



Lecture 2

Review on Digital Logic (Part 1)

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<http://classes.engineering.wustl.edu/ese461/>

Grading



- Engagement 5%
- Review Quiz 10%
- Homework 10%
- Labs 40%
- Final Project 35%

- Policy:
 - 90% or above A
 - 80% - 89% B
 - 65% - 79% C
 - 45% - 64% D
 - 44% or below F

Number Systems



- Decimal (radix $r=10$)
 - digits 0-9
- Binary (radix $r=2$)
 - $(1011.01)_2 = (\quad . \quad)_{10}$
- Octal
 - radix = 8
- Hexadecimal
 - each HEX digit can represent 4 bits
 - $(10011110.0101)_2 = (\quad . \quad)_{16}$
 - $(11110010.0001)_2 = (\quad . \quad)_{16}$
 - $(11010100.1111)_2 = (\quad . \quad)_{16}$

Base Conversion



- Convert the integer part
- Convert the fraction part
- Join the two results with a radix point

- Example: $(325.64)_{10}$ to $(\quad . \quad)_5$
 - $2 \cdot 5^3 + 3 \cdot 5^2 = 325$
 - $3 \cdot 5^{-1} + 1 \cdot 5^{-2} = 0.64$

$$\sum A_i r_1^i = \sum B_i r_2^i$$

Range of Numbers



- Integer (n-bit Number)
 - 2^n different numbers
 - Min: 0
 - Max: $2^n - 1$
- Fraction (m-bit Number)
 - Min: 0
 - Max: $(2^m - 1) / 2^m$

- Diminished Radix Complement of N
 - defined as $(r^n - 1) - N$, with n = number of digits or bits
 - 1's complement for binary (radix = 2)
- Radix Complement
 - defined as $r^n - N$
 - 2's complement for binary
- Why
 - subtraction as addition of complement
 - if negative?

Binary Complement



- 1's complement
 - complement each individual bit (bitwise NOT)

- 2's complement
 - 1's complement plus 1
 - alternative
 - start from the least significant bit (LSB)
 - copy all least significant 0's
 - copy the first 1
 - complement all bits thereafter

Subtraction with 2's Complement



- For n-digit unsigned numbers M and N
- $M - N = ?$
 - add 2's complement of N to M
 - $M + (2^n - N) = M - N + 2^n$
- Example
 - carry 1
 - carry 0

Signed Binary Numbers



- To represent a sign (+ or -)
 - need one more bit
 - sign + magnitude
 - signed-complements
- Positive numbers unchanged
- Negative numbers use one of the two methods

- #	sign+	1's	2's
- +2	010	010	010
- -2	110	101	110
- +3	011	011	011
- -3	111	100	101
- +0	000	000	000
- -0	100	111	000

2's Complement Arithmetic



- Addition
 - represent negative number by its 2's complement
 - add the number including the sign bits
 - discard a carry out of the sign bits
 - e.g. $M + (-N) \rightarrow M + (2^n - N)$
- Subtraction
 - $M - N \rightarrow M + (2^n - N)$
 - form the complement of the subtrahend
 - follow the rules for addition

Overflow



- Error occurs when out of range
 - $[-2^{n-1}, 2^{n-1}-1]$
 - carry-in into the sign bit different from the carry-out

• $-39 + 92 = 53$:

$$\begin{array}{r}
 \boxed{1} \quad 1 \quad 1 \quad 1 \\
 \quad 1 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 0 \quad 1 \\
 + \quad 0 \quad 1 \quad 0 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0 \\
 \hline
 \quad 0 \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1
 \end{array}$$

Carryout without overflow. Sum is correct.

• $-19 + -7 = -26$:

$$\begin{array}{r}
 \boxed{1} \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 \quad 1 \quad 1 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \\
 + \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0 \quad 1 \\
 \hline
 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0 \quad 1 \quad 1 \quad 0
 \end{array}$$

Carryout without overflow. Sum is correct.

• $44 + 45 = 89$:

$$\begin{array}{r}
 \quad 1 \quad 1 \quad 1 \\
 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 0 \\
 + \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 0 \quad 1
 \end{array}$$

No overflow nor carryout.

• $104 + 45 = 149$:

$$\begin{array}{r}
 \quad 1 \quad 1 \quad 1 \\
 \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \quad 0 \quad 0 \quad 0 \\
 + \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 \quad 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1
 \end{array}$$

Overflow, no carryout. Sum is not correct.

• $-75 + 59 = -16$:

$$\begin{array}{r}
 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 \quad 1 \quad 0 \quad 1 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \\
 + \quad 0 \quad 0 \quad 1 \quad 1 \quad 1 \quad 0 \quad 1 \quad 1 \\
 \hline
 \quad 1 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0 \quad 0 \quad 0
 \end{array}$$

No overflow nor carryout.

• $-103 + -69 = -172$:

$$\begin{array}{r}
 \boxed{1} \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 \quad 1 \quad 0 \quad 0 \quad 1 \quad 1 \quad 0 \quad 0 \quad 1 \\
 + \quad 1 \quad 0 \quad 1 \quad 1 \quad 1 \quad 0 \quad 1 \quad 1 \\
 \hline
 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 0
 \end{array}$$

Overflow, with incidental carryout. Sum is not correct.

• $10 + -3 = 7$:

$$\begin{array}{r}
 \boxed{1} \quad 1 \quad 1 \quad 1 \quad 1 \\
 \quad 0 \quad 0 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \\
 + \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 0 \quad 1 \\
 \hline
 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 1 \quad 1 \quad 1
 \end{array}$$

Carryout without overflow. Sum is correct.

• $127 + 1 = 128$:

$$\begin{array}{r}
 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 \quad 0 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 + \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 1 \\
 \hline
 \quad 1 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0
 \end{array}$$

Overflow, no carryout. Sum is not correct.

• $-1 + 1 = 0$:

$$\begin{array}{r}
 \boxed{1} \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \\
 + \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 1 \\
 \hline
 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0 \quad 0
 \end{array}$$

Carryout without overflow. Sum is correct.

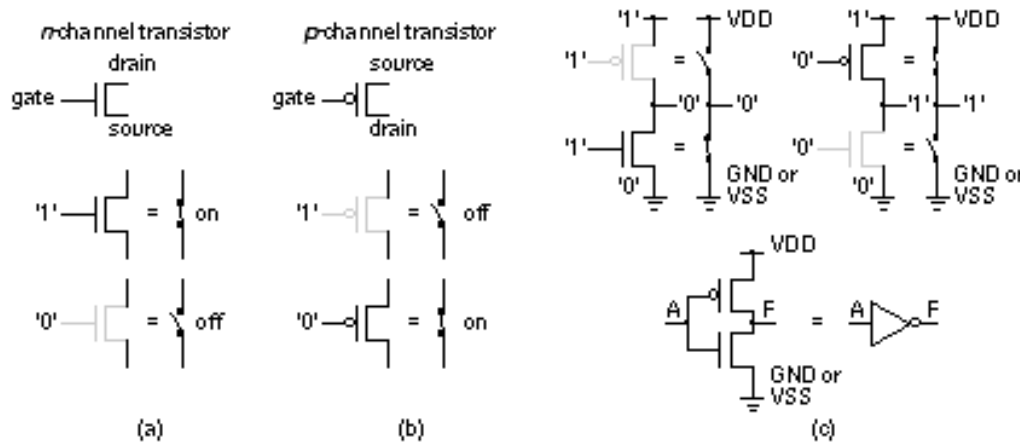


Number Representation

Boolean Logic and Gates

Combinational Logic

- Basic logic operations
 - AND: $X \cdot Y$
 - OR: $X + Y$
 - NOT: \bar{X} or $/X$
- Truth table
 - all possible input combinations
- CMOS as a switch



- Basic identities

1. $X + 0 = X$

3. $X + 1 = 1$

5. $X + X = X$

7. $X + \bar{X} = 1$

9. $\overline{\overline{X}} = X$ Involution

2. $X \cdot 1 = X$

4. $X \cdot 0 = 0$

6. $X \cdot X = X$

8. $X \cdot \bar{X} = 0$

} Existence 0 and 1 or operations with 0 and 1

} Idempotence

} Existence complements

Replace “+” by “.”, “.” by +,

Dual “0” by “1” and “1” by “0”

- Multiple Variables

10. $X + Y = Y + X$

Commutative

11. $XY = YX$

12. $(X + Y) + Z = X + (Y + Z)$

Associative

13. $(XY)Z = X(YZ)$

14. $X(Y + Z) = XY + XZ$

Distributive

15. $X + YZ = (X + Y)(X + Z)$

16. $\overline{X + Y} = \bar{X} \cdot \bar{Y}$

DeMorgan's

17. $\overline{X \cdot Y} = \bar{X} + \bar{Y}$

- Other useful theorems

$$XY + \bar{X}Y = Y$$

Minimization

Dual

$$(X + Y)(\bar{X} + Y) = Y$$

$$X + XY = X$$

Absorption

$$X(X + Y) = X$$

$$X + \bar{X}Y = X + Y$$

Simplification

$$X(\bar{X} + Y) = XY$$

$$XY + \bar{X}Z + YZ = XY + XZ$$

Consensus

$$(X + Y)(\bar{X} + Z)(Y + Z) = (X + Y)(\bar{X} + Z)$$

Standard (Canonical) Forms



- For comparison of equality
- Correspondence to the truth table
- Sum of Products (SOP)
 - sum of minterms
- Product of Sum (POS)
 - product of maxterms

Index	Minterm	Maxterm
0 (00)	$\bar{x}\bar{y}$	$x + y$
1 (01)	$\bar{x}y$	$x + \bar{y}$
2 (10)	$x\bar{y}$	$\bar{x} + y$
3 (11)	xy	$\bar{x} + \bar{y}$

- Relationship between min and MAX term

$$M_i = \bar{m}_i \quad m_i = \bar{M}_i$$

Circuit Optimization

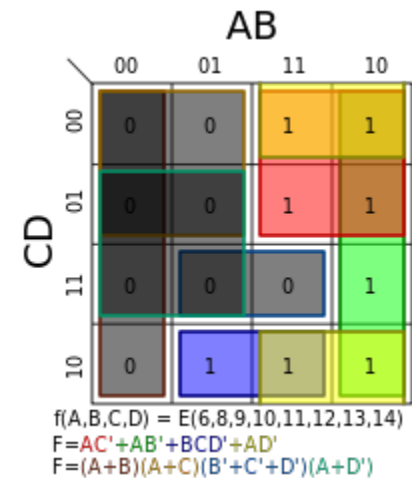
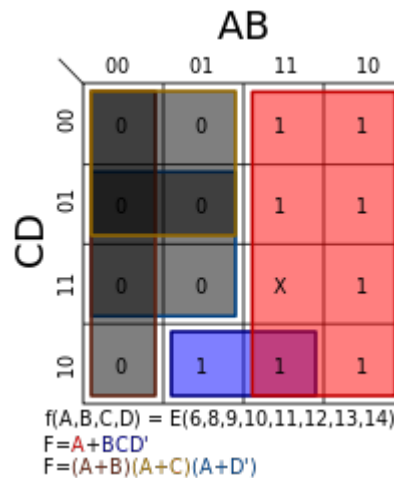
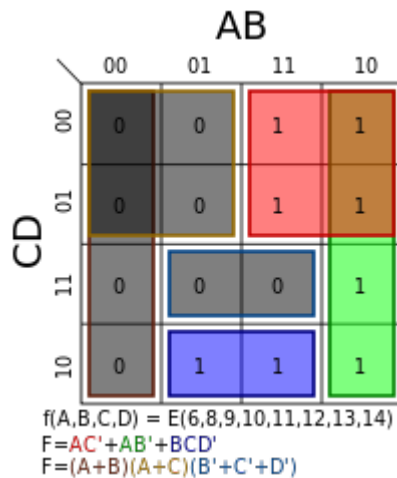


- Simplest implementation
- Cost criterion
 - literal cost (L)
 - gate input cost (G)
 - gate input cost with NOTs (GN)
- Examples (all the same function):
 - $F = BD + AB'C + AC'D'$ L =
 - $F = BD + AB'C + AB'D' + ABC'$ L =
 - $F = (A + B)(A + D)(B + C + D')(B' + C' + D)$ L =
 - Which solution is best?

Karnaugh Maps (K-map)








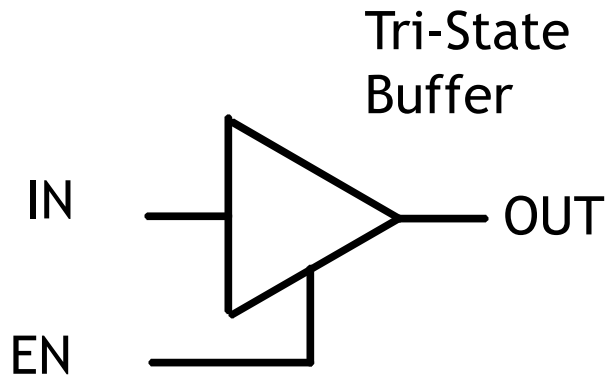
- Simplify Boolean algebra
- Transfer truth table to 2D grid
 - ordered in Gray code
- Human's pattern recognition capability
 - don't cares
 - race hazards



Other Gate Types



						
A	B	BUF	NAND	NOR	XOR	XNOR
0	0	0	1	1	0	1
0	1	0	1	0	1	0
1	0	1	1	0	1	0
1	1	1	0	0	0	1

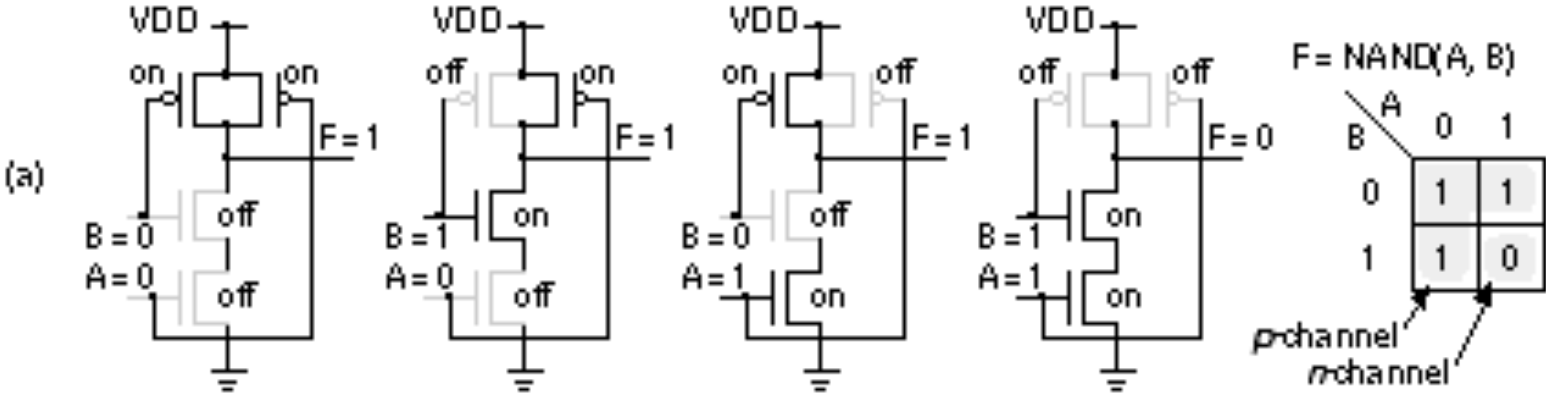


OUT = IN · EN

Truth Table

EN	IN	OUT
0	X	Hi-Z
1	0	0
1	1	1

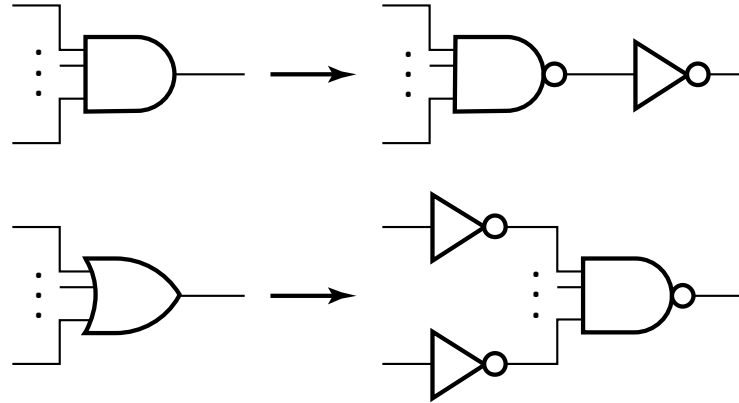
CMOS NAND and NOR Gates



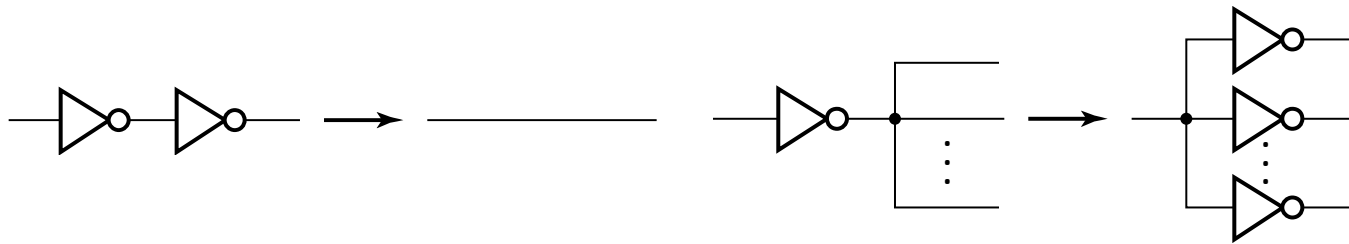
NAND Mapping Algorithm



- Replace ANDs and ORs



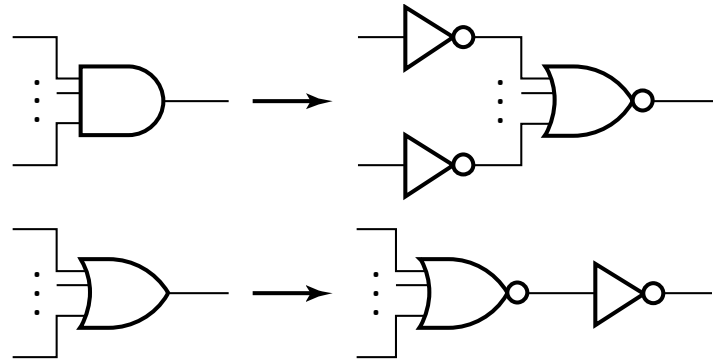
- Repeat until there is at most one inverter between
 - a circuit input or driving NAND gate output
 - the attached NAND gate inputs



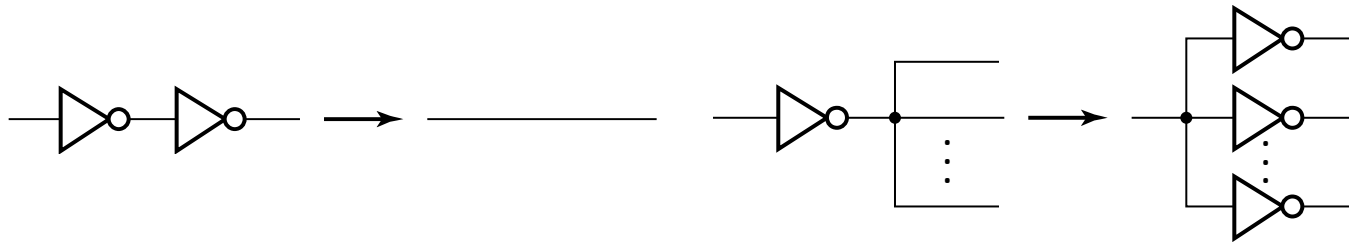
NOR Mapping Algorithm



- Replace ANDs and ORs



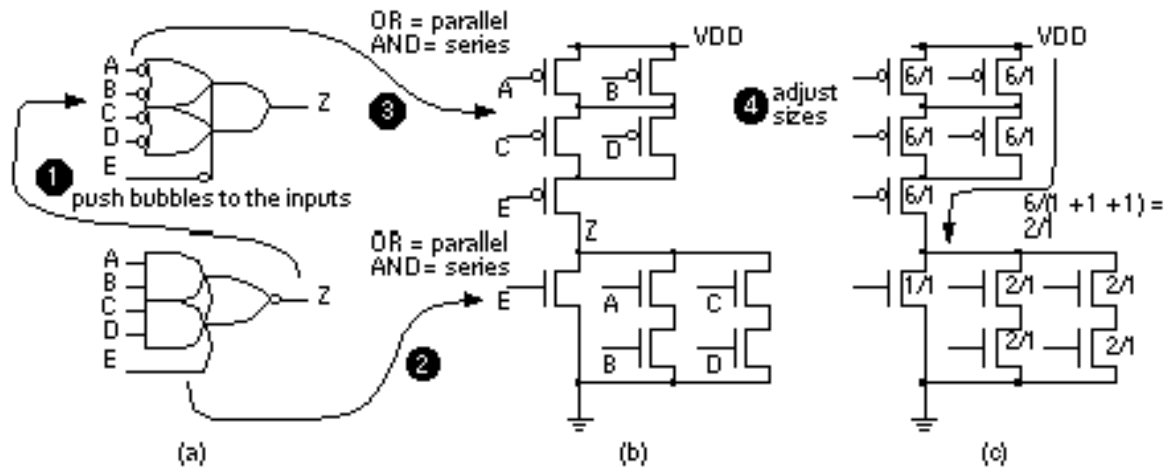
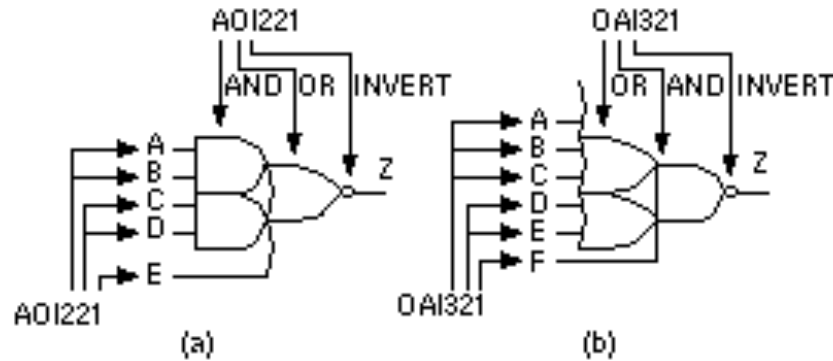
- Repeat until there is at most one inverter between
 - a circuit input or driving NOR gate output
 - the attached NOR gate inputs



Construct AOI and OAI Gates



- AND-OR-INVERT (AOI)
- OR-AND-INVERT (OAI)





Number Representation

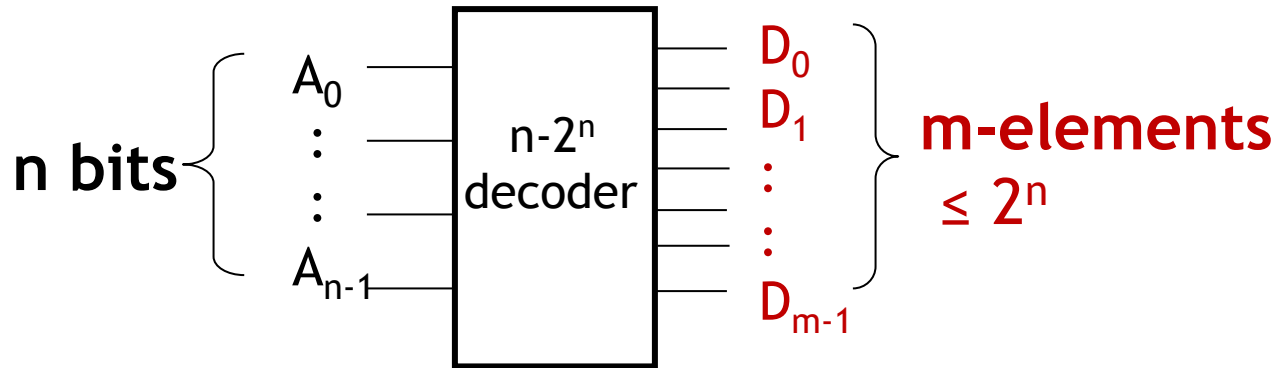
Boolean Logic and Gates

Combinational Logic

Decoding



- n-bit to represent up to $m=2^n$ elements
- convert n-bit input to m-bit output code
 - $n \leq m \leq 2^n$



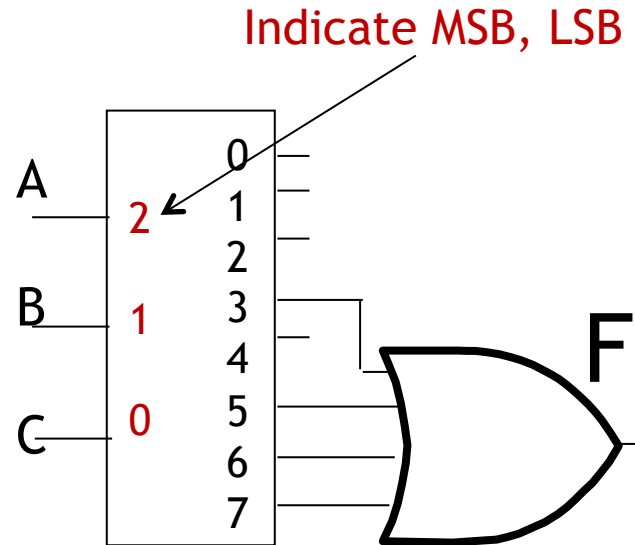
- outputs correspond to minterms: $D_i = m_i$
- divide into smaller decoders

Arbitrary Combinational Logic



- Decoder and OR gates
 - implement m functions of n variables
 - SOP expressions
 - one n -to- 2^n line decoder
 - m OR gates, one for each output

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1



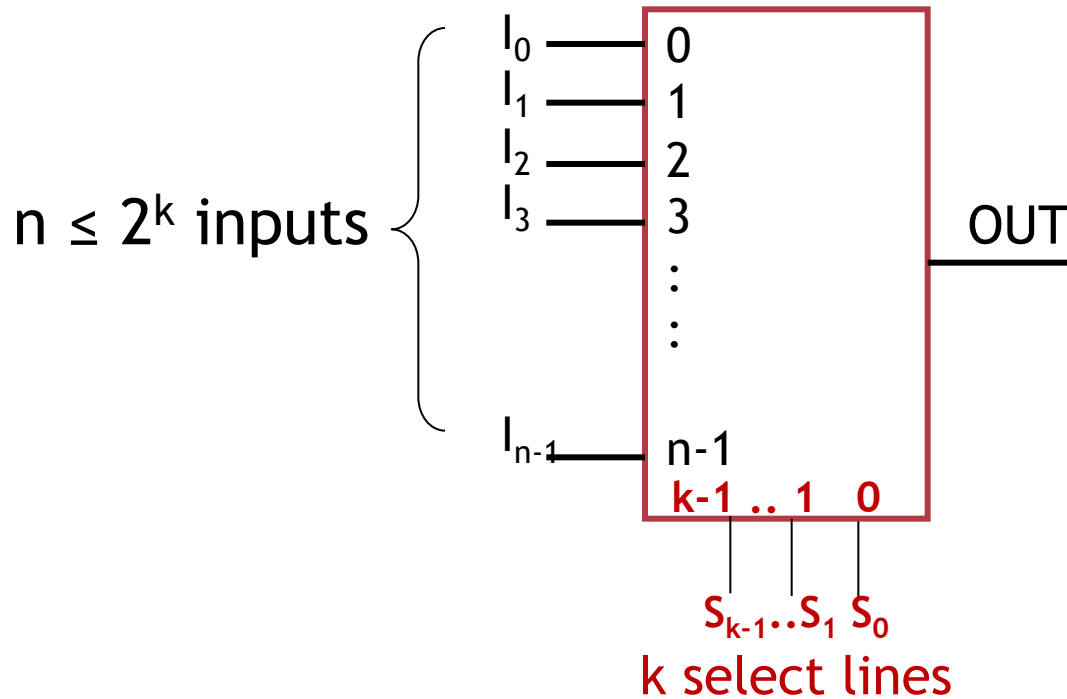
$$F = \sum m(3, 5, 6, 7)$$

- Convert one-hot code to its position
 - assume exactly one bit is 1
 - need to know its position
- Priority
 - more than one input value is 1

Multiplexing



- Select data
 - a set of n information inputs to select from
 - a set of k control (select) lines to make selection
 - a single output



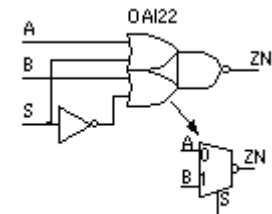
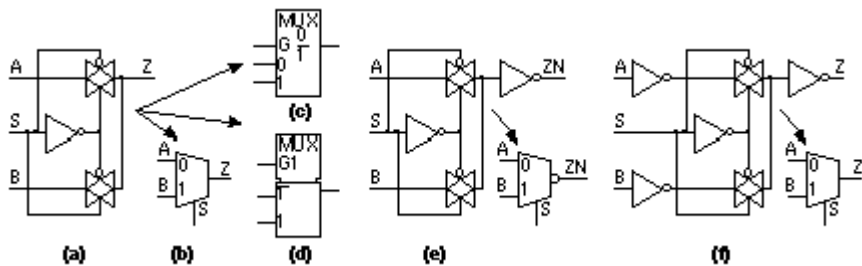
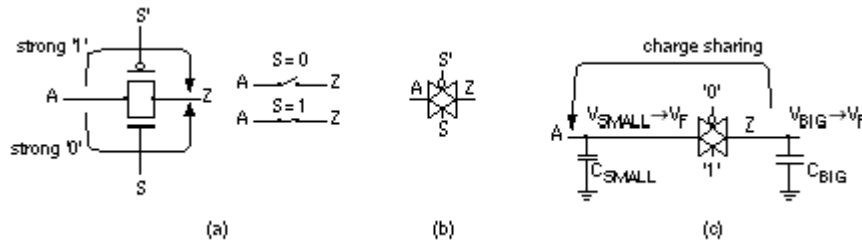
MUX Realization



- SOP Expression

$$\text{OUT} = \sum_{i=0}^{2^k-1} m_i I_i$$

- Transmission gates





Questions?

Comments?

Discussion?

Homework #1



- Download problem sets from class website
- Due 09/07 (Wednesday) in class
- No grace period