Reference Materials

1. RTN Description of SRC

Unified RTN for SRC

Below is the entire RTN description for SRC with reset and exception handling integrated into it.

Memory

Processor state

PC 310 :	program counter (address of the next instruction)
IR 310 :	instruction register
Run:	one bit run/halt indicator
Strt:	start and hard reset signal
Rst:	soft reset signal
R[031] 310 :	general purpose registers

Processor interrupt mechanism

ireq:	interrupt request signal
iack:	interrupt acknowledge signal
IE:	one bit interrupt enable flag
IPC 310:	storage for PC saved upon interrupt
II 310 :	interrupt info.: about source of last interrupt
Isrc_info 150 :	information from interrupt source
Isrc_vect 70 :	type code from interrupt source
Ivect 310 := 20@0#	#Isrc_vect 70 #4@0:

Main memory state

Mem[02 ³² - 1] 70 :	2^{32} addressable bytes of memory
M[x] 310 := Mem[x]#Mem	n[x+1]#Mem[x+2]#Mem[x+3]:

Formats

Instruction formats

op 40 := IR 3127 :	operation code field
ra 40 := IR 2622 :	target register field
rb 40 := IR 2117 :	operand, address index, or branch target
rc 40 := IR 1612 :	2nd operand, conditional test, or shift count
c1 210 := IR 210 :	long displacement field
c2 160 := IR 160 :	short displacement or immediate field
c3 110 := IR 110 :	count or modifier field

Branch condition format

cond := (c3 20)	0 = 0 0:	never
c3 20 =1	1:	always
c3 20 =2	R[rc]=0:	if register is zero
c3 20 =3	R[rc] 0:	if register is nonzero
c3 20 =4	R[rc] 31 =0:	if register is positive or zero
c3 20 =5	R[rc] 31 =1):	if register is negative

Shift count format

$n := ((c3 \ 40 = 0))$	R[rc] 40 :	shift count is register or
(c3 40 0)	c3 40):	constant field of instruction

Effective address calculations

disp 310 := ((rb	e=0) c2 160	{sign extend}:	disp.
(rb 0)	$R[rb] + c2 \ 160$	{sign extend, 2's complement})	: addr.
rel 310 := PC 3	10 + c1 210	{sign extend, 2's comp.}:	el. adr.

Instruction interpretation

instruction_interpretation :=

$(\neg \text{Run Strt} (\text{Run 1: PC, R[031]} 0;$	Hard reset
instruction_interpretation):	
Run Rst (Rst 0: IE 0: PC 0;	Soft reset
instruction_interpretation):	
Run ¬ Rst (ireq IE) (IPC PC 310:	Interrupt
II 150 Isrc_info 150 :	
IE 0: PC Ivect 310 :	
iack 1; iack 0;	
instruction_interpretation):	
Run \neg Rst \neg (ireq IE) (IR M[PC]:	Normal fetch
PC $PC + 4$; instruction_execution):	

Instruction execution instruction_execution := (

Load and store instruct	tions					
ld (:= op= 1)	R[ra]	M[disp]:	lo	oad register	:	
ldr (:= op= 2)	R[ra]	M[rel]:	lo	oad register	relative	
st (:= op= 3)	M[disp]	R[ra]:	S	tore registe	r	
str (:= op= 4)	M[rel]	R[ra]:	S	tore registe	r relative	
la (:= op= 5)	R[ra]	disp:	lo	oad displac	ement addre	ess
lar (:= op= 6)	R[ra]	rel:	le	oad relative	address	
Branch instructions						
br (:= op= 8)	(cond	PC R[rb]):		cond. bra	anch
brl (:= op= 9)	(R[ra]	PC: cond	(PC	R[rb])):	branch &	z link
Arithmetic instructions	s (assume	d to be 2's co	mpleme	nt arithmet	ic)	
add (:= op= 12)) R[ra]	R[rb] + R	[rc]:			
addi (:= op= 13	B) R[ra]	R[rb] + c	2 160	{2's comp.	. sign ext.}:	
sub (:= op= 14)	R[ra]	R[rb] - R[[rc]:			
neg (:= op = 15)) R[ra]	-R[rc]:				
and $(:= op = 20)$) R[ra]	R[rb] R	[rc]:			
andi (:= op= 21) R[ra] R[rb] c	2 160	{sign exte	nd}:	
or (:= op= 22)	R[ra]	R[rb] R[ı	rc]:			
ori (:= op= 23)	R[ra]	R[rb] c2	160 {	sign extend	d}:	
not (:= op= 24)	R[ra]	$\neg R[rc]$:				
Shift instructions						
shr (:= op= 26)	R[ra]	310 (n (@ 0) # R	[rb] 31n :	: r	ight
shra (:= op= 27) R[ra]] 310 (n	@ R[rb]	31) # R[1	cb] 31n : a	arith.
shl (:= op= 28)	R[ra]	310 R[rb)] 31-n() #(n @ 0):	: le	eft
shc (:= $op=29$)	R[ra]	310 R[rl	o] 31-n	0 #R[rb] 3	132-n : c	irc.
Interment instructions						
$\frac{\text{Interrupt instructions}}{\text{con}(1-\text{con}-10)}$	(IE 1).		08.00	ntion anabl	
een(op - 10)).).		CALE	ption disabl	
eut(op - 11)	$(\mathbf{IE} \ \mathbf{U})$). $C = I = 1$		exce	ption disabi	C
m(op - 30)	$(\Gamma C) = \Gamma$	$\begin{array}{cccc} \mathbf{L} & \mathbf{I} \\ 0 & \mathbf{I} & \mathbf{I} \\ \mathbf{I} & \mathbf{I} \\ \mathbf{I} & \mathbf{I} \\ \mathbf{I} \\$	$0 \cdot \mathbf{P}[\mathbf{r}\mathbf{b}]$		1 0	ntorrupt state
svi(op - 10)	(K[14] 15 II 15 0	$\mathbf{D}[ro] 15 0$	$0 \cdot \mathbf{DC}$	$ \begin{bmatrix} \mathbf{I} & \mathbf{C} \\ \mathbf{I} \\ \mathbf{I} \end{bmatrix} $	[rb] (rb)	interrupt state
II(0p-17) (11 150	K [1a] 130	· IFC :	510 K	[10]).1051010	interrupt state
Miscellaneous instruct	ions					
nop (:= op= 0)	:		N	lo operation	n	
stop (:= op= 31) Run	0	S	top instruc	tion	
);			E	and of instruction	uction_exec	ution
instruction_inte	erpretation	1.				

A.19 RegisterTransfer Notation - RTN

Register transfer: register on LHS stores value from RHS. [] Word index: selects word or range from a named memory. Bit index: selects bit or bit range from named register. Index range: from left index n to right index m; can be decreasing. n..m If-then: true condition on left yields value and/or action on right. Definition: text substitution with dummy variables. := # Concatenation: bits on right appended to bits on left. Parallel separator: actions or evaluations carried out simultaneously. : Sequential separator: RHS evaluated and/or performed after LHS. ; **(***a*) Replication: LHS repetitions of RHS are concatenated. { } Operation modifier: describes type of preceding operation. ()Nested grouping of operations or values: operators or separators. Comparisons: produce 0 or 1 (true or false) logical value. > = < Arithmetic operators: also , , and mod. $+ - \times \div$ Logical operators: and, or, not, exclusive or, equivalence. _

Notes:

Expressions can be values and/or actions. Actions can be considered side effects if a value is present.

A list of conditional expressions need not have disjoint conditions. Right hand sides of conditionals are evaluated for all conditions which are true. No sequencing is implied unless there are sequential separators between conditional expressions. There is no *else* equivalent.

Pseudo-operations.

.org	Value	Load the program starting at address Value.
.equ	Value	Define the Label symbol to be the constant Value.
.dc	Value [,Value]	Allocate memory words and set to the 32 bit Values.
.dcb	Value [,Value]	Allocate bytes and load them with the 8 bit Values.
.dch	Value [,Value]	Allocate halfwords and load with the 16 bit Values.
.db	Count	Allocate storage for Count bytes.
.dh	Count	Allocate storage fot Count 16 bit halfwords.
.dw	Count	Allocate storage for Count 32 bit words.

Table A.1 SRC assembly language pseudo operations

Values are assumed to be decimal unless terminated by B (binary) or H (hexadecimal).

SRC Instruction Mnemonics	I. MC68000 Reference Materials
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ld	ra, c2	Load from absolute address. rb is register 0.
ld	ra, c2(rb)	Load from displacement address.
ldr	ra, cl	Load from relative address.
st	ra, c2	Store into absolute address. rb is register 0.
st	ra, c2(rb)	Store into displacement address.
str	ra, cl	Store into relative address.
la	ra, c2	Load value of absolute address into ra. rb is reg. 0.
la	ra, c2(rb)	Load value of displacement address into ra.
lar	ra, cl	Load value of relative address into ra.
add	ra, rb, rc	Add rb to rc and put result in ra.
addi	ra, rb, c2	Add rb to immediate constant and put result in ra.
sub	ra, rb, rc	Subtract rc from rb and put result in ra.
neg	ra, rc	Place two's complement negative of rc into ra.
or	ra, rb, rc	OR rb and rc and put result in ra.
ori	ra, rb, c2	OR rb and immediate constant and put result in ra.
and	ra, rb, rc	AND rb and rc and put result in ra.
andi	ra, rb, c2	AND rb and immediate constant and put result in ra.
not	ra, rc	Place logical NOT of rc into ra.
shr	ra, rb, c3	Shift rb right into ra by constant shift count c3 31.
shr	ra, rb, rc	Shift rb right into ra by count in rc. c3 is 0.
shra	ra, rb, c3	Shift rb right with sign extend into ra by constant c3.
shra	ra, rb, rc	Shift rb right with sign extend into ra by count in rc.
shl	ra, rb, c3	Shift rb left into ra by constant c3.
shl	ra, rb, rc	Shift rb left into ra by count in rc. c3 is 0.
shc	ra, rb, c3	Shift rb left circularly into ra by constant c3.
shc	ra, rb, rc	Shift rb left circularly into ra by count in rc. c3 is 0.
br	rb, rc, c3	Branch to target in rb if c3 satisfies condition c3.
brl	ra,rb,rc,c3	Branch to rb if rc satisfies c3 and save PC in ra.
br	rb	Branch unconditionally to rb.
brl	ra, rb	Branch unconditionally to rb and save PC in ra.
brlnv	ra	Do not branch, but save PC in ra.
brzr	rb, rc	Branch to rb if rc is zero.
brlzr	ra, rb, rc	Branch to rb if rc is zero and save PC in ra.
brnz	rb, rc	Branch to rb if rc is non-zero.
brlnz	ra, rb, rc	Branch to rb if rc is non-zero and save PC in ra.
brpl	rb, rc	Branch to rb if rc is positive or zero (sign is plus).
brlpl	ra, rb, rc	Branch to rb if rc is positive and save PC in ra.

 Table A.2
 SRC instructions—assembly language form

brmi	rb, rc	Branch to rb if rc is negative (sign is minus).				
brlmi	ra, rb, rc	Branch to rb if rc is negative and save PC in ra.				
nop		No operation. Used to insert pipeline bubble.				
stop		Set Run to zero, halting the machine.				
een		Exception enable. Set overall exception enable bit.				
edi		Exception disable. Clear overall exception enable.				
rfi		Return from interrupt. PC IPC; enable exceptions.				
svi	ra, rb	Save II and IPC in ra and rb, respectively.				
ri	ra, rb	Restore II and IPC from ra and rb, respectively.				

 Table A.2
 SRC instructions—assembly language form

5 4 3 2 1 0 ‡ Mode Reg							
Addressing mode Name	Mode #	Reg #	Notation [*]	Extra Word	Operand location		
Data register direct	0	0-7	Dn	0	Dn		
Address register direct	1	0-7	An	0	An		
Address register indirect	2	0-7	(An)	0	Mem[An]		
Autoincrement	3	0-7	(An)+	0	Mem[An]; An An+WS		
Autodecrement	4	0-7	-(An)	0	An An-WS; Mem[An]		
Based	5	0-7	disp16 (An)	1	Mem[An + disp16]		
Based indexed short	6	0-7	disp8(An, XnLo)	1	Mem[An + XnLo + disp8]		
Based indexed long	6	0-7	disp8(An, Xn)	1	Mem[An + Xn + disp8]		
Absolute short	7	0	addr16	1	Mem[addr16]		
Absolute long	7	1	addr32	2	Mem[addr32]		
Relative	7	2	disp16(PC)	1	Mem[PC + disp16]		
Relative indexed short	7	3	disp8 (PC, XnLo)	1	Mem[PC + XnLo + disp8]		
Relative indexed long	7	3	disp8(PC, Xn)	1	Mem[PC + Xn + disp8]		
Immediate	7	4	#data	1-2	no location, data is value		

Motorola MC68000 addressing modes

‡ When the 6-bit field specifies the dst (destination) operand of a MOVE instruction, the mode and reg fields are reversed.

* An and Dn denote one of the 8 address or data registers respectively

WS = word size in bytes: 1, 2, or 4.

disp8 and disp16 are 8- and 16-bit displacements.

Xn is one of D0-D7, or A0-A7.

XnLo is the low order 16 bits of register Xn, sign extended to 32 bits.

All values smaller than 32 bits are sign extended to 32 bits before addition.

data is an 8-, 16- or 32-bit value as indicated by .B, .W, or .L in the instruction.

Mnemonic	Operands	Opcode Word [‡]	XNZVC Operation		Operand Size	
MOVE.B	EAs,EAd	0001dddddsssss	-xx00	dst	src	byte
MOVE.W	EAs,EAd	0011ddddddsssss	-xx00	dst	src	word
MOVE.L	EAs,EAd	0010ddddddsssss	-xx00	dst	src	long
MOVEA.W	EAs,An	0011rrr001ssssss		An	src	word
MOVEA.L	EAs,An	0010rrr001sssss		An	src	long
LEA.L	EAc,An	0100aaa111ssssss		An	EAc	Addr.
EXG	Dx, Dy	1100rrr1mmmmyyy		Dx	Dy	long

1 Notes:

MC68000 Data movement instructions

EAS: Source EA-any addressing mode, except cannot move byte to address register

EAd: Destination EA-any addressing mode except immediate or relative.

EAc : Control EA-all modes except register, auto increment, autodecrement, or immediate

SSSSS, dddddd = src and dst addressing mode specifiers. see Table top.

rrr, yyy = one of eight registers.

aaa = one of the eight address registers

An, Dn, one of the eight address or data registers, respectively.

mmmmm = mode field: 01000-exchange data regs.; 01001-exchange address regs.; 10001-exchange data and addr regs., where xxx specifies the data reg., and yyy specifies the address reg.

Condition codes: - = unched. from previous value, x = ched. by the operation. 0, 1 = value.

Mnemonic	Operands	Opcode word	XNZVC	Operation	Oprnd size
ADD	EA,Dn	1101rrrmmmaaaaaa	xxxxx	dst dst+src	b,w,l
SUB	EA,Dn	1001rrrmmmaaaaaa	xxxxx	dst dst-src	b,w,l
CMP	EA,Dn	1011rrrmmmaaaaaa	-xxxx	dst-src	b,w,l
CMPI	#dat,EA	00001100wwaaaaaa	-xxxx	dst-immed.data	b,w,l
MULS	EA,Dn	1100rrr111aaaaaa	-xx00	Dn Dn*src	1 w*w
DIVS	EA,Dn	1000rrr111aaaaaa	-xxx0	Dn Dn/src	1 l/w
AND	EA,Dn	1100rrrmmmaaaaaa	-xx00	dst dst src	b,w,l
OR	EA,Dn	1000rrrmmmaaaaaa	-xx00	dst dst src	b,w,l
EOR	EA,Dn	1011rrrmmmaaaaaa	-xx00	dst dst src	b,w,l
CLR	EAs	01000010wwaaaaaa	-0100	dst 0	b,w,l
NEG	EAs	01000100wwaaaaaa	xxxxx	dst 0-dst	b,w,l
TST	EAs	01001010wwaaaaaa	-xx00	dst-0	b,w,l
NOT	EA	01000110wwaaaaaa	-xx00	dst ¬dst	b,w,l

MC68000 Integer arithmetic and logic instructions[‡]

‡ Notes: rrr is a D register number.

mmm is a 3-bit mode field specifying the dst as EA or Dn, and operands as b, w, or 1:

Byte Word Long Destination

000 001 010 Dn

100 101 110 EA

EA is an effective address.

aaaaaa is a 6-bit address specifier. Not all modes are available to all instructions. See the manufacturer's literature for details.

ww is a word-size specifier field: 00-byte; 01-word; 10-long.

CMPI is followed by one or two words containing the immediate data to compare.

Mnemonic	Operands	Opcode Word [‡]	Operation				
Conditional instructions							
Bcc	disp	0110ccccddddddd DDDDDDDDDDDDDDDDD	if (cond) then PC PC+ disp				
DBcc	Dn,disp	0101cccc11001rrr DDDDDDDDDDDDDDDDD	<pre>if ¬(cond) then (Dn Dn-1 if Dn -1 then PC PC+ disp) else PC PC+ 2</pre>				
Scc	EA	0101cccc11aaaaaa	if (cond) then (EA) FFH else (EA) 00H				
Unconditional instructions							
BRA	disp	01100000ddddddd DDDDDDDDDDDDDDDD	PC PC+ disp				
BSR	disp	01100001ddddddd DDDDDDDDDDDDDDDDD	-(SP) PC; PC PC+ disp				
JMP	EA	0100111011aaaaaa	PC EA				
JSR	EA	0100111010aaaaaa	-(SP) PC; PC EA				
Subroutine return instructions							
RTR		0100111001110111	CC (SP)+; PC (SP)+				
RTS		0100111001110101	PC (SP)+				

MC68000 Program control instructions

‡ Notes:

rrr is one of An

If 8-bit displacement dddddddd is zero, then displacement is DDDDDDDDDDDDDDDDDD. EA is an effective address.

aaaaaa is a 6-bit effective address specifier. Not all addressing modes are available to all instructions. See the manufacturer's literature for details.

ww is a word-size field: 00-byte; 01-word; 10-long.

cccc is defined in the Table below:

Name	Meaning	Code	Logic	Name	Meaning	Code	Logic
Т	true	0000	1	F	false	0001	0
CC	carry clear	0100	С	LS	LS low or same		C+Z
CS	carry set	0101	С	LT	less than	1101	$N \cdot V + N \cdot V$
EQ	equal	0111	Z	MI	minus	1011	Ν
GE	greater or equal	1100	$N \cdot V + N \cdot V$	NE	not equal	0110	Ζ
GT	greater than	1110	$N \cdot V \cdot Z + N \cdot V \cdot Z$	PL	plus	1010	Ν
HI	high	0010	C·Z	VC	overflow clear	1000	V
LE	less or equal	1111	$N \cdot V + N \cdot V + Z$	VS	overflow set	1001	V

MC68000 Conditions