is connected to router X by a 100 Mb/s link, that the corresponding receiver B is connected to router Y by a 100 Mb/s link and that the link connecting X and Y is 10 Mb/s. Also, assume that the roundtrip propagation time between A and B is 50 ms, that the MSS is 1250 bytes and that $ssthresh=64$ KB. Suppose that A starts in the slow-start state. At what rate is A sending after 400 ms?

$N \times 1250\text{kB} \geq 64\text{ kBytes}$
 $N = 52$
We are looking for
 $\text{cwnd} = 52$

Slow Start with $\text{sssthresh} = 64\text{kB}$

$\text{cwnd} = 1\text{ MSS}$

51.2 ms

RTT 1

Time = 50 ms +

0.1 ms

$\text{cwnd} = 2\text{ MSS}$

RTT 2

Time = 100 ms +

1 ms

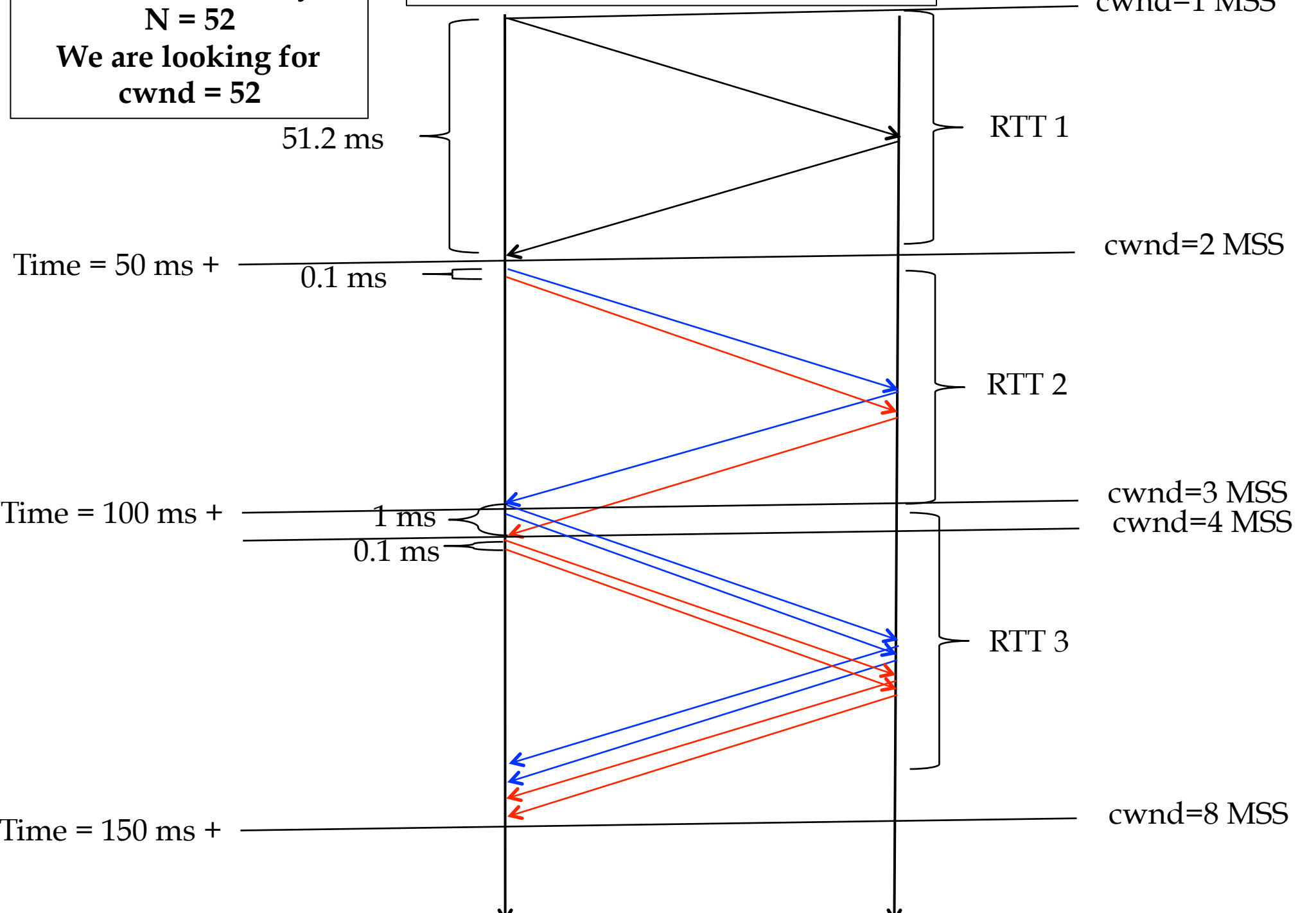
0.1 ms

$\text{cwnd} = 3\text{ MSS}$
 $\text{cwnd} = 4\text{ MSS}$

RTT 3

Time = 150 ms +

$\text{cwnd} = 8\text{ MSS}$



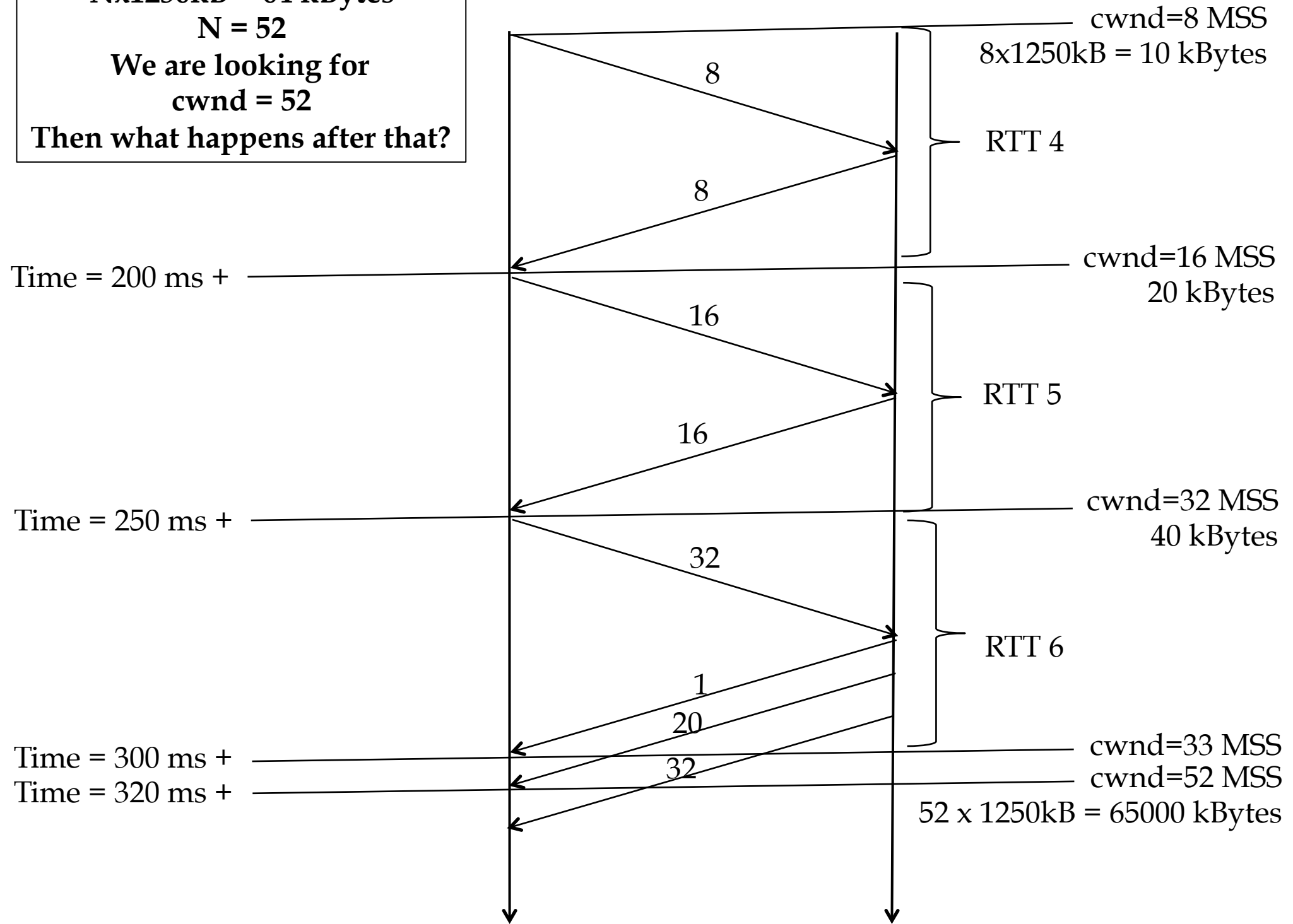
$N \times 1250\text{kB} = 64\text{ kBytes}$

$N = 52$

We are looking for

$\text{cwnd} = 52$

Then what happens after that?



After that...

- What happens once we have reached 64KB?
 - » Congestion control is geared toward controlling the sending rate when it is the network limiting the sender.
 - » What if it is the application limiting it?
 - This is when we don't have enough data from the application on the sending host so that the send buffer empties before we run out of *cwnd* sending space. In this case we are application limited and *cwnd* will not be incremented for each ACK.
 - » What if it is the receive window limiting it?
 - This is when $cwnd \geq rwnd$ and it is no longer *cwnd* that controls the sending rate of the sender. Again, *cwnd* should not continue to be incremented if it is not *cwnd* that limits our sending rate.
 - I'm still looking for a definitive reference for this...
- So, if we assume that the receive window has the default maximum of 64KBytes that *cwnd* will not increment higher than that.
 - » The second part of question turns off that assumption...