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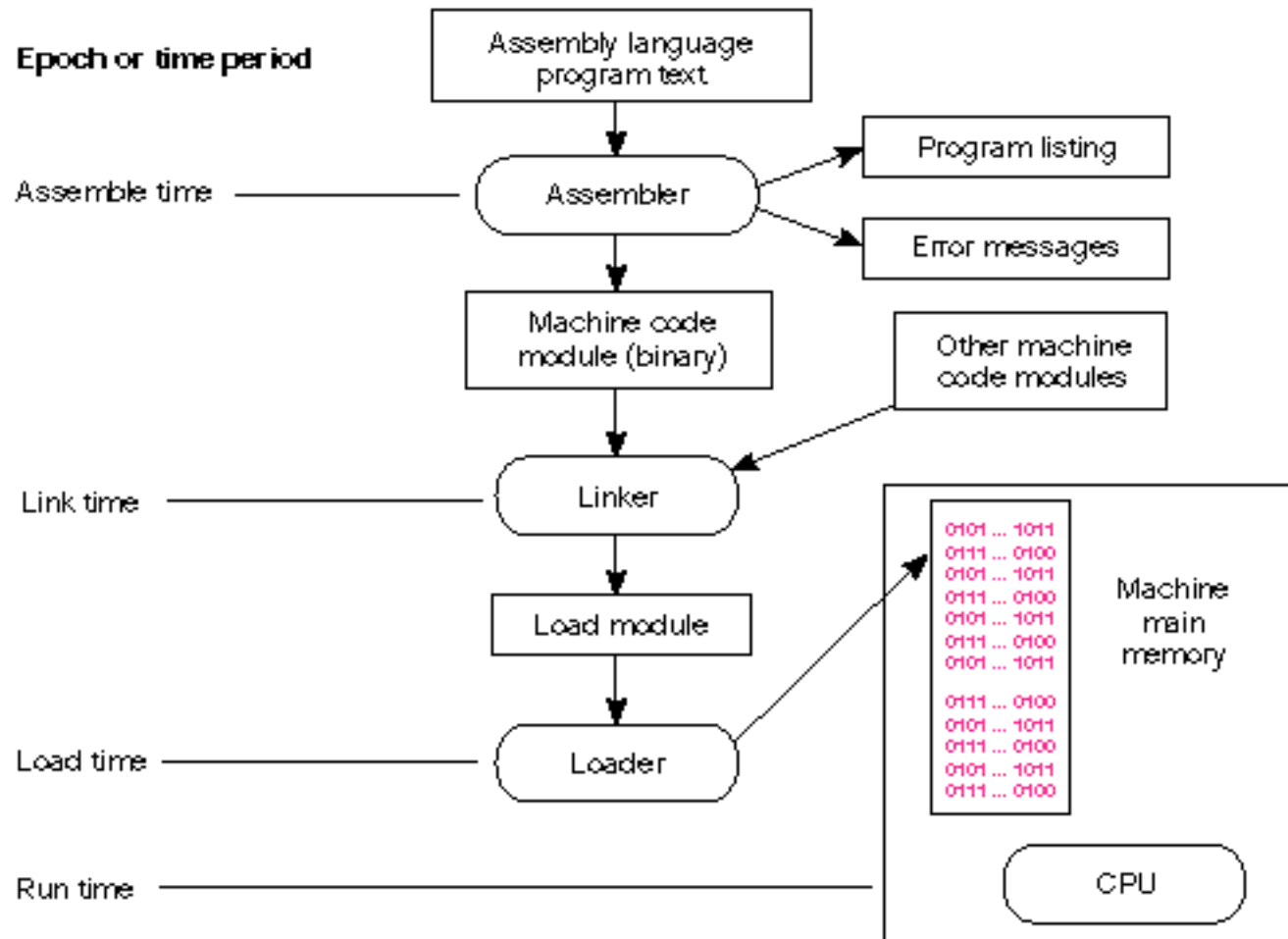
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Assembly and Assemblers

- What is an assembler?
 - Translates from assembly language to machine language
- Assembly language structure
- Tasks of the assembler

Fig C.1 Time Periods of Various Processes in Program Development



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The Assembler Provides—

- Access to all the machine's resources by the assembled program. This includes access to the entire instruction set of the machine.
- A means for specifying run-time locations of program and data in memory.
- Provide symbolic labels for the representation of constants and addresses.
- Perform assemble-time arithmetic.
- Provide for the use of any synthetic instructions.
- Emit machine code in a form that can be loaded and executed.
- Report syntax errors and provide program listings
- Provide an interface to the module linkers and program loader.
- Expand programmer defined macro routines.

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Assembler Syntax and Directives

- Syntax: Label OPCODE Op1, Op2, ... ;Comments
- Pseudo Operations (sometimes called “pseudos,” or directives are “Opcodes” that are actually instructions to the assembler, and that do not result in code being generated.
- Assembler maintains several data structures
 - Table that maps text of opcodes to op number and instruction format(s)
 - “Symbol Table” that maps defined symbols to their value

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Table C.1 Assembler Directives

Directive Type	Example	Action
Define a symbolic constant	<pre>NOV .equ 11 MAY .equ "May" Size .equ 10</pre>	Set the symbol <code>NOV</code> equal to the constant 11. Set the symbol <code>MAY</code> equal to the string "May". Set the symbol <code>Size</code> equal to the constant 10.
Fix a memory location	<pre>Start .org 0x2000 main: ld r0, Size</pre>	Fix the location at which the following program or data word will load to 2000_{16} . Begin program execution at this location.
Reserve a block of storage	<pre>Array: .dcb 20 Harray: .dch Size Warray .dc 20</pre>	Reserve space for 20 bytes. Base address is <code>Array</code> . Reserve space for 10 halfwords. Base address is <code>Harray</code> . Reserve space for 20 words. Base address is <code>Warray</code> .
Initialize memory location(s)	<pre>Minus1: .dch 0xffff Colors: .dc 0, 1, 2, 3 Hi .dcb "Hello"</pre>	Reserve a halfword at location <code>Minus1</code> ; initialize it to the hexadecimal value <code>ffff</code> . Reserve space at location <code>Colors</code> for 4 words, and initialize them to 0, 1, 2, and 3. Reserve space at location <code>Hi</code> for 5 bytes; init. to 'Hello'
Describe [†] module linkage	<pre>Out: ... PUBLIC ... Val EXTERN</pre>	Make the value of <code>Out</code> available to the linker for linkage to other assembled modules. The value of <code>Val</code> is defined externally in another module.

[†]Not available in SRC, but available in production assemblers.

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Figure C.2 Example Program fib.asm

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; fib.asm. Compute Fibonacci numbers.
; The Fibonacci sequence is defined as follows:
; fib(1) = 1, fib(2) = 1,
; fib(n) = fib(n-1) + fib(n-2) n > 2.
cnt:    .equ 8          ; No. to compute after first two
        .org 0          ; Store sequence at addr. 0
seq:    .dc 1, 1        ; Init. the first two Fib. Nos.
        .dw cnt         ; Storage for the next 8 Fib. Nos.
        .org 0x1000    ; Begin ass'y. at hex. addr. 1000
main:   lar r31, loop   ; Pgm start. Init. branch address
        la r0, cnt      ; Init. count
        la r1, 0        ; Init r1 to index of seq[0]
        la r2, 4        ; Init r2 to index of seq[1]
loop:   ld r3, seq(r1)  ; Get fib(n-2)
        addi r1, r1, 4  ; Increment index
        ld r4, seq(r1)  ; Get fib(n-1)
        add r3, r3, r4  ; compute fib(n)
        addi r2, r2, 4  ; fib(n) = fib(n-1) + fib(n-2)
        st r3, seq(r2)  ; Store fib(n)
        addi r0, r0, -1 ; Decrement count
        brnz r31, r0    ; loop untill done
        stop

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Figure C.3 Assembler listing file, fib.lst

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*****SRC Assembler***** (SRCTools Version 2.1.0)

HexLoc      DecLoc      MachWord    Label      Source Code      Comments
00000000    0000000000    00000000    ; fib.asm. Compute Fibonacci numbers.
00000000    0000000000    00000000    ; The Fibonacci sequence is defined as follows:
00000000    0000000000    00000000    ; fib(1) = 1, fib(2) = 1,
00000000    0000000000    00000000    ; fib(n) = fib(n-1) + fib(n-2) n > 2.
00000000    0000000000    00000000    cnt: .equ 8      ; No. to compute after first two
00000000    0000000000    00000000    .org 0        ; Store sequence at addr. 0
00000000    0000000000    00000001    seq: .dc 1    ; Init the first two numbers
00000004    0000000004    00000001    .dc 1        ;
00000008    0000000008    00000000    .dw cnt      ; Storage for the next 8 Fib. Nos.
00000028    0000000040    00000000    .org 0x1000  ; Begin ass'y. at hex. addr. 1000
00001000    0000004096    37c0000c    main: lar r31, loop ; Pgm start: init branch address
00001004    0000004100    28000008    .la r0, cnt   ; Init. count
00001008    0000004104    28400000    .la r1, 0    ; Init r1 to index of seq[0]
0000100c    0000004108    28800004    .la r2, 4    ; Init r2 to index of seq[1]
00001010    0000004112    08c20000    loop: ld r3, seq(r1) ; Get fib(n-2)
00001014    0000004116    68420004    .addi r1, r1, 4 ; Increment index
00001018    0000004120    09020000    .ld r4, seq(r1) ; Get fib(n-
1)
0000101c    0000004124    60c64000    .add r3, r3, r4 ; compute fib(n)
00001020    0000004128    68840004    .addi r2, r2, 4 ; fib(n) = fib(n-1) + fib(n-2)
00001024    0000004132    18c40000    .st r3, seq(r2) ; Store fib(n)
00001028    0000004136    6801ffff    .addi r0, r0, -1 ; Decrement count
0000102c    0000004140    403e0003    .brnz r31, r0  ; loop untill done
00001030    0000004144    f8000000    .stop

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The 2-Pass Assembly Process

Pass 1:

- Init. Location Counter (Assemble-Time “PC”) to 0
- Pass over program text: enter all symbols into symbol table
 - May not be able to map all symbols on first pass
 - (Definition before use is usually allowed)
- Determine size of each instruction, map to a location
 - Uses pattern matching to relate opcode to pattern
 - Increment location counter by size
 - Change Location Counter in response to ORG pseudos.

Pass 2

- Insert binary code for each opcode and value
- “Fix up” forward references and variable-sizes instructions
 - Examples include variable-sized branch offsets and constant fields.

Table C.2 Snapshot of the Symbol Table Generated by the Assembler During Pass 1 at Address 0x1008

Key	Symbol	Type	Value	Defined
0	"cnt"	constant	8	defined
1	"seq"	label	00000000	defined
2	"main"	label	00001000	defined
3	"loop"	unknown		undefined