

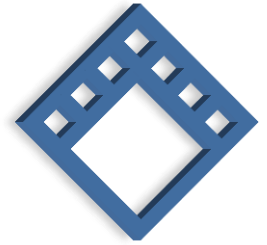
Metastability (What?)

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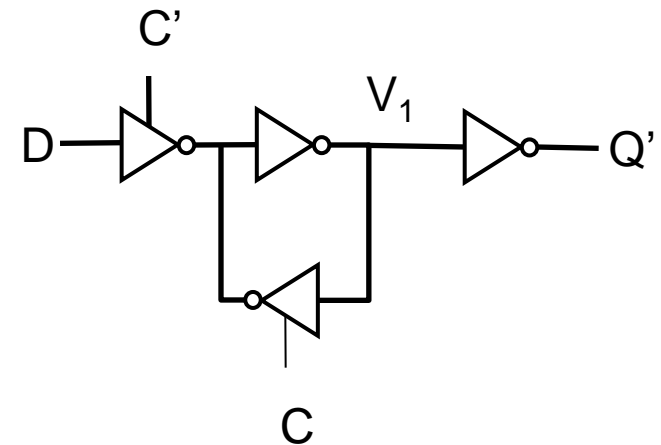
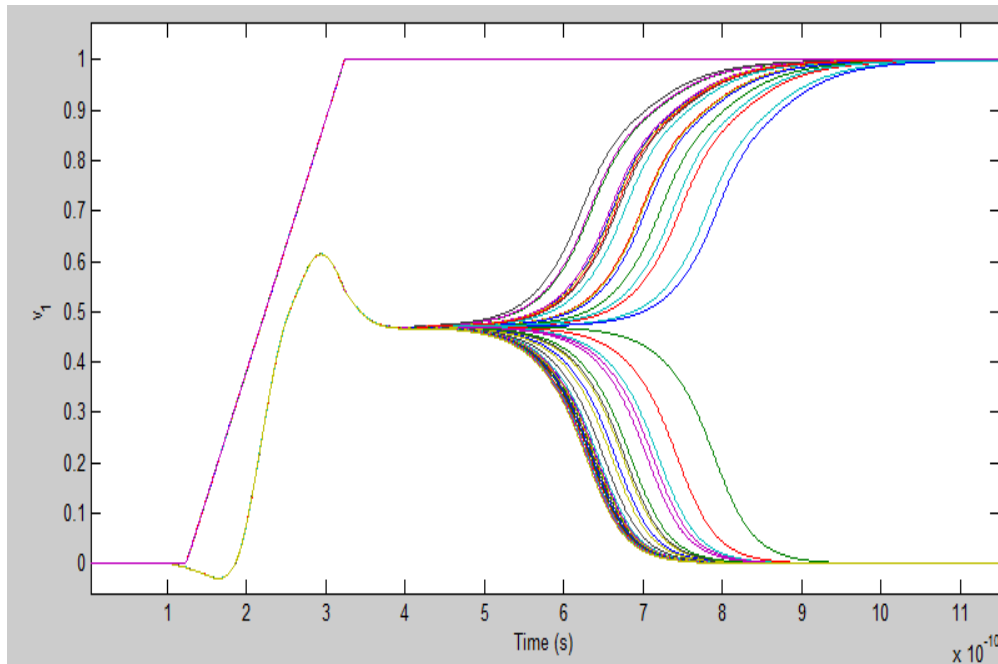
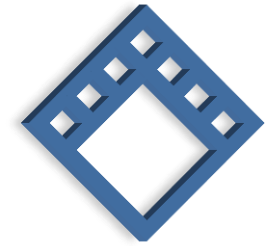
Metastability Is



- a fundamental property of all bi-stable circuits (flip-flops and arbiters)
- the cause of ambiguous output voltages and unpredictable behavior
- the reason for setup & hold-time constraints on flip-flops
 - When observed they eliminate metastability
 - When violated may lead to circuit malfunction
 - Satisfying constraints perfectly between multiple independent clock domains is not possible



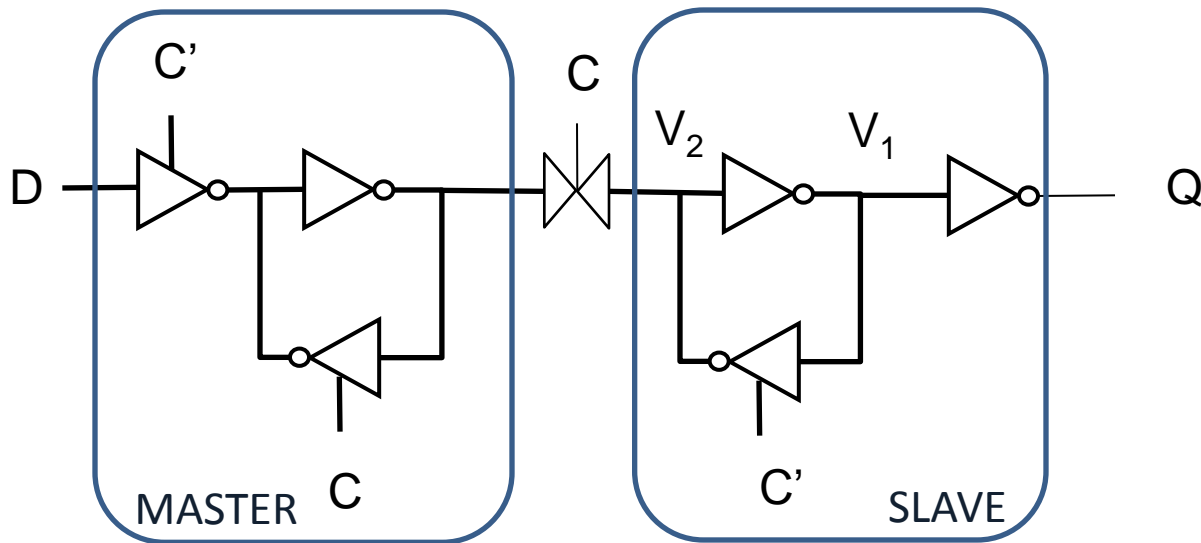
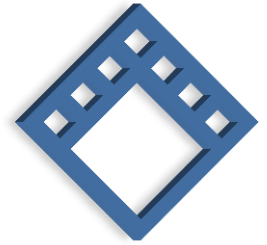
Results for a D-Latch



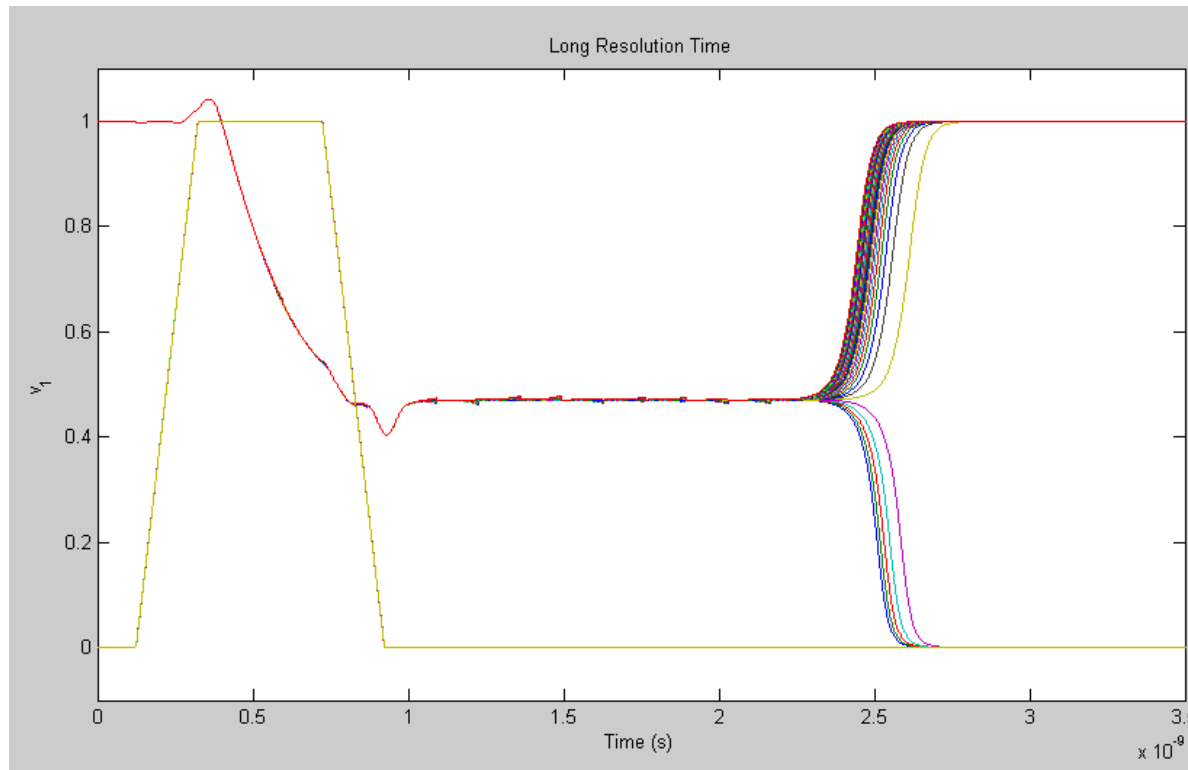
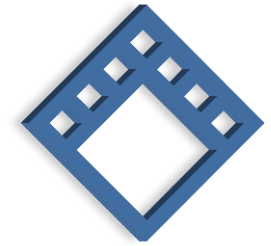
- Latch output before final inverter (clock is also shown).
- Rightmost two traces bracket unbounded metastable point



Prototypical Master-Slave DFF



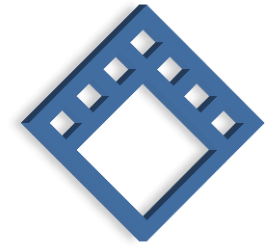
Results for a Master-Slave



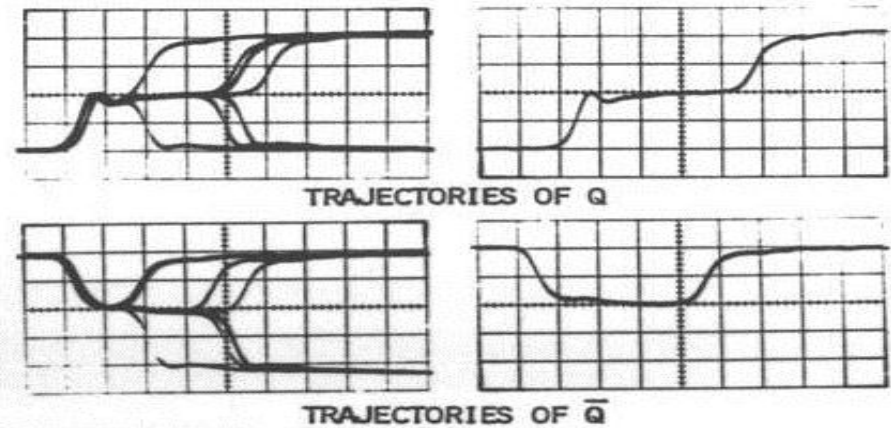
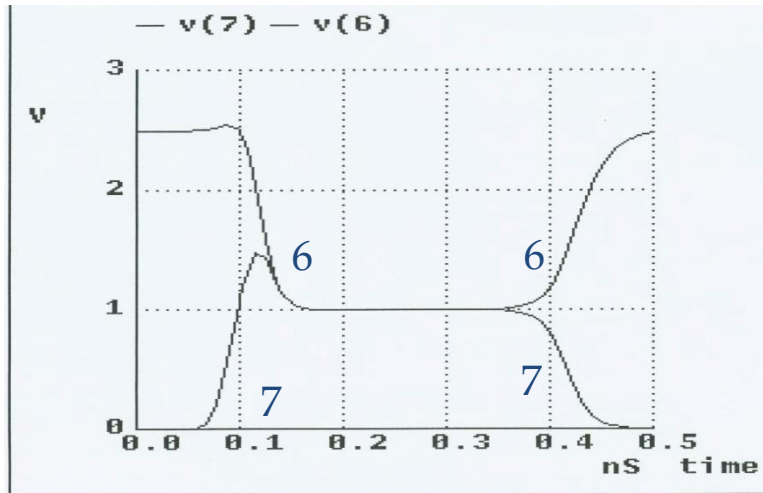
- Clock is shown in yellow.
- Other traces are obtained by varying the data-clock separation and observing the output of the FF before the output inverter.



Real plots!



Times and voltages far from normal experience
And History Dependent! – must collect data slowly



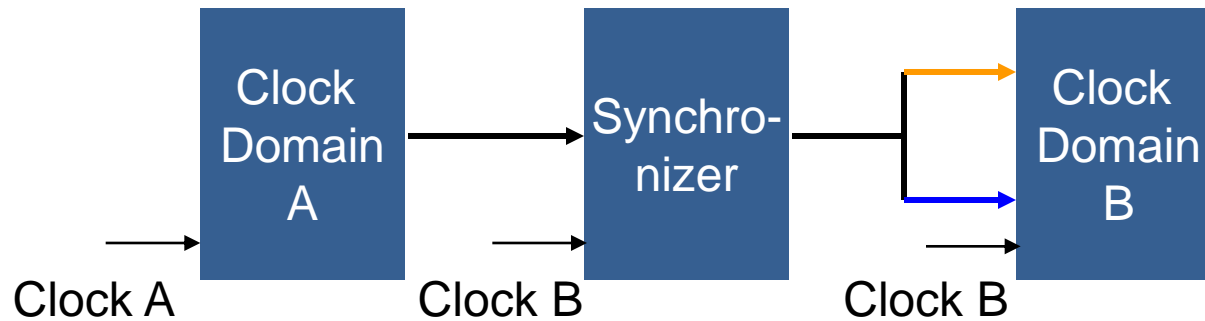
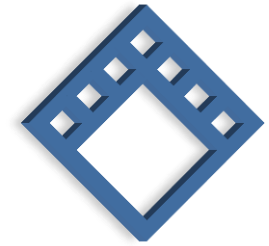
Photos of ECL circuits taken about 45 years ago.

VCLK 1 0 pulse (2.5 0 0.063028851134n 0.5n 0.5n 100n 200n)
 Vdata 2 0 pulse (0 2.5 0n
 ↑ ↑ ↑ ↑
 psec fsec asec zepto sec

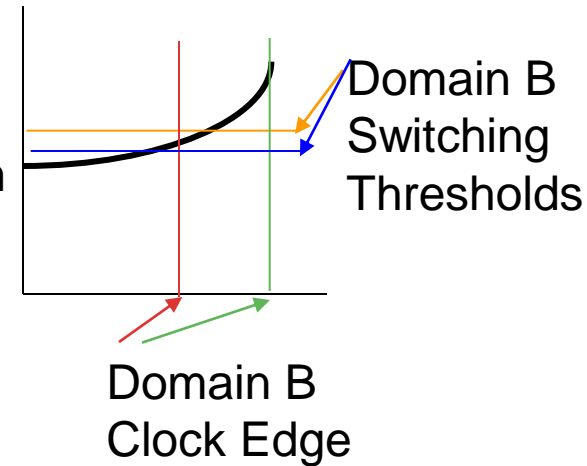
These measured waveforms represent an input timing resolution of about 100 asec.



A Synchronizer Failure



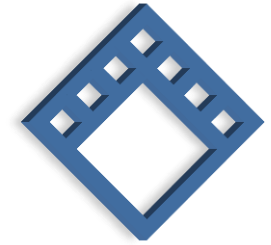
Synchronizer
Output Voltage
Recovering from
Metastability



**Synchronizer fails and future
behavior of Domain B unknown**



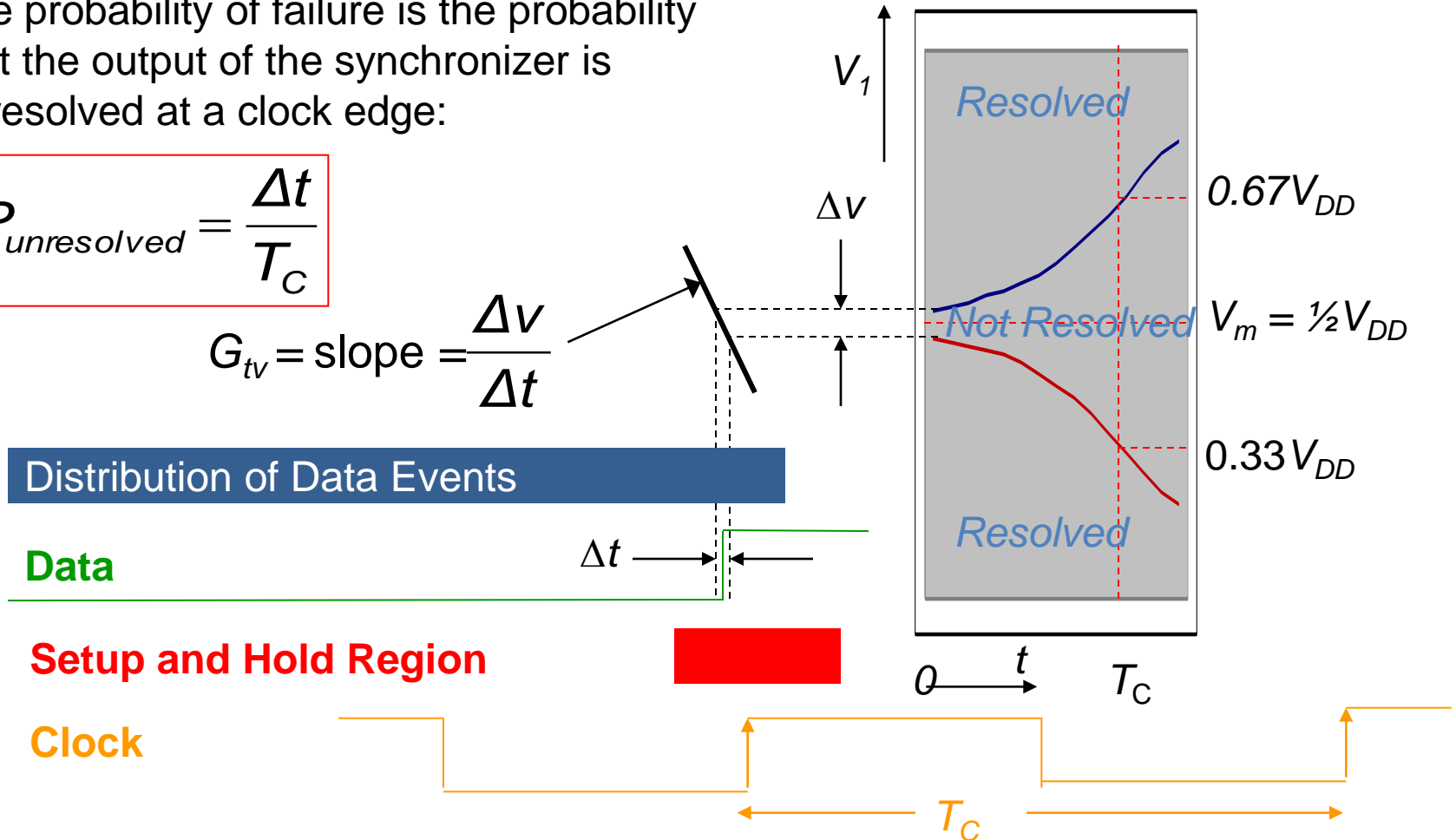
Probability of Synchronizer Failure (Noise Free Case First)



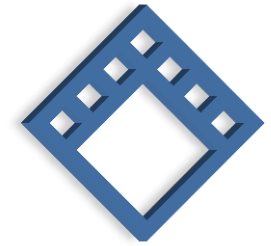
The probability of failure is the probability that the output of the synchronizer is unresolved at a clock edge:

$$P_{unresolved} = \frac{\Delta t}{T_C}$$

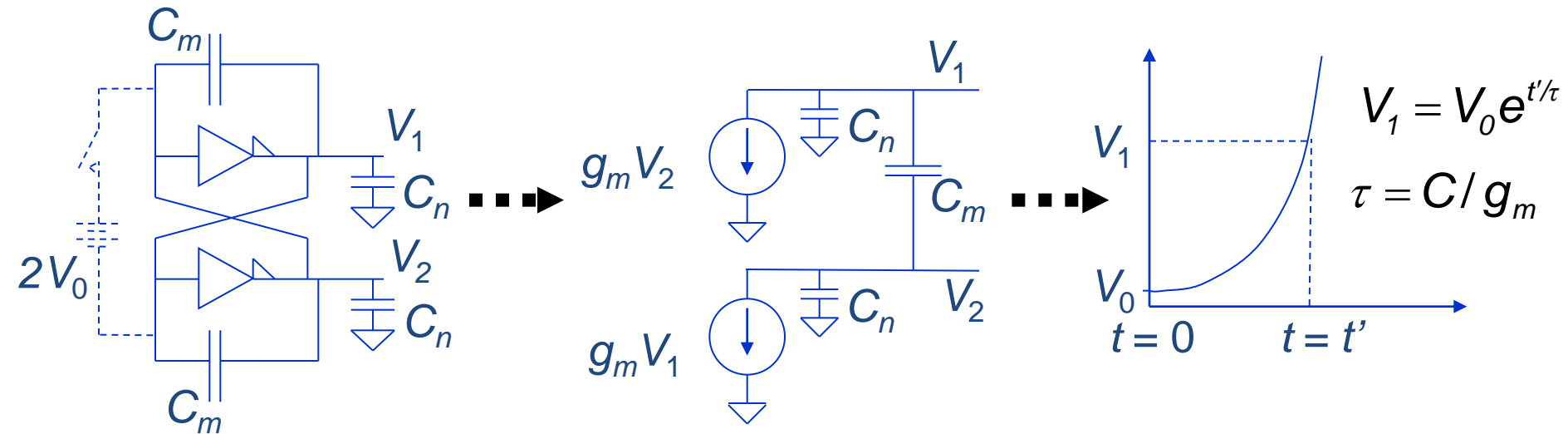
$$G_{tv} = \text{slope} = \frac{\Delta v}{\Delta t}$$



Circuit Model Analysis



Use small signal analysis

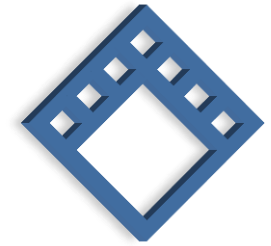


For V_0 small

Result



MTBF for Synchronizers



The probability of failure is the probability that the synchronizer output is unresolved at the next clock edge:

$$P_{unresolved} = \frac{\Delta t}{T_C} = f_C \Delta t$$

With a uniform distribution of data events in a clock period

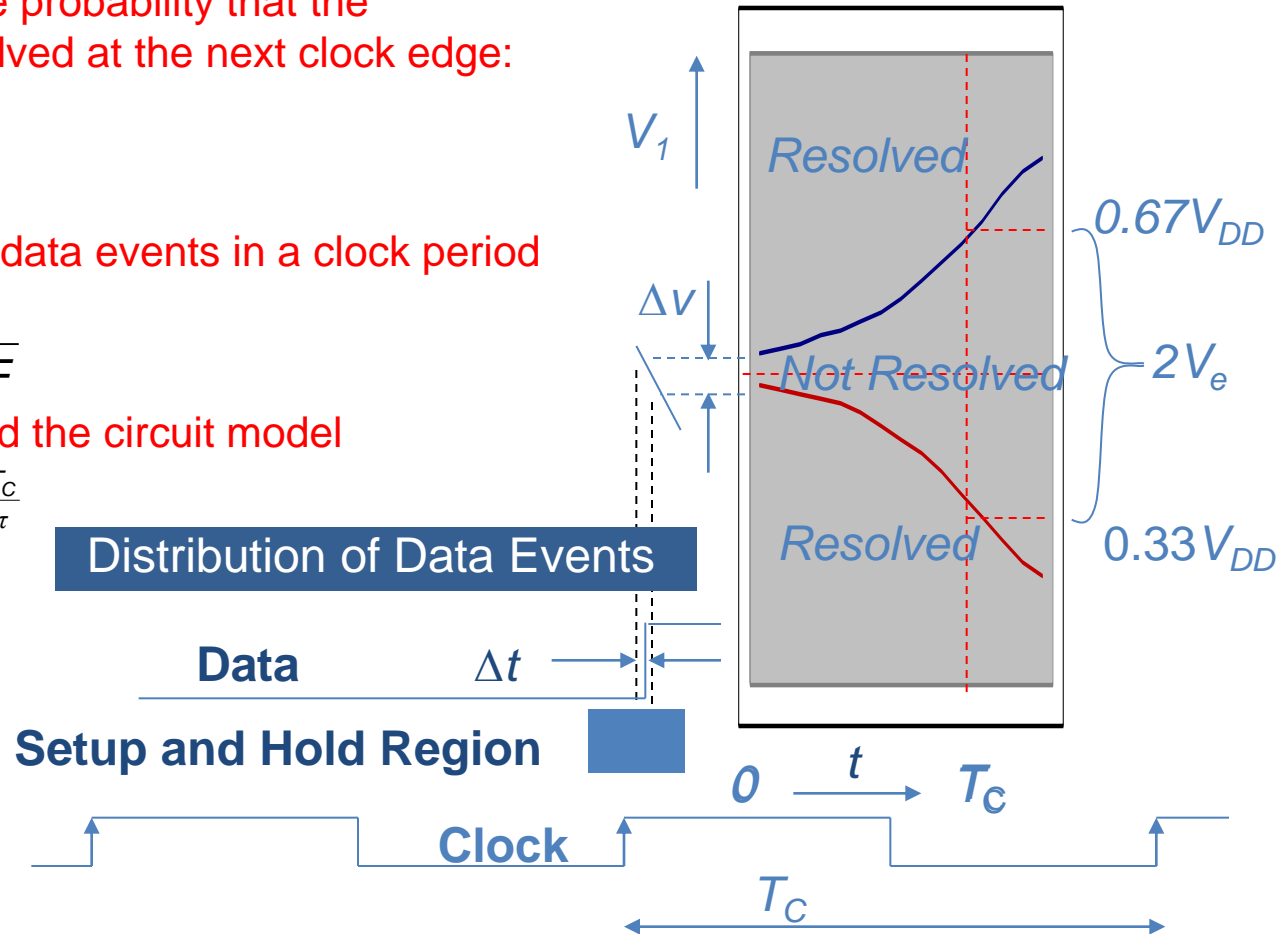
$$P_{unresolved} = f_D f_C \Delta t = \frac{1}{MTBF}$$

From the definitions of G_{tv} and the circuit model

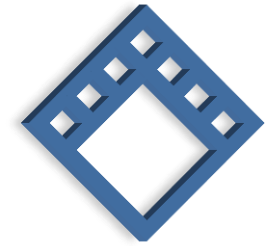
$$\Delta t = \frac{\Delta v}{G_{tv}}; \quad \Delta v = 2V_e e^{-\frac{T_C}{\tau}}$$

we see that

$$MTBF = \frac{G_{tv} e^{\frac{T_C}{\tau}}}{2V_e f_D f_C}$$



MTBF Based on Aperture Time



The probability of failure is the probability that the synchronizer output is unresolved at the next clock edge:

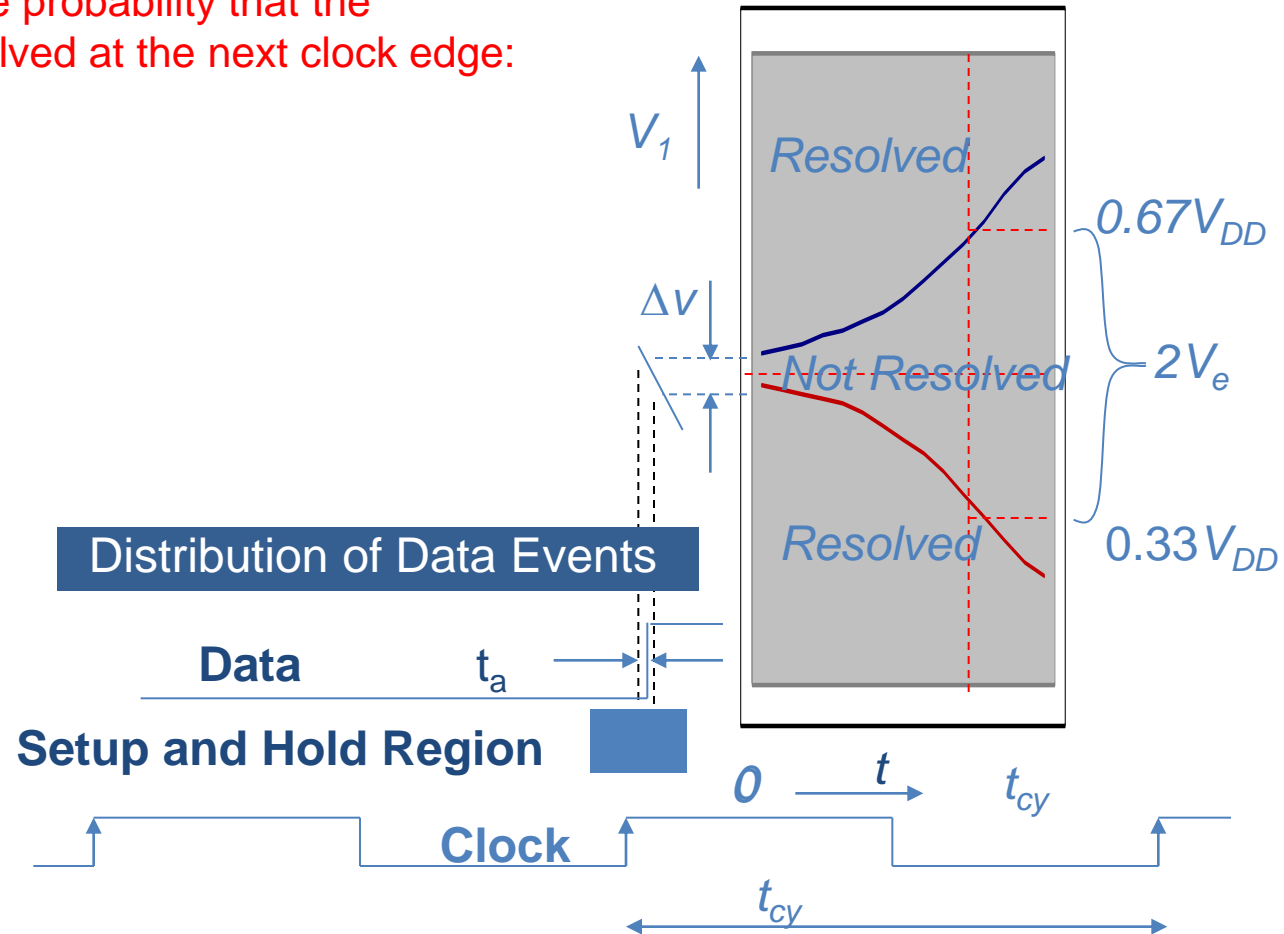
$$P_E = \frac{t_a}{t_{cy}} = t_a f_{cy}$$

$$P_S = e^{-\frac{t_w}{\tau}}$$

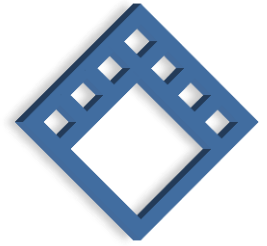
$$P_F = P_E P_S = t_a f_{cy} e^{-\frac{t_w}{\tau}}$$

$$f_F = f_e P_F = t_a f_e f_{cy} e^{-\frac{t_w}{\tau}}$$

$$MTBF = 1/f_F = \frac{e^{\tau}}{t_a f_e f_{cy}}$$



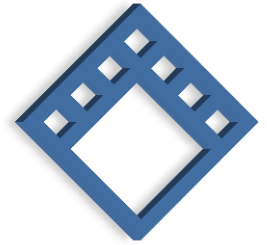
Synchronizer Failure Trend



- System failures due to synchronizer failures have been rare, but will be more likely in future
 - Many more synchronizers in use (Moore's Law)
 - Systems with 100s of synchronizers, perhaps 1000s soon
 - Systems with synchronizers in million-fold production
 - Small changes in V_t cause large changes in τ
 - Growing parameter variability in nano-scale circuits
 - In an IBM 90 nm process V_t varies for 0.4 to 0.58 volts
 - Transistor aging increases vulnerability
 - An ASU model shows V_t increasing by 5% over 5 years
 - Clock domains may not have uncorrelated clocks



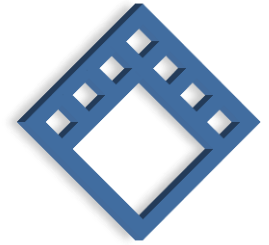
Is There A Perfect Solution?



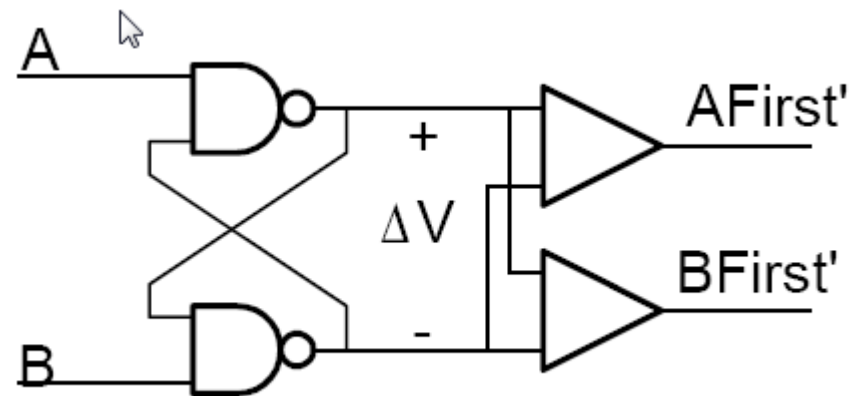
- Theoretical results show metastability is a fundamental problem of all bi-stable circuits
- Failures caused by metastability are always a possibility
 - between two independently clocked domains
 - between a clock domain and outside world
- One solution uses asynchronous circuits, but real-time applications may still be problematic
- Another solution uses synchronizer circuits and designers must hope failures are rare



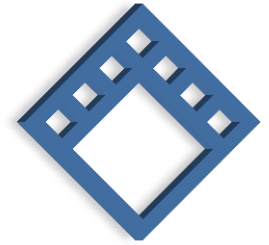
Completion Detection



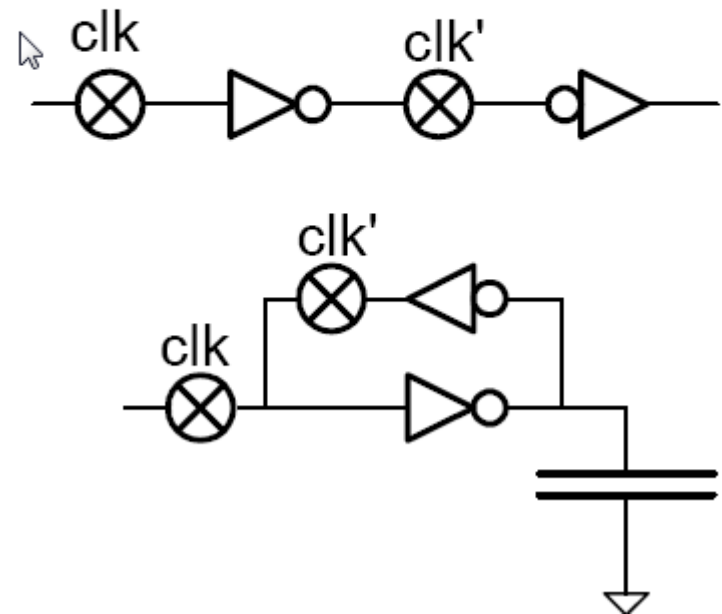
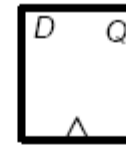
- It is not possible to bound the amount of time needed for a synchronizer to settle.
- It is, however, possible to detect when the synchronizer has settled!
- This is only useful if the downstream logic can use this asynchronous completion signal



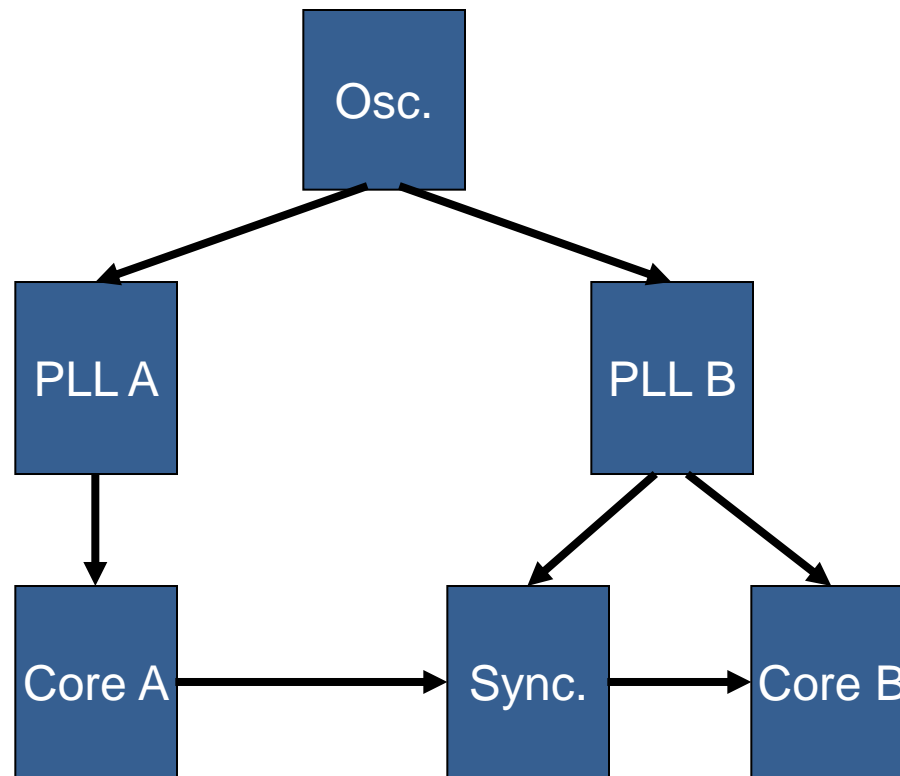
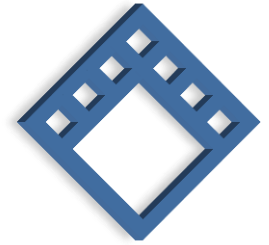
What Could Go Wrong?



- It's easy to get a synchronizer design wrong
- The three most common pitfalls are:
 - using a non-restoring (or slowly restoring) flip-flop
 - τ needs to be small
 - not isolating the flip-flop feedback loop
 - Using two flip-flops in parallel
- The last pitfall is doing everything “right” but not understanding that influences MTBF!



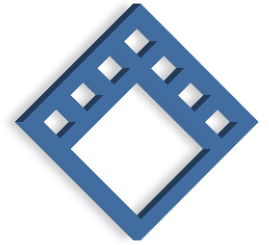
Correlated Clocks



Although Cores A and B may be clocked at different rates, these rates are based on the same oscillator and are thus correlated. This relationship between the synchronizer's clock and data inputs can be very malicious.

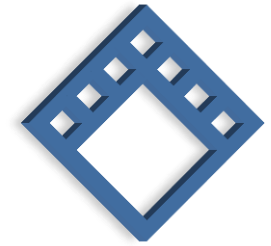


Correlated Clocks & Noise



- The effects of correlated clocks and the effects of noise can be approached similarly.
- As we will see, circuit noise may be treated as one case of correlated clocks.

Region of Vulnerability: Δt

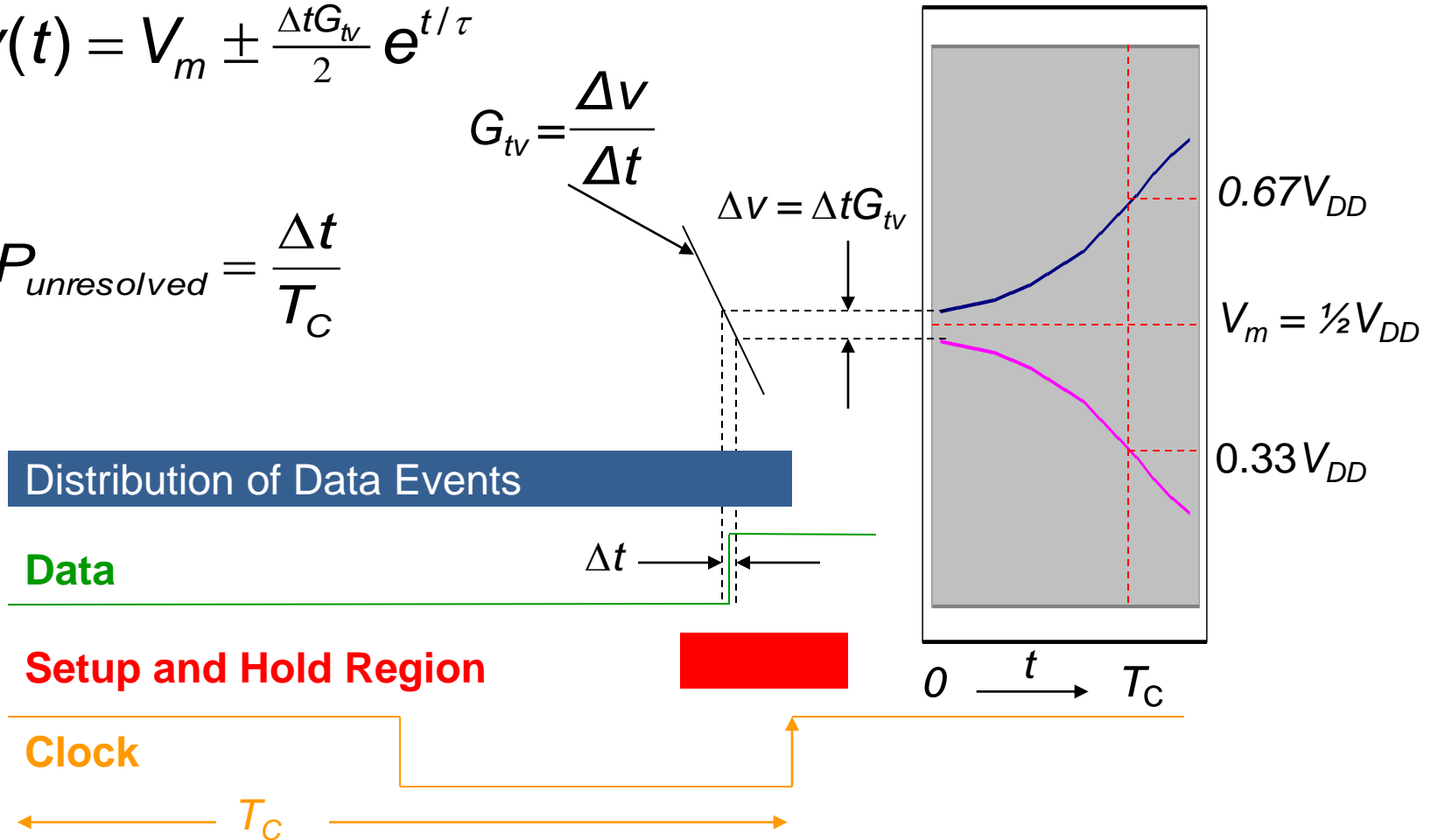


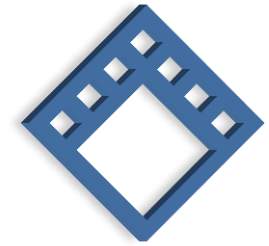
$$v(t) = V_m \pm \frac{\Delta t G_{tv}}{2} e^{t/\tau}$$

$$P_{unresolved} = \frac{\Delta t}{T_C}$$

$$G_{tv} = \frac{\Delta v}{\Delta t}$$

$$\Delta v = \Delta t G_{tv}$$





Malicious Data Events

$$v(t) = V_m \pm \frac{\Delta t G_{tv}}{2} e^{t/\tau}$$

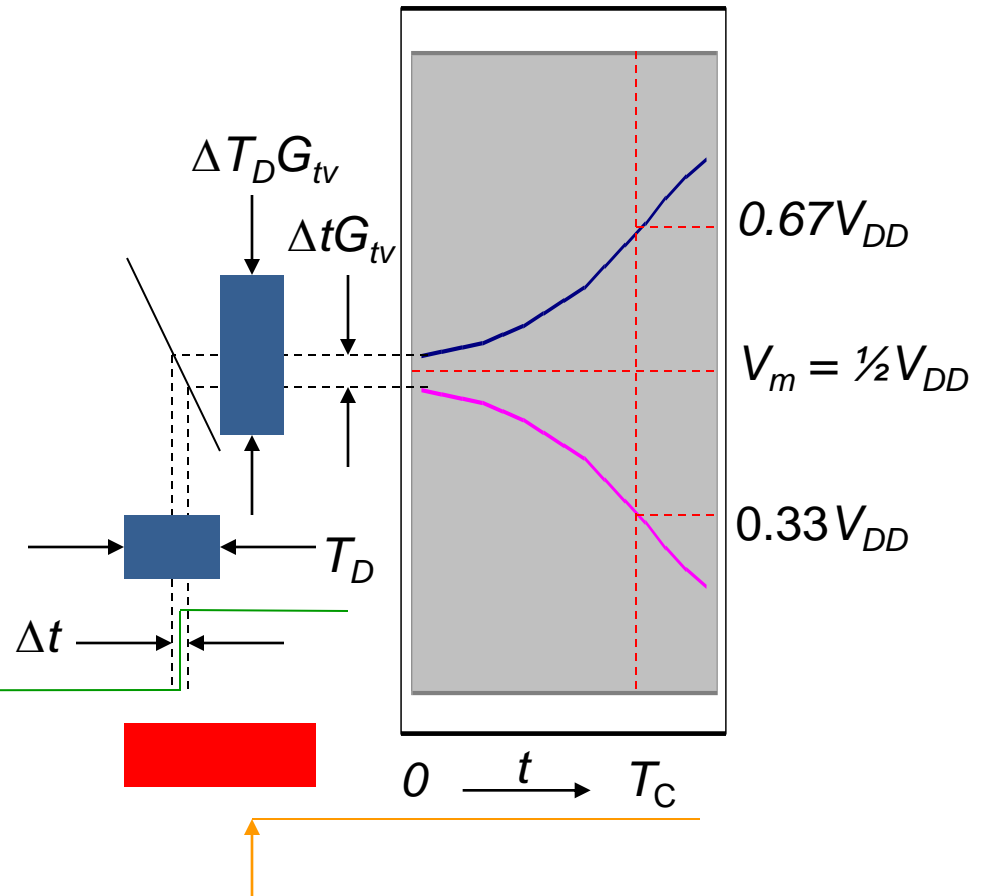
$$P_{unresolved} = \frac{\Delta t}{T_D} > \frac{\Delta t}{T_C}$$

Distribution of Data Events

Data

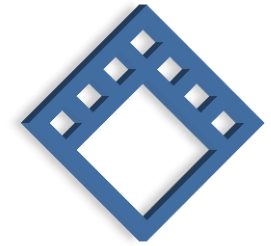
Setup and Hold Region

Clock



Malicious Data Events

Even More Malicious



$$v(t) = V_m \pm \frac{\Delta t G_{tv}}{2} e^{t/\tau}$$

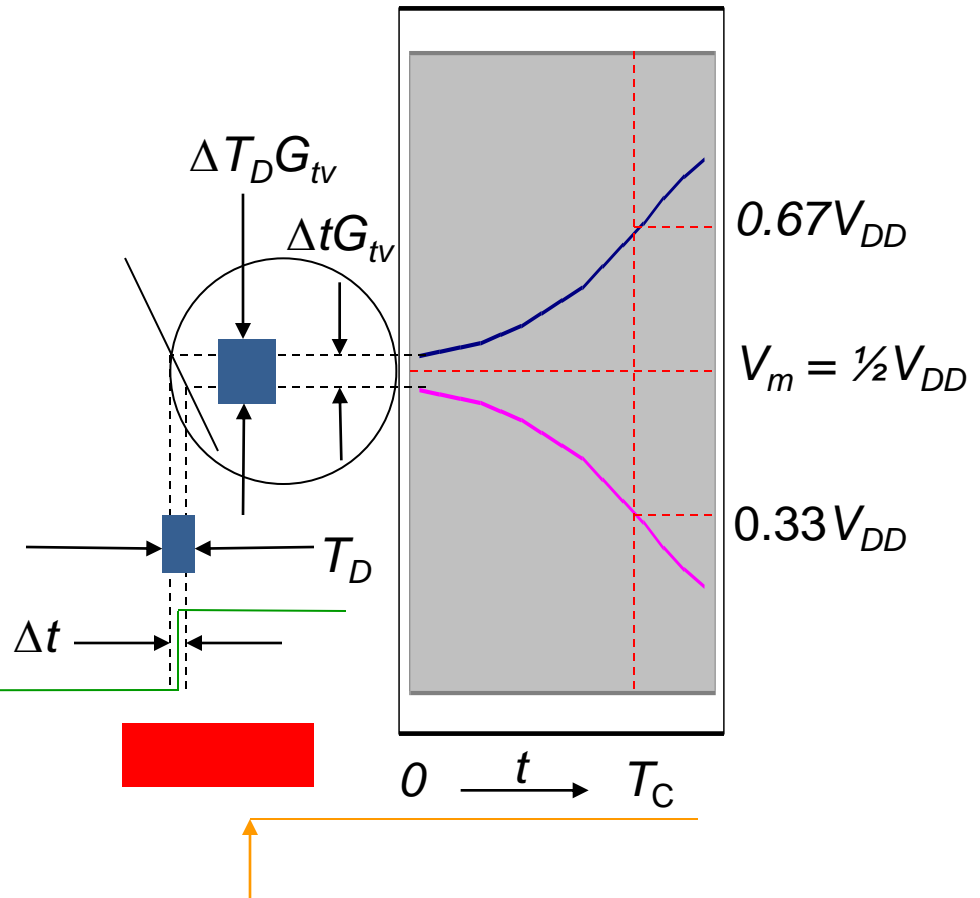
$$P_{unresolved} = \frac{\Delta t}{T_D} \gg \frac{\Delta t}{T_C}$$

Distribution of Data Events

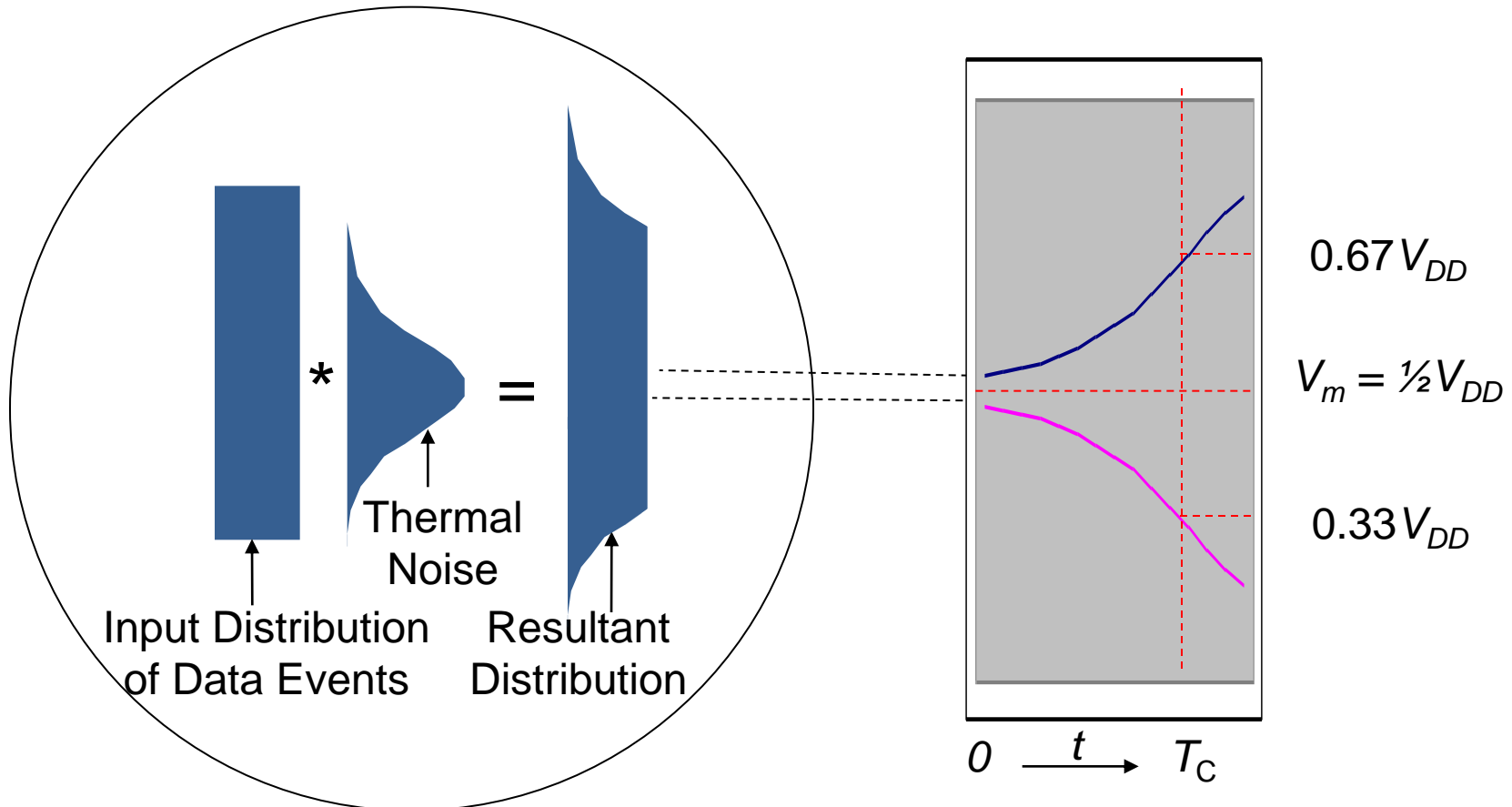
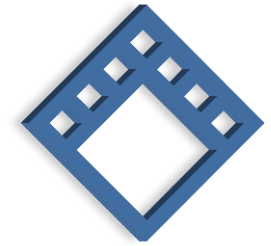
Data

Setup and Hold Region

Clock

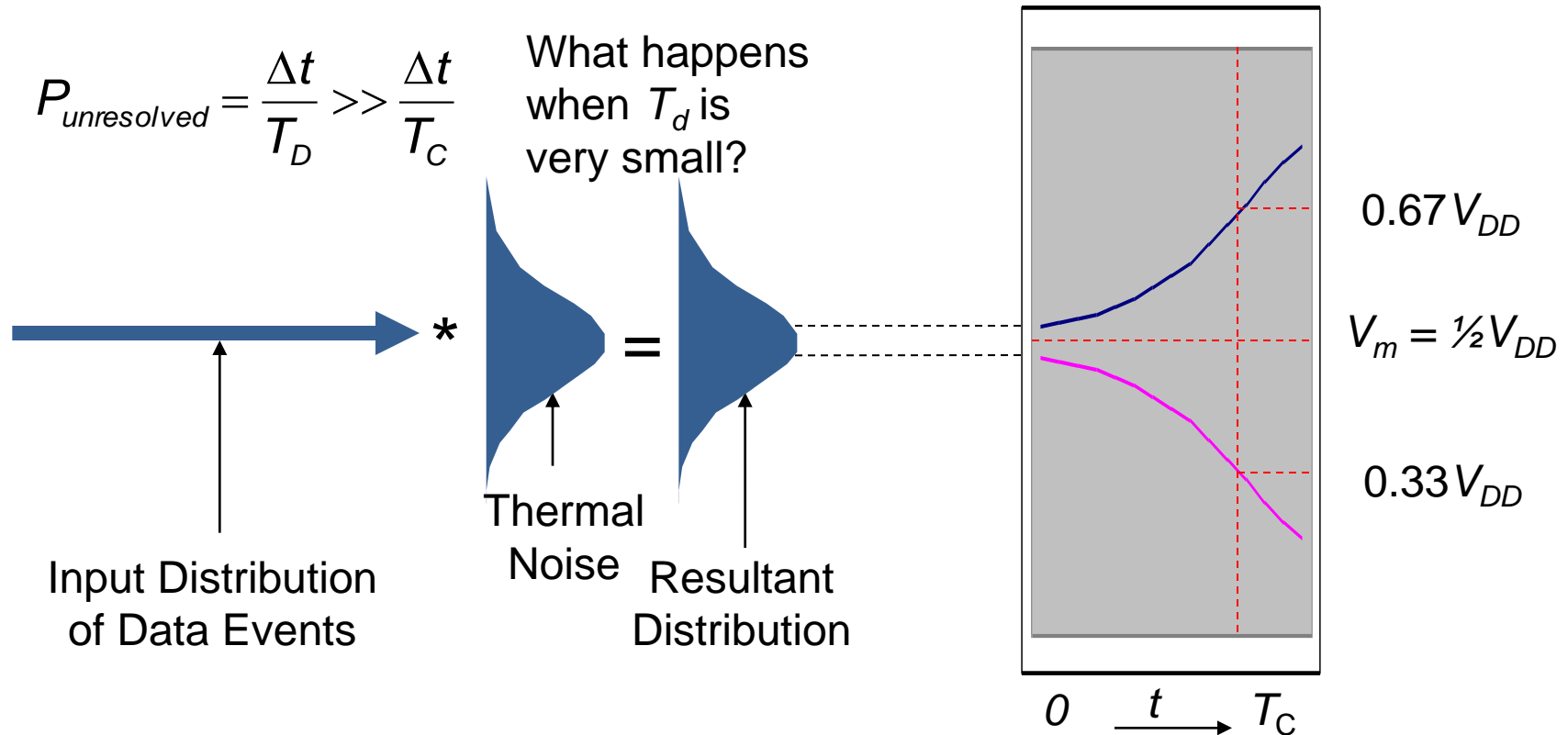
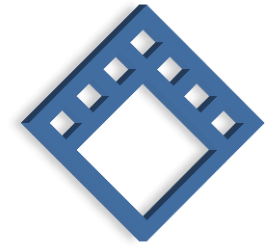


Effects of Thermal Noise



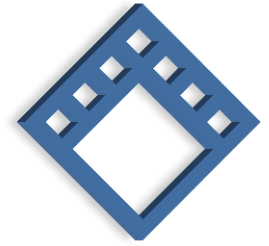
Bottom Line: Thermal noise pushes as many events into the window of vulnerability as is pushed out.

Upper Bound on $P_{unresolved}$



Bottom Line: Thermal noise establishes an upper bound on $P_{unresolved}$ and a lower bound on MTBF

Calculating MTBF



- Always a stochastic calculation

- Assume clock and data unrelated

$$MTBF(FF \text{ unresolved at } t) = \frac{G_{tv} e^{t/\tau}}{2V_e f_D f_C}$$

- If related, thermal noise gives lower bound

- E.g. clock and data from same source or clockless

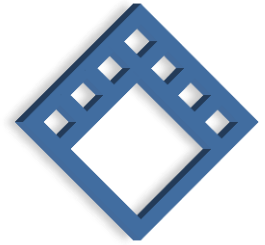
$$MTBF(FF \text{ unresolved at } t) \approx \frac{\sigma}{V_e f_D} e^{t/\tau}$$

- Thermal noise voltage standard deviation: $\sigma = \sqrt{2kT/C}$

- This lower bound is 2 to 3 orders of magnitude smaller than when clock and data are unrelated

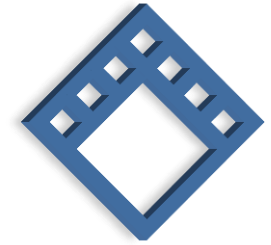


MTBF Affects System Behavior



- Assume:
 - Desired probability of system failure = 1 : 2,000,000
 - System lifetime is 30 years ($\sim 10^9$ sec)
 - System has 50 processors with 10 synchronizers each
- Then:
 - Need MTBF of 30 billion years ($3 \cdot 10^{10}$) per synchronizer
- But:
 - Corner cases can further reduce needed MTBF
 - If clock and data are related, must use lower bound set by thermal noise: $MTBF_n$
- **Unwise to use conventional MTBF formula without understanding its limitations**

Master-Slave DFF MTBF Examples



Clock Frequency (MHz)	MTBF (yrs)	MTBF _n (yrs)
200	9.7E+37	2.1E+35
300	4.3E+19	1.4E+17
500	7.5E+04	4.1E+02
750	2.7E-03	2.2E-05

90 nm process

$\tau=39.83$ ps, $G_{tv}=0.375$ V/ns, $f_d = 133$ MHz

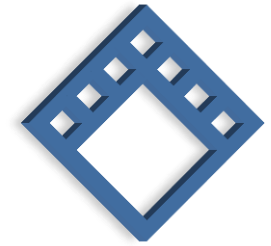
125 ps setup time assumed

MTBF ranges from 1 day to $9.7 \cdot 10^{37}$ years

MTBF_n ranges from 11.5 minutes to $2.1 \cdot 10^{35}$ years



Parameter Variations in Master-Slave Process-Voltage-Temperature 200 MHz



	τ (ps)	Gtv (V/ns)	MTBF (yrs)	MTBF _n (yrs)
-3 sigma	106.49	0.369	5.07E+04	1.12E+02
-1 sigma	55.50	0.543	1.37E+23	2.06E+20
Nominal 0 degrees	39.30	0.751	1.00E+39	1.04E+36
Nominal 27 degrees	39.83	0.375	9.79E+37	2.13E+35
Nominal 70 degrees	41.01	0.301	2.29E+36	6.65E+33
1 sigma	28.98	0.866	1.80E+58	1.70E+55
3 sigma	16.69	0.031	4.16E+110	1.09E+109

200 MHz Clock; 90 nm process, 125 ps setup time

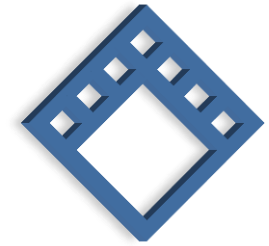
MTBF ranges from $5.07 \cdot 10^4$ years to $4.16 \cdot 10^{110}$ years

MTBF_n ranges from 112 years to $1.09 \cdot 10^{109}$ years



Latch Versus Master-Slave FF

MTBF @200 MHz



	τ (ps)	Gtv (V/ns)	MTBF (yrs)
Master-Slave FF	39.83	0.375	9.8E+37
Latch	40.54	4.729	1.4E+38

200 MHz Clock; 90 nm process, 125 ps setup time

